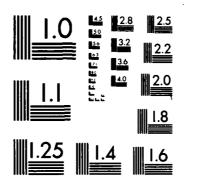
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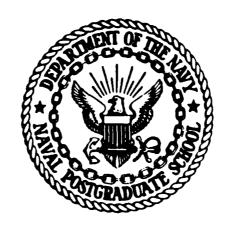


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NAVAL POSTGRADUATE SCHOOL

Monterey, California





THESIS

ADAMEASURE:
AN IMPLEMENTATION
OF THE HALSTEAD AND HENRY METRICS

bу

Paul M. Herzig

June 1987

Thesis Advisor:

Daniel L. Davis

Approved for public release; distribution is unlimited

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AdaMeasure An Ada Software Metric Implementation of the Henry Metric

bу

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ABSTRACT

A software metric is a tool that should be used in the development of quality software. The properties that define good software vary but encompass reliability, complexity, efficiency, testability, understandability, and modifiability. The Henry metric measures the complexity of data flow within a module and the complexity of inter-module communication. This thesis is an extension of a previous thesis titled 'AdaMeasure' that calculated the Halstead metric. The present design and implementation is a tool that computes the Halstead and Henry metrics for Ada programs.

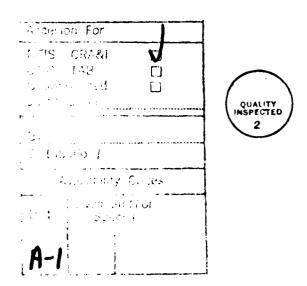


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I. INTRODUCTION AND BACKGROUND

A. DEFINITIONS

A metric is an assignment of indices of merit to programs in order to evaluate and predict software quality [Ref.1: p.6-2]. The qualities to measure are, at present, subjectively chosen but in general encompass reliability, complexity, efficiency, testability, understandability and modifiability [Ref.2: p.1-3]. The predictive nature of a metric allows it to be used to say "when" to proceed to the next phase in the software life cycle model. Another aspect of the predictive nature of a metric would be for it to provide management with a rough guess of the outcome of a particular path of development, provide an acceptance index, or provide an immediate feedback loop to the implementors while in the unit test phase [Ref.2: p.5]. How the metric is implemented will dictate its primary use from the above selections.

B. SALLIE HENRY'S METRIC

Sallie Henry's metric attempts to measures data flow complexity. It is intended to be used as a tool to establish a module's quality or to enforce particular modularization standards [Ref.2: p.6]. She argues that quality control of software is the result of software reliability and that reliability comes about through well designed modules that do not have complex data flow.

The hierarchical structure of a program should be layered modules. Each layer should function as a virtual machine and be composed of modules. This approach to modularization gives each module characteristics that can be exploited so that each module can be independently developed, more easily comprehended, assembled so that the system is more stable and designed so that the system is a great deal more flexible. This schema of development extols two primary tenets that are stated by D. Parnas in [Ref.3: p.339] and quoted here:

. . . provide the user of a module with all the information to use the module correctly, and nothing more. Provide the implementor of the module with all the information to implement the module correctly, and nothing more.

All this implies that a good design will have high module cohesion, good module strength and low module coupling [Ref.4: p.330].

C. INFORMATION FLOW

Information flow complexity is a twofold process the flow of data within a module and the flow of data external to the module. The measurement of these criteria is dependent on two premises: (1) that there is a capability to measure this data and (2) the data obtained can be used to evaluate software design. The seemingly obvious nature of the first premise runs into problems in implementation and applicability, but if it is accepted that the first deficiency can be surmounted, then the second part remains to be shown as reasonable. Applicability is a debated concept that is still not resolved. It revolves around whether the data gathered is related to the property under consideration. It is

further exacerbated by the human element that defines an environmental bubble and then programs within this bubble. How to measure this bubble without destroying its foundations is the problem of measuring human performance. The problem of what to measure is the problem of applicability.

The more specific the metric's application the less the applicability property is questioned, but, the problem of "what" to measure is still not clearly defined. This thesis will not argue the applicability question because the approach of Sallie Henry is reasonable and the results obtained from the metric appear to adequately encompass the area of data flow complexity. If the reader will accept that the properties measured are related to data flow complexity then the results obtained are also related to complexity.

The second premise is even more thorny. If the data is obtained and it seems reasonable can it be shown to be truly the result of the property under measurement? Any human endeavor will never be clearly and objectively quantified. Thus, the answer to the efficacy of the second premise is, proceed and maybe the amassing of results will eventually show the correlation.

The above analysis is far from a convincing argument to utilize metrics to measure programs however as this thesis was developed the applicability of measuring data flow complexity in order to determine code quality became more apparent although not proven. Nothing will be learned if no attempt is made to measure data flow complexity. This thesis attempts to measure data flow

complexity in the light of learning and the hope that the data gathered will prove the applicability of the process.

Consider first a simple module: a procedure in a structured language. Each procedure defines certain relations between itself and other procedures. These include:

- formal input/output parameters
- function call input and return data
- local data structures
- global data structures

These relations will generate a particular information flow structure similar to a hierarchical tree structure. This tree structure is peculiar to the procedure and will reflect its complexity of structure. It is reasonable to analyze this tree to determine derived calls, local data flow and global data flow.

D. RELATIONS

Some definitions are now in order. Global data flow exists from procedure 1 to procedure 2 if procedure 1 deposits data in the global data structure and then procedure 2 reads that data. Local data flow comprises direct and indirect species.

A local direct flow, from procedure 1 to procedure 2, results when procedure 1 calls 2 passing parameters. An indirect data exchange from procedure 1 to procedure 2 exists if procedure 2 calls 1, which returns a value used by 2, or procedure 3 calls both 1 and 2, and passes an output value from 1 to 2.

Figure 1 represents data flow from procedure to procedure or from a procedure into a data structure. Parameter passing within this scheme is represented by the arrows. A hidden data exchange through modification of a variable is represented by the dashed flow arrow. Module A retrieves data from the data structure then calls B passing a parameter; module B updates the data structure. C calls D passing a parameter. D calls E with a parameter and E returns a value to D which is used by D and passed to F. The function of F updates the data structure.

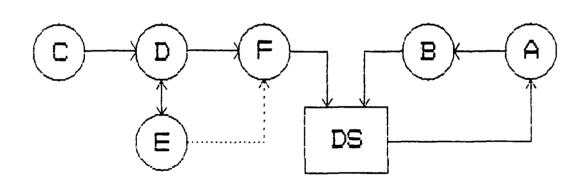


Figure 1. Data Flow Structure

The direct data flows represented are:

A -> B, C -> D, D -> E, D -> F.

These are simply the calls.

The indirect local flows are:

 $E \rightarrow D, F \rightarrow A.$

The global flows are:

 $B \rightarrow A, F \rightarrow A.$

Both B and F update the data structure while A retrieves data from the structure.

The implications of data flow for procedure and function calls will be discussed later with derived calls.

The calling notation $A(x) \to B()$ or $A() \to B(x)$ is used to connote a data flow transmission from A to B either by direct parameter passing or side-effect. In the first condition the variable x is returned to procedure A and in the second example the variable x is sent to B. A condition that leads the Henry metric to not detect a procedure or function call's data flow (labeled a missed call) is for the condition where $A(x) \to B()$ and variable x is a returned value from B not modified within procedure A's code. An example of this would be a conditional statement within A that depends on the returned value from function B. The data flow detection problem leads to two key ideas, effective parameters and data utilization.

Calls that are detected by information flow analysis are dependent upon how the information is passed. If the conditions $A() \rightarrow B(x)$ exists where parameter x

is passed to B or condition A() -> B() where no parameters are exchanged then the calls will not be missed if B receives information in one of the following formats:

- a formal parameter
- a data structure
- a constant
- an actual parameter from a third procedure whose value is modified within A prior to the call to B

An effective parameter will define the call structure in such a way that the data flow will not be missed. It is a parameter that receives information from one of the calling procedure's parameters, a data structure, a constant, or a third procedure's returned actual parameter that is modified within the calling modules structure. What the effective parameter implies is that side- effect data flow is difficult to effectively analyze. Another construct that will cause a missed call is the condition $A(z) \rightarrow B(x)$ where B is a function. This condition means A uses data from B. A uses data from B if (1) B updates a data structure used by A; (2) A receives a constant from B; (3) A receives an output parameter from B; or (4) B updates a return value to A. Thus information flow will be detected if A passes B an effective parameter or if A uses data from B.

Appendix A gives all the rules that are applicable to the data flow relationships. Some notation is now needed to simplify the descriptions that follow. The form of a relation is L <- R1, R2, R3,...Rn; where L is the resultant from the application of the relationships R1, R2, ... Rn. An example would be:

A.D3 < -A.D1, A.D2, A.constant.

This series notation represents the code line that begins with D3 below.

A()
begin
.
.
.
D3 := D1 + D2 + 1;
end procedure A;

In words, the A.D3 means procedure A updates data structure D3 by first applying relationship A.D1 then A.D2 and finally A.constant. This format shows that data flows into procedure A's data structure D3 from the noted relationships. A thorough discussion of the notation for the relations is given in Appendix A but a short discussion follows to aide in the immediate understanding of Figure 2.

The notation B.1.I defines the first input parameter in the actual parameter list of procedure B and an O would refer to an output parameter. All possible data flow paths are considered even if a B.1.I parameter is not an input parameter. Thus, if procedure B has an output actual parameter in position B.1 and the Henry metric attempts to analyze this parameter as an input flow an error condition would result from the attempted evaluation (depicted as B.ERROR). B.NULL means that no relationship exits for this parameter or that there is no data flow into or out of the parameter being considered.

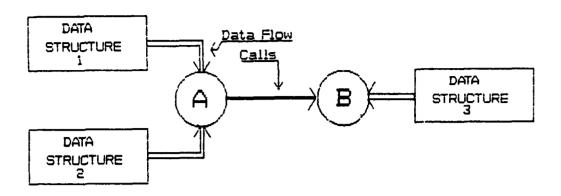


Figure 2. Data Flow With Call Structure

Code

```
A()
begin
    X := D1 + 1;
    Y := D2;
    B(X,Y);
end;
```

```
B(P,Q)
begin
D3 := P + Q;
end;
```

Relations Set

A1 B.1.I < -A.D1, A.CONSTANT

A2 B.2.I < - A.D2

B1 B.1.O <- B.NULL

B2 B.2.0 <- B.NULL

B3 B.D3 <- B.1.I, B.2.I

The relation sets were derived by looking at the data flow into and out of procedure A. That is, since procedure A has no parameters there can only be local data flows into or out of the procedure. These flows are described in terms of the procedure call to B. B.1.I stands for procedure B's first input parameter. This parameter is fed from procedure A's data structure D1 and a constant. Analyzing procedure B's second input parameter yields the A2 relationship. Relationship B1 describes the first parameter in procedure B as an output parameter to procedure A that receives no data for transfer. Relationship B3 describes how the two input parameters to procedure B constitute the data flow to this data structure.

The data flow analysis deals primarily with the analysis of parameters which are direct data flow and indirect data flow as defined above. Modifying Figure 2 and incorporating some local variables will illustrate some more data flow analysis techniques as seen in Figure 3.

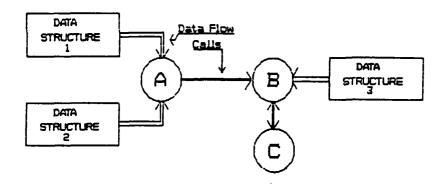


Figure 3. Data Flow And Inter-dependent Procedures

```
A()
begin
 X := D1 + 1;
 Y := D2;
 B(X,Y)
end;
B(P,Q)
begin
R := Q;
C(P,R,S);
D3 := S;
end;
C(I,J,K)
begin
 K := I + J;
 J := J + 1;
end;
```

Code

Relation Set

A1, A2 same.

```
B1 B.1.O <- C.1.O
B2 B.2.O <- B.NULL
B3 C.1.I <- B.1.I
B4 C.2.I <- B.2.I
B5 C.3.I <- B.ERROR
B6 B.D3 <- C.3.O
```

```
C1 C.1.I <- C.NULL
C2 C.2.O <- C.2.I, C.CONSTANT
C3 C.3.O <- C.1.I, C.2.I
```

In the relation set B1 receives data from procedure C's output parameter. B2 is the same. B3 through B5 describe the parameter list of procedure C. However B5 denotes an error or a condition that is not allowed. That is, the data direction was in error as variable S is an output from procedure C as indicated by relation B6. It should be noted that this relation set building considers all possible data flow paths without regard to the possibility that the parameters could be assigned only particular directions as Ada formal parameters are. Figure 4 shows the effects of a function call.

Code

```
A()
begin
    X := D1 + 1;
    Y := F(X);
    B(X,Y)
end;
```

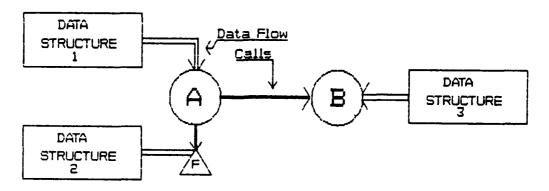


Figure 4. Data Flow With Function Call

```
F(M) return integer;
begin
N := D2 * M;
return N;
end;
```

F2 F.O <- F.D2, F.1.I

Relation Sets are changed as follows:

```
A1 F.1.I <- A.D1, A.CONSTANT
A2 B.1.I <- A.D1, A.CONSTANT, F.1.O
A3 B.2.I <- F.O
F1 F.1.O <- F.NULL
```

Relation A1 has changed to reflect the analysis of the function call to F. The input to the function call is analyzed as well as its output and any possible modification of its input parameter. This analysis can be seen to cover all possibilities of hidden data transfers except the missed calls described earlier.

E. INFORMATION FLOW STRUCTURE

Once the relation set has been built the relations are sorted alphabetically and stored for future use in the Information Flow Structure (IFS). A recursive algorithm is employed to build the information tree structure for the flow analysis. The IFS is then analyzed to find the derived calls, the local flows and, finally, the global flows.

The IFS will have leaves that are data structures; the root is the initial call from the highest level procedure. Each node of the tree will have the relational form of X.DS, X.O. X.k.I, or X.k.O. See Appendix A for all the possibilities of derived calls. The local flows are described in Appendix A as derived calls. The global flows for a particular data structure are all the possible paths from leaf elements of the form A.DS to the root.

F. INDICES OF MERIT

The calculations of the indices of merit use the idea that the complexity of a module comprises the complexity of the code plus the complexity of the connections of the code to other modules. The formula describing the complexity of a module is

Complexity = length * (fan-in * fan-out) ^ code index.

Length is defined as the number of executable statements. The expression fan-in ' fan-out represents all the combinatorial possibilities for each input to produce an output. The code index is an exponent that represents the code

difficulty. Nominal code difficulty for operating systems is 2. This index needs more data for other types of programming.

The purpose of this computation is to produce comparative numbers of merit that point out and isolate specific areas within the code that have the potential for problems. A high fan-in/fan-out implies a large interconnection to outside modules. This leads to the assessment that the code in question is most likely not properly modularized or, more succinctly, that the code has more than one function. The other form of data flow is global data flow to data structures. It is calculated as follows:

The term write refers to a change to the data within the structure through an assignment statement and a read is an access to the data structure that does not change the data. The identifier read-write is the sum of reads and writes. A high global flow implies overworked data structures and represents a stress point in the program. A stress point is the weak link in the chain. The presence of high flow is not automatically an indicator of poor programming but it is a juncture in the program that is highly susceptible to problems. Once the metric has assembled all the different components, such as fan-in or global reads and calculated the above equations it performs module analysis.

G. MODULE ANALYSIS

Module analysis revolves on the outputs of each of the above equations and their respective components. The numbers generated are symptomatic of certain problems. The analysis is first conducted with the equations output defining the particular categories of problems then the components refine the analysis. Examples of the first level of analysis follow:

A high global flow calculation implies an overworked data structure. These structures are overworked because of the need for continuous accessing. This implies a better decentralized design is in order, that is, distribute the information to the procedures that it serves. A high module complexity index indicates not enough modularization. This number is to be treated with respect but should be analyzed in context with global data flow. Together these indices represent the in's and out's of the modules data. A corrective action based solely on complexity should be avoided. A procedure should be analyzed for singularity of purpose and non-duplication within a module. Simply put, a procedure should be in one place, have one purpose and have minimal external references. These properties are quantified by the Complexity and Global flow metric numbers.

Next the interim cases where one aspect is high and the other component is low. A module with high global flow and low complexity shows poor internal structure. This structure will most likely have excessive numbers of procedures with extensive use of data structures outside the module. Low global flow with

high module complexity implies either poor decomposition into procedures or extremely complicated interface.

H. INTERFACE MEASUREMENTS

Interface between procedures comprise protocol interface, coupling and binding of procedures. Protocol interface from module A to B is defined as those procedures that are not in any other module and which receive information from A for passing to B. Binding is the sensitivity measure between modules, that is, tightly bound modules have a high sensitivity. A tightly bound module is difficult to change without adversely affecting the other module. Coupling is the strength of binding. Figure 5 depicts the interface structure.

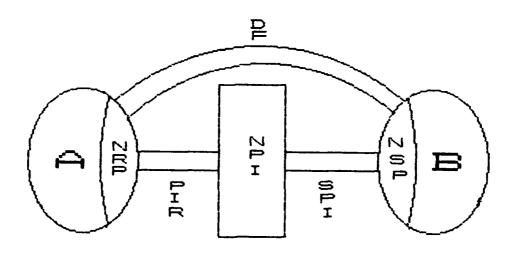


Figure 5. Interface Structure

Protocol interface, since it is not symmetrical between procedures, requires the construction of a tabular cross reference table with all possible procedures on both the X and Y axes. The internals of the table are the data flow complexity indexes for one way transmission from each procedure to the other.

Figure 5 shows that binding is sectioned into five components; the number of procedures sending information from A (NSP), the number of procedures receiving information from B (NRP), the number of procedures in the protocol interface (NPI), the number of paths to the interface from A (SPI), and the number of paths from the interface to B (PIR). The Direct Flow paths represented as the outer loops are data transmissions without the interim procedures. [Ref.2: p.85] lists the binding calculations as follows:

Binding =
$$(NSP + NPI) * SPI + (NPI + NRP) * PIR$$

The term (NSP + NPI) * SPI is the coupling strength.

All the direct path binding is calculated by

$$DF Binding = (NSP + NRP) \cdot DF$$

Modules that are tightly bound are extremely difficult to maintain and modify. This difficulty stems from their lack of independence and the "ripple effect" of changes to one module flowing into the other.

I. THEORY SUMMARY

The purpose of the Henry metric is to provide designers and implementors with a method to quantify the quality of the code that they are developing. The goal is to produce reliable code that is interconnected in as logical a fashion as possible. Information flow complexity produces reliability through enforcement of design rules that lead to well connected code. The measurements will point out lack of functionality, improper modularization, poorly designed modules, poorly designed data structures, system stress points, inadequate refinement, strength of binding, modifiability, missing levels of abstraction and, will produce comparative indices to assess changes.

II. DESIGN CRITERIA

A. INTRODUCTION

The design of the implementation of the Henry metric was a two step process. First, a thorough understanding of the previous work by Neider and Fairbanks [Ref.5: pp.1-164] was undertaken to determine what data were available for importation into the Henry metric. The intention of the study and the basic design issues are (1) modify the output of the parser portion of their thesis to include the necessary data passes to the Henry metric (2) to initially analyze only the Ada package as a unit (3) encapsulate as much of the Henry metric into one Ada package as possible and (4) calculate the necessary Henry metric numbers transparently to the user but present the user with an output that is easily understood.

The underlying premise, of this program, is that the code presented for analysis has been successfully compiled. If the code does not compile and is presented to the parsers it will most likely fail to parse but in the event it does escape detection it will be erroneously analyzed.

The design criteria of encapsulation of the Henry metric was modified during the implementation due to the unwieldy length of the code. The division yielded three packages: one that holds all the global data, another for the analysis portion and a third for the interface to the user. Although this division violated one of the basic design issues it was necessary in order to achieve solid data transfers between the Naval Post Graduate School computer and the Naval Weapons Station computers.

While the level of what was available from the parsers was being determined the data structures of the Henry metric were being layed out. The data structures and data gathering procedures for the Henry metric were designed to be as simple, yet as flexible as possible. Once the layout of the linked list data structure to hold the raw information for the Henry metric was decided upon, the tedious procedure of inserting the appropriate calls to the Henry package was undertaken. Basically, this reduced to exercising all the possible data flow characteristics from function or procedure calls that the Henry metric could expect to encounter and then ensuring that an appropriate call to the Henry data collection procedure was placed in the Neider/Fairbanks parser.

Next, the analysis procedure was designed. The analysis was separated from the display because computers have very different capabilities in their output devices. The first package analyzes the collected raw data and the second displays the finished, smooth data.

The program is menu driven. It initially gives the user a choice to parse a new program or view old data. If the parse choice is selected the parsers feed both the Henry and the Halstead packages with data. After a successful parse the user is presented with a choice of either viewing the Halstead or Henry metric data. This is when the analysis portion of the Henry metric is called. It is not

until the user decides which data to view is there a distinction made as to how to process the parsed data. This feature was designed for several reasons:

- Use the Halstead metric to refine the code
- then analyze the code via the Henry metric for data flow
- but the user should still have the option to select either metric

The data for both metrics are essentially different as are the purposes for gathering the data but the final goal of this dual metric system is to produce good code.

Finally, the presentation data module was designed. The Halstead and the Henry metric both produce numbers which are essentially meaningless unless a thorough understanding of the particular metric implementation is undertaken. Thus, both metrics, in differing fashions present "help" data to aid understanding of the metric output. These files are both verbally and graphically presented to the user.

The overriding design issue was to modularize the Henry metric as much as possible. The most significant exception to the Parnas' ideal [Ref.4: p.330] are the numerous calls to the data gathering procedure from the parser modules. These calls depend on details of how the data will be analyzed in their sequencing and data passing scheme. A conscientious effort was made to minimize global data and isolate procedures into nearly stand alone modules.

B. SPECIFIC ISSUES

The housekeeping routines of the global package are used to adjust the parsed data into a more palatable form for the analysis and presentation procedures. The output from the parser is stored as a linked list of raw data produced by procedures WRITE_HENRY_DATA and CREATE_NODE. This data is then analyzed by the ANALYSIS PACKAGE for the particular data constructs that represent a data flow. The output is an array of tabulated data that is a set of all the relations necessary for the Henry metric to detect local and global data flow.

The display module presents data in a tabular format or as a graphical representation for relative merit analysis. The intent was for the user to see the effect of changes, or to select a more verbose description of the meaning of the results. The modules that accomplish tabular and graphic displays are separated again because of the varying capabilities of machines. The purpose of these procedures is to provide some form of relative measure to the user so that the improvement or results of a change could be more objectively weighed. The overall purpose of the display modules is to show the data in such a fashion that an intelligent assessment is possible.

C. DESIGN ISSUES

Design issues encountered in the implementation of Henry's metric involve the efficient use of the Ada language's structures and data analysis techniques. Sallie Henry developed a metric process in which the constructs of a particular language are ignored or not put to specific use. That is, some languages have extremely thorough type and range checking facilities. This is not considered in the basic design of her metric. These powerful features are incorporated in Ada and provide the application programmer more analysis capability.

This design issue concerns the collection of 'all possible paths' data for analysis of actual parameters. The approach taken by Henry is biased to a language where input, output, and combination input output parameters are treated as if they could be modified by the particular procedure regardless of their type. Ada is very picky about the manipulation of formal and actual parameters and goes to great lengths to ensure that parameter consistency is maintained by means of strong type/range checking. The explicit declaration of a parameter type was used to select which component of the complexity equation should be updated. The appropriate fan-in or fan-out number was also correctly updated from default declarations such as the undeclared default formal parameter.

The data analysis technique issue encountered was the need to analyze the data via the IFS. Henry's IFS was designed so that a traversal of its nodes analyzing parent-child pairs will capture all the transitive relational data flows. The transitive flow analysis designed into the present parser will account for the first two layers. The reasoning behind this approach stems from a program review. This review, albeit not extensive, was conducted looking for the predominant use of transitive relations. The review revealed that transitivity is not often used and if used is at most two layers deep. There was little use of deep

TO SOURCE MANAGEMENT

transitive constructs. Thus, the design approach selected will detect the majority of the transitive data flow paths without the need for an extensive tree structure. The "normal" program has few transitive relations but the capability to analyze this style of program would add more accuracy to the metric.

Another design anomally of the metric is the problem of detecting the difference between a function call outside the package declaration and a global data structure manipulation. Ada libraries or packages inhibits the proper analysis of a function call as opposed to a data structure read unless a full compiler's output is available. The present metric was designed so that local function calls (within the package being analyzed) are properly valued but the function calls outside the package are treated as data structure manipulations or more specifically as global data flows.

D. CONCLUSION

The design and implementation phase was driven by the analysis of the Neider/Fairbanks parser portion of their thesis followed by the modularization of the Henry metric. The tradeoffs considered were: the strong typing and range checking of the Ada language, the need for an information flow tree, the need for relative output for the user and, and most importantly the desire to incorporate all of the Henry metric into one Ada package.

III. DESIGN AND TESTING

A. THE EMBEDDED CODE

The previous work done by Neider/Fairbanks had to be modified to output the necessary data for the Henry metric. This was accomplished through embedding calls to the Write Henry Data procedure in Parser0, Parser1, Parser2, Parser3 and Bypass Function (See Appendix C). The writing of the Lexeme, or identifier's name, was controlled by a Boolean that was turned on or off according to the position of the parse of a particular package. The design criteria was to keep the data gathering as simple as possible. If time permitted, a more thorough and sophisticated scheme could be developed. The embedded code was thoroughly tested by two test harnesses that simulated a series of Current Token Records in the form of an input Ada package.

B. THE HENRY PACKAGE

The first package to be implemented was Henry.pkg. It was conceived to be a stand-alone construction that would initialize the data collection process, receive data from the other parsing packages and store the raw incoming data in a linked list. (See Appendix B). Minimal variables and foreign procedures from other packages are used. The Henry package's only "withed" packages are TEXT IO, HENRY GLOBAL. HENRY ANALYSIS—and—HENRY DISPLAY.—This

approach was considered necessary so that the subsequent changes or upgrades would not affect other modules (ripple effect). The design was to implement a basic Henry metric first for Ada packages then to improve and more fully develop the Henry analysis techniques if time allowed. The Main Menu module sequences the user into the analysis and display support packages. The modularization was considered necessary because the analysis and display packages are separate entities and the separation will ensure maintainability.

The initialization is conducted by procedure Initialize Henry and the declaration statements that assign initial values to various Boolean variables. Initialize Henry creates two head nodes, one for the raw data linked list, the other for the procedure or function length records. The raw data linked list storage is a straight line of Henry records. These records have five fields that identify whether this is (1) local or global declaration. (2) the variable/procedure's name. (3) an action class. (4) a parameter class and (5) a pointer to the next record. The action class is comprised of various identifiers that range from procedure type to end parameters declare. Their purpose is to delineate the actions within the parsed program so that the Henry analysis package can look for the data flow. The parameter type field is used to define input, output or combination input output formal parameters. The variable Henry Line count is purposely initialized within this procedure to draw attention to it's initial value. The array of length records is initially a parallel construct not directly tied to the procedure or function it holds the data for. In the analysis package a sequential process is produced where the records are linked to the data manipulation array. The purpose of the length record is to hold the begin and end line counts of each procedure or function. These line counts are used later to compute the specific modules length for inclusion in the complexity equation.

The receipt incoming accomplished primarily Write Henry Data. This procedure supported boole**a**n is Write Henry Enable. This boolean turns on or off the recording of the incoming records from the Get Current Token Record procedure. Specifically, the boolean will allow recording only selected data from the incoming record stream selected by the place within the recursive descent parser that the boolean is activated. This control is necessary to pick and choose the data that is critical and to ignore the remainder.

The procedures Create Node and Clear Henry Lexeme support the data gathering scheme. The "in out" pointers within Create Node serve the purpose of allowing a view of the last record in the incoming stream or to work on the current record. It is arranged so that New Node points to the newly created blank record and Last Record points to the just filled in trailing record. Procedure Clear Henry Lexeme is necessary because of the way Ada handles strings. Create Line Node procedure functions identically to Create Node.

The incoming data is chosen from within the Bypass Function and from Parser0 to Parser4 [Ref.5: pp.102-160] by where the calls to the Write henry data procedure is positioned. The purpose of this approach is to assure the Henry

metric receives sufficient information but more importantly that the records writen into the linked list are delimited in a particular fashion for ease of analysis. There is still considerable data that can be collected for analysis from the parsers but the Henry metric is not to the stage of development where it would be useful. The added depth of information could be used in two areas; analysis and a more informative output from both metrics.

The Write Henry Data procedure selectively enters the field data into the raw data linked list records as dictated by the incoming actual parameters. That is, the incoming data has default settings but if the data is to be ignored then the "null setting" is passed as an actual parameter. This assists in the gathering process. The design of the data gathering modules is such that modifications could be easily implemented. This was purposely designed into them so that upgrades would be fairly painless.

The Henry pkg was constructed with modularization and maintainability in mind. It was meant to be a stand alone entity that receives data from the Neider/Fairbanks Bypass Function and Parser packages. It performs the functions of initialization, data receipt and data storage besides defining the data structures used throughout the Henry metric packages. There are a number of improvements that could be added to the actual parameter analysis. These improvements all concern the wealth of options Ada provides in parameter passing schemes, such as aggregates, dot notation to access hidden variables and allocators. Further, the present Henry metric does not analyze the incoming

actual parameters for expressions but the variables are all considered for inclusion in the complexity calculation by the transitivity analysis.

C. HENRY ANALYSIS PACKAGE

The Henry Analysis package comprises three procedures to set up the raw linked list data and a fourth procedure to actually analyze the code for metric calculations. The Analysis procedures are called sequentially from the Henry package and function as support for the Henry package. They operate on the data in sequential discrete steps. They first determine the formal parameters, then search and identify procedures and variables and then determine the metric numbers. The approach used was to nibble each piece of the tremendously complex data flow calculation down into minute sub-steps until all that is left is to simply count the marks on each record for determination of the complexity or global flow metric numbers. This approach removed the necessity for an arduous single pass calculation.

The set up procedures are CLEAN_UP_HENRY_DATA.

SET UP HENRY ARRAY, and SPRUCE_UP_HENRY_DATA. A support function, LOCAL NAME, assists in the setting-up process. These procedures end product are two metrics, the complexity metric and the global flow metric.

The Clean Up procedure ensures that all parameter type records have all their fields properly filled. It scans for their parameter lists all the procedures and functions that are declared in the analyzed package. The field of most

importance is the classification of either "in, out or in-out" type parameters.

These fields are checked up to the colon delimiter within the formal parameter list and then entered into all parameter type records correctly.

The Set_Up procedure scans through the entire linked list setting up another array of pointers to facilitate the analysis process. The Henry_Array records have identifier, beginning pointer and line_length record pointer entries. This procedure's purpose is to break up the long linked list into another array. It actually does not sub-divide the list it merely arranges an array of pointers into the linked list that delineate each function or procedure. The resulting array is called the Henry_Array. The line length record pointers are records that hold the stop and stop line numbers. These records are eventually used to compute procedure/function lengths.

The Spruce_Up procedure goes through the Henry array data and sorts out the local and global data flow paths. It does this through the use of the LOCAL NAME function. This function searches either the Henry array for a particular procedure name or the package and appropriate procedure's declaration sections for the variable name in question. Its purpose is to sort out the local procedure or function calls from the global data structure manipulations. It cannot completely solve this problem but defers final resolution to the Calculation procedure.

The Calculate Metric procedure will again process the Henry array data looking for the final resolution to local procedure or function calls as opposed to a

global data structure manipulation. It proceeds in small increments to finally arrive at the complexity metric calculation and a global data flow calculation. The complexity metric number is arrived at by first considering all the in, out, in-out formal parameters to calculate the fan-in and fan-out numbers. After the initial cut the fan-in, fan-out numbers are incremented upward by the numbers of identified procedure actual parameters that feed these formal parameters and then by the the Transitivity In and Transitivity Out functions.

An example of this process would be for procedure A with formal parameters X. Y. First process parameters X and Y for their explicit type adding 1 to fan-in if its an input parameter or 1 to fan-out if its an output parameter. Next process all the assignment expressions looking for a modification of the formal parameter. If procedure A modifies parameter X prior to a call to another function increment the fan-out count by the number of statements after the assignment delimiter. Then go through an analysis of transitivity incrementing fan-in or fan-out accordingly. Finally, call up the appropriate record of Henry Line count and calculate the length of the procedure or function in question.

The equation that the process is working toward solving is:

This equation represents the local data flows within the analyzed procedure. Sallie Henry set the exponent of the bracketed expression to 2 because of her experience with operating system code analysis. This program will continue with

this number until enough data can be compiled to support a change. Once this calculation is done then the global data flows are analyzed.

The global flows are arrived at by first eliminating all other possibilities. Then the remaining choices have to be foreign data flows. This process is started in the Spruce_Up procedure and completed within the Metric_Calculation procedure. The process is used to find whether the data structure is being read from or written to or both.

The equation that the analysis is striving to solve is:

```
Global flow = write * (read + read-write) +
read-write * (read + read-write - 1)
```

This equation represents how and by what means the global data structures are manipulated. The global data analysis procedure goes across procedure or function boundaries whereas the previous complexity metric calculations remain within the particular procedure or function under scrutiny. This across- the-border calculation is accomplished through the text file that is discussed next.

Within the calculation procedure the initial entries for the display package are started. This amounts to constructing a text file of descriptive terms and indices of merit for output in the Display Package. It also provide a temporary storage bank for the global data information. This across-boundary analysis of

global flows was necessary because of the implications of not being able to detect the difference within the Ada code of an access to a data structure or a function call to a "withed" package.

In summary, the Analysis package is a series of analytical steps. The purpose of these steps is to arrive at the complexity and global flow metrics. These indices and additional data are stored in a text file for output to the user within the Display package.

D. HENRY DISPLAY PACKAGE

The Henry Display package is the user interface portion of the metric program. It provides the user with four different aspects of viewing the analyzed data. The purpose of this package is to show the user the data flow characteristics of the particular parsed input program. The output data will be the fan-in, fan-out, length, complexity, and four global flow numbers. These numbers can be presented in a listing format, viewed with a help file of informative paragraphs or compared by means of the other portions of the analyzed package to gain a relative sense of merit.

The procedures that comprise the Display package are LIST_METRIC_DATA and WRITE_RELATIVE_DATA and GRAPH_RELATIVE. The LIST_METRIC_DATA procedure will output the data file compiled while in the calculation portion of the previously discussed package. It will be a straight listing of information that will be grouped by each

element in the calculation of the complexity or global flow numbers, such as all procedures are grouped under the head of FAN-IN. The purpose of this listing is to show each procedure or functions component figure in the calculation of the final complexity and global figures. If the programmer is in a compile, test, edit, recompile mode of operation this will provide a spotlight on where to improve the data flow "choke- points". These data flow critical points will be seen as either global high flow complexity high numbers. In short. the LIST METRIC DATA is designed for a more sophisticated programmer wishing to edit-and-run and see the results of particular programming style changes.

The WRITE_RELATIVE_DATA display will provide the same format of data but the numbers will have been normalized. Accompanying each number set will be a short narrative keyed to the relationships of the particular numbers. That is, if the user sees a complexity number of 125 beside the procedure X he will be provided with an explanation that that number is not too far out of line in comparison to the other procedures or functions analyzed within this package. The purpose of this approach is to normalize the output numbers to provide a relative comparison for a more user friendly approach to the mysterious metric number generation.

There is an additional procedure within the Display package that provides a complete listing of the raw input data. This procedure will most likely be of no use to anyone except those programmers who are extremely interested in the factors that lead to the particular numbers presented.

The final package for viewing the data is the graphical presentation module. It takes the relative data and manipulates the floating point numbers to achieve a bar chart display.

In summary, the Display package will provide the user with information in a variety of formats so that he can reach a conclusion from relative merit or absolute input numbers. The data flow numbers will point out the critical data flow points within a procedure so that the programmer can better see where to improve or expend the most effort. The purpose of the output data is to show the user where to improve, not how to improve.

E. TESTING

The testing of the design was conducted as the modules were being built and at the integration step prior to the final product. This was accomplished through the use of test harnesses that simulated the particular module's inputs and through test input programs that were hand analyzed to verify the metric's outputs.

The testing of the Henry package was accomplished by gradually building a more thorough test harness as each previous test was successful. The final test harness encompassed over 200 input records that simulated a myriad of token record inputs. The testing of the Henry module presented some difficulty because it is so intimately tied to the parsers. This was overcome by simulating the Bypass Support package as a partial input and the test harness as the balance of

the test vehicle. The package performed well within the test harness and functions adequately within the context of the entire program.

The testing of the analysis package of the Henry metric again was an iterative build of the harness. The testing accomplished after all the Henry metric packages were integrated was accomplished on the same group of programs provided by the NWC programmers to test the Halstead metric integration. The harness testing was comprised of a 50 step program that simulated a package with three independent procedures/functions utilized within its scope. There were intentional references outside the scope of the test harness package to determine if the global call detection scheme functioned properly. In all, the test harness exercised every possible data flow scheme analyzable by the Henry metric including one that would be a missed call. The analysis package performs adequately within the scope of the harness. Testing revealed that the code within the analysis phase was non-reentrant which required a the use of a boolean to define the status of the call to the package. This boolean will protect the data structure and effectively make the code reentrant.

page personal addition excellent establish and analysis and and

The display module was tested with the same driver harness as the analysis package. The results were used to fine tune the package and to debug the problems. The process used was to call the analysis package from the display package and drive the display package with the test harness. This is also how the integrated program performs. The results were adequate from the stand point of the test harness but need some refinement when using the whole program.

The summary of testing would be extensive use of complicated test harnesses. Since a test harness has to simulate all the inputs to a module that the tested code could possibly see during integration, they are difficult to build much less debug. The debugging problem comes from the question 'is it the code or is it the harness?'. The test harness approach is quite fruitful from two orientations: (1) it forces the programmer into a thoroughly understanding his code and (2) the harness construction will lead the programmer into optimizing his code. Why doesn't the programmer already understand his code? He does but the ramifications of a certain approach does not come surface until the design of a test harness is considered. The optimization is driven by the need to get accurate, fast results so that the troubleshoot-repair-compile-troubleshoot regimen can proceed fairly rapidly. This is a real concern with the tremendous

IV. CONCLUSIONS

A. IMPLEMENTATION

STREET STREET STREET STREET STREET STREET

The Henry metric was implemented in as modularized a fashion as possible. The intent was to allow for improvements through a more thorough use of the parser's information in AdaMeasure's first revision. Also, certain aspects of the Henry metric were not implemented, but it is now felt that they would add depth to the analysis process. In particular, the first change should be that the Information Flow Structure be added. This tree-like structure will allow the analysis of hidden calls but will still not detect the missed call problem discussed earlier. The missed call problem will most likely only be solved through the use of the Program Counter Register, but this approach defeats the idea of a high level language. The final improvement would be to add analysis of the "withed" packages so that an interface table could be constructed.

The program was incorporated into the previous work by Neider and Fairbanks. Their work was extensive and deserves favorable mention because it made the implementation of the Henry metric considerably easier. The output of the program is still in need of sophistication and improvement. In particular, two improvements are needed: (1) explaining the theory behind the metric and (2) conveying the ideas to the user. For an example, a high global data flow indicates an overworked global data structure. What should the metric present to the user?

The average programmer might not see the relevance of this and would miss the indication that a critical point in the data flow should probably be revised.

B. THE FUTURE OF METRICS

Metrics are tools. They point out areas of weakness. The metric will show a direction to proceed even in the absence of an absolute answer as to the correctness of the analyzed code.

The importance of metrics will grow as the size of programs grow. We do not know how important metrics will become but it does seem clear that there is a need for something that helps improve code quality and is fairly painless to use. The emphasis on "good" code will continue to be in the forefront of the Armed Service's concerns because of their intense involvement in real time embedded programs. These programs present a real challenge for incorporation of changes, improvements or any other form of maintenance programming. The purpose of metrics in this environment would be to point the way to good modularized design.

The metric should be part of the test scenario besides being an integral member of the life cycle of the program. The metric will force quality control without the painful process of formal inspections. The formal process has its place but the metric tool could perform more than the inspection. The metric tool should be incorporated into the test cycle as a meter of improvement. This immediate feedback to the programmer will be beneficial. The manager could

also use the absolute number as a goal for acceptance. This approach will provide the manager with the data needed at decision points in the life cycle of a program. The absolute number could also be tied to the program throughout its life as a measure of improvement or degradation over time. The uses are many, as the reader can see. The importance of the metric cannot be overstated when the future holds programs that will span millions of lines of code.

Metrics are important. They hold out the hope of an automated tool that will guide, interpret, and assess progress for programmers and management alike. I hope that the work of this metric will assist in advancing metrics and the use of the Ada language.

APPENDIX A: INFORMATION FLOW MECHANISMS

MECHANISMS FOR INFORMATION FLOW ANALYSIS

As Sallie Henry so succintly states:

The information flow analysis takes place in three phases. The first phase involves generating a set of relations indicating the flow of information through input parameters, output parameters, returned values from functions and data structures.

General Format of a Relation

The generation of relations is first prefaced with a quick review of relational format.

L - R1, R2 ... Rcount:

Where L may be in any one of the following forms:

- 1. P.DS P is the name of a procedure and DS the data structure.
- 2. P.O. P is the procedure name and O is the return value.
- 3. P.j.O P is the procedure j is an integer representing the formal parameter position, and O is the jth Output parameter.
- 4. P.j.I P is the procedure, j is an integer representing the jth parameter, and I is the jth input parameter.

Ri may be in one of the following forms:

- 1. S.DS S is a procedure name and DS is the name of a data structure.
- 2. S.O. S is a procedure name and O is the returned value.
- 3. S.j.I S is the procedure name, j is the jth parameter and I is the jth input parameter.
- 4. S.j.O. S is the procedure name, j is an integer representing the jth parameter in the list and O is the output parameter.
- 5. Smull S is the procedure name, null represents no data

- 6. S.cons. S is the procedure name and constant a value used within S.
- 7. Serror S is the procedure name and error represents an invalid flow of information through procedure S.

RULES

- L is of the form P.DS then this form is used only to generate the relations from procedure P that updates DS with Ri.
- L is of the form P.O then
 This is used only in generating the relations from procedure P that produce an output.
- 3. L is of the form P.j.O then
 This is used when generating the relations that produce an input of the jth parameter in the procedure's formal parameter list. There must be a unique relation for each of P's parameters.
- 4. L is of the form P.j.I then
 This is used when generating the relations for procedure P
 that produce an input for the jth parameter. Another
 procedure T calls P to indicate that the jth parameter of P
 receives the input update.
- 5. Ri is of the form S.DS then
 Procedure S reads information from DS this format is used to
 indicate a read only.
- 6. Ri is of the form S.O then
 Relations are generated that come from procedure T that are
 return values to T from S.
- 7. Ri is of the form S.j.I then
 For generating relations for procedure S that indicates S's
 jth input parameter passes information to L.
- Ri is of the form S.cons.
 Then S causes a constant number or string to flow to L.
- 9. Ri is of the form S.Null then
 This is used to indicate when S does not update a parameter, that is, the parameter was strictly input only.
- 10. Ri is of the form S.error then S calls T and one of the parameters to T is an output only thus if S attempts to input a value this would be an error.

ANALYSIS OF CALLS

The following two procedures X and Y, exhibit all possible calling structures in the light of information flow analysis. NP stands for not possible, NC for no calls and the numbers beneath are used for later reference.

The pairs 1, 3, 5, 7, 13, and 15 cannot appear in a flow of data path because for DS's the only assignment and reads allowed are from procedures or functions. The other not possibles stem from input parameters not flowing into DS's and not flowing into output parameters. Entries 2 and 4 indicate X calling Y, receiving information from Y and using this information to update a DS. The rest of the possibilities can be reasoned in like manner except entries 10 and 12, which represent calls via a third procedure. Here procedure Z calls Y and passes the returned value from Y to X. This represents a no call between X and Y but there is a data flow.

TABLE 1.

LOCAL CALL TABLE					
	x.DS	x.o	X.k.I	X.K.O	
Y.DS	NP 1	NP 5	Y calls X 9	NP 13	
Y,O	X calls Y	X calls Y	NC 10	X calls Y	
Y.K.I	NP 3	NP 7	Y calls X	NP 15	
Y.K.O	X calls Y	X calls Y 8	NC 12	X calls Y 16	

All data flows from the highest calling structure, eventually being deposited in the data structures. Analysis of the table's data confirms this premise.

MEMORYLESS PROCEDURES

Some procedures keep no record of their data passing or the data supplied. These procedures are used to do housekeeping for memory management, for example, but their analysis for data flow would produce a false amount of data transactions. Another area that these procedures appear in are arithmetic operations that are sometimes duplicated in hardware such as double precision math etc. This discussion leads to the problem that these procedures would be difficult to discern in an automated process. That is, if memoryless procedures are not to be considered in data flow analysis some form of human decision making is required. It should be noted that this is another premise that the automation of the Henry metric is based on. The absolute numbers for the Henry metrics would

TABLE 2.

	 	GLOBAL CAL	L TABLE	
	X.DS	x.o	X.K.I	x.k.o
Y.DS	NP 1	NP 5	Y flows X	NP 13
Y.O	Yflows X	Yflows X	Y flows X 10	Y flows X
Y.K.I	NP 3	NP 7	Y flows X	NP 15
Y.K.O	Y flows X	Y flows X 8	Y flows X	Y flows X 16

be inflated if memoryless procedures are not eliminated from the analysis. In short a memoryless procedure should be removed from the code to be analyzed if a more accurate assessment or if the absolute numbers produced are being used for a comparative study.

APPENDIX B: HENRY METRIC CODE

```
-- TITLE:
              AN ADA SOFTWARE METRIC
 MODULE NAME: PACKAGE HENRY GLOBAL
-- DATE CREATED: 09 MAY 87
-- LAST MODIFIED: 19 MAY 87
-- AUTHOR:
               LCDR PAUL M. HERZIG
-- DESCRIPTION: This package contains the data declarations
    and basic procedures used throughout the Henry metric.
with GLOBAL, TEXT IO;
use GLOBAL TEXT 10;
package HENRY GLOBAL is
 package INTEGER IO is new TEXT IO.INTEGER IO(INTEGER);
 use INTEGER 10;
 package REAL IO is new TEXT IO.FLOAT IO(FLOAT);
 use REAL 10;
-- Real IO produces floating point output
MAX ARRAY SIZE : constant integer := 50;
MAX LINE SIZE : constant integer := 76;
DUMMY9s
            : constant integer := 9999;
             : constant character := '';
NULL CHAR
--DUMMY9s are used for false data input to the line length calculations
type DECLARED TYPE is (BLANK, LOCAL DECLARE, GLOBAL DECLARE);
type ACTION TYPE is (UNDEFINED,
               HENRY HEAD NODE.
               PACKAGE TYPE,
               PROCEDURE TYPE.
               FUNCTION TYPE,
               PARAM TYPE.
               ASSIGN TYPE,
               IDENT TYPE.
               DATA STRUCTURE.
               FUNCALL OR DS.
```

```
PROCALL OR DS.
              END PARAM DECLARE.
              END ACTUAL PARAM,
               END DECLARATIONS,
              END ASSIGN TYPE,
               END PACKAGE DECLARE,
              END PACKAGE TYPE,
               END FUNCTION TYPE,
              END PROCEDURE CALL):
type PARAM CLASS
                    is (NONE, IN TYPE, OUT TYPE, IN OUT TYPE,
              ACTUAL PARAM);
subtype FORMAL PARAM CLASS is PARAM CLASS range IN TYPE. IN OUT TYPE;
subtype LEXEME TYPE is string (1.. MAX LINE SIZE);
subtype END UNITS is ACTION TYPE range
             END FUNCTION TYPE..END PROCEDURE CALL:
--Declared, action and parameter classes or types are used
--in the Henry record data collection process
type HENRY RECORD;
type POINTER is access HENRY RECORD;
type HENRY RECORD is record
  IDENTITY
             : DECLARED TYPE:
   NOMEN
             : LEXEME TYPE:
  TYPE DEFINE : ACTION TYPE:
   PARAM TYPE : PARAM CLASS:
            : POINTER;
   NEXT1
end record:
--Henry record is the workhorse storage medium
type HENRY LINE COUNT RECORD:
type LINE POINTER is access HENRY LINE COUNT RECORD;
type HENRY LINE COUNT RECORD is record
   ID NAME
                : LEXEME TYPE:
   START COUNT
                    : INTEGER;
  STOP COUNT
                   : INTEGER;
   NEXT REC
                  : LINE POINTER;
end record:
--Henry line count record is used to claculate the length of procedures
--or functions
type HENRY DATA is record
 NAME OF DATA
                     : LEXEME TYPE:
 BEGIN POINTER
                     : POINTER;
 LINE LENGTH POINTER: LINE POINTER:
end record:
```

--Henry data records are used to delineate the functions and procedures

```
--for easier data calculations
```

```
type HENRY DATA ARRAY is array (1...MAX ARRAY SIZE) of HENRY DATA:
type OUTPUT DATA is record
 TYPE OF
             : ACTION TYPE := UNDEFINED;
 NAME OF
              : LEXEMĒ TYPE;
 TYPE FAN IN : FLOAT := 0.0;
 TYPE FAN OUT : FLOAT := 0.0;
 TYPE COMPLEXITY: FLOAT := 0.0;
 TYPE READ : FLOAT := 0.0:
 TYPEWRITE : FLOAT := 0.0;
 TYPE\_READ\_WRITE : FLOAT := 0.0;

TYPE\_FLOW : FLOAT := 0.0;
 CODE LENGTH : INTEGER := 0;
end record,
--Output data records hold the final calculation numbers for storage into
-- an output 'input file
type OUTPUT ARRAY is array (1...MAX ARRAY SIZE) of OUTPUT DATA;
NEXT HEN. LAST RECORD, NEW RECORD,
HEAD, NAME POINTER
                             : POINTER:
HENRY ARRAY
                          : HENRY DATA ARRAY:
HENRY LINE COUNT
                            : integer := 0;
OUT PUT DATA
                           : OUTPUT ARRAY;
LINE COUNT RECORD
                            : HENRY LINE COUNT RECORD;
HEAD LINE, NEXT LINE, LAST LINE : LINE POINTER;
PACKAGE BODY DECLARE,
ASSIGN MARKER.
GLOBAL MARKER,
NAME TAIL SET.
ASSIGN STATEMENT.
FUNCTION PARAM DECLARE.
                                : BOOLEAN := FALSE:
FORMAL PARAM DECLARE
                            : BOOLEAN := TRUE;
FIRST HENRY CALL
DUMMY LEXEME
                           : LEXEME TYPE:
procedure CREATE NODE(NEW NODE, LAST RECORD : in out POINTER);
procedure CREATE LINE COUNT NODE(NEXT_LINE,
                 LAST LINE : in out LINE POINTER):
procedure INITIALIZE HENRY(HEAD : in out POINTER:
              HEAD LINE in out LINE POINTER);
procedure CLEAR HENRY LEXEME(HENRY LEXEME : in out LEXEME TYPE);
end HENRY GLOBAL:
```

```
package body HENRY GLOBAL is
--procedure creates Henry record nodes for data storage
procedure CREATE NODE(NEW NODE, LAST RECORD: in out POINTER) is
TEMP POINTER: POINTER;
begin
 put(result file, "in create henry node"); new line(result file);
 TEMP POINTER := new HENRY RECORD;
 TEMP POINTER.IDENTITY := BLANK;
 for I in 1..MAX LINE SIZE loop
 TEMP POINTER.NOMEN(I) := NULL CHAR;
 end loop;
 TEMP POINTER. TYPE DEFINE := UNDEFINED:
 TEMP POINTER PARAM TYPE := NONE;
 NEW_NODE.NEXT1 := TEMP_POINTER:
 LAST RECORD
                     := NEW NODE;
NEW NODE := TEMP POINTER;
end CREATE NODE;
--creates line count nodes to hold the length data for each procedure or
--function
procedure CREATE LINE COUNT NODE(NEXT LINE,
                  LAST LINE : in out LINE POINTER) is
TEMP POINTER: LINE POINTER;
 put(result file, "in henry create line node"); new line(result file):
 TEMP POINTER
                       := new HENRY LINE COUNT RECORD;
 for I in 1...MAX LINE SIZE loop
 TEMP POINTER.ID NAME(I) := NULL CHAR:
 end loop:
 TEMP POINTER.START COUNT := DUMMY9s;
 TEMP POINTER.STOP COUNT := DUMMY9s;
 NEXT_LINE.NEXT_REC := TEMP_POINTER;
NEXT_LINE := NEXT_LINE;
NEXT_LINE := TEMP_DOD:
                   := TEMP POINTER;
end CREATE LINE COUNT NODE:
--sets all of the variables to their initial values besides
--creating the first Henry record and line count record
procedure INITIALIZE HENRY(HEAD : in out POINTER:
               HEAD LINE: in out LINE POINTER) is
```

```
HEAD STRING: STRING(1..9) := "HEAD NODE";
                : INTEGER := 9;
begin
 CREATE(HENRY FILE, out file, HENRY FILE NAME);
 put(HENRY FILE, "in INITIALIZE HENRY"); new line(HENRY FILE);
 CREATE(HENRY OUT, out file, HENRY OUT NAME):
             := new HENRY RECORD:
 HEAD.NOMEN(1..SIZE) := HEAD.STRING;
 HEAD IDENTITY := BLANK;
 HEAD.TYPE DEFINE := HENRY HEAD NODE;
 HEAD PARAM TYPE := NONE;
 NEXT HEN
                := HEAD;
 CREATE NODE(NEXT HEN. LAST RECORD);
 HENRY LINE COUNT := 0;
 DUMMY LEXEME(1) := NULL CHAR;
 HEAD LINE .= new HENRY LINE COUNT RECORD:
 HEAD LINE.ID NAME(1..SIZE) := HEAD STRING:
 HEAD LINE.START COUNT := DUMMY9s;
 HEAD LINE STOP COUNT
                          := DUMMY9s;
 NEXT LINE
                    := HEAD LINE;
 CREATE LINE COUNT NODE(NEXT LINE, LAST LINE);
end INITIALIZE HENRY:
--clears the input string to null characters
procedure CLEAR HENRY LEXEME(HENRY LEXEME : in out LEXEME TYPE) is
begin
put(HENRY FILE, "IN CLEAR HENRY LEXEME"); NEW LINE(HENRY FILE);
FOR I in 1 .. MAX LINE SIZE loop
HENRY LEXEME(I) := NULL CHAR:
end loop;
END CLEAR HENRY LEXEME:
END HENRY GLOBAL;
-- TITLE: AN ADA SOFTWARE METRIC
-- MODULE NAME: PACKAGE HENRY METRIC
-- DATE CREATED: 06 APR 87
-- LAST MODIFIED: 15 MAY 87
-- AUTHORS LCDR PAUL M. HERZIG
```

```
-- DESCRIPTION: This package contains the Henry metric data
         collection and program control routines.
with GLOBAL, HENRY GLOBAL, HENRY ANALYSIS, HENRY DISPLAY, TEXT 10;
use GLOBAL, HENRY GLOBAL, HENRY ANALYSIS, HENRY DISPLAY, TEXT 10;
package HENRY is
procedure WRITE HENRY DATA(ID : in DECLARED TYPE := BLANK;
              IN NAME: in LEXEME TYPE := DUMMY LEXEME;
              DEFINE : in ACTION TYPE := UNDEFINED;
              PARAM : in PARAM CLASS := NONE;
              LINK : in POINTER);
procedure UPDATE LINE COUNT:
procedure WRITE LINE COUNT(IN NAME: in LEXEME TYPE:= DUMMY LEXEME:
               FIRST COUNT : in INTEGER := DUMMY9s:
               LAST COUNT : in INTEGER := DUMMY9s:
               PTR
                       : in LINE POINTER):
end HENRY:
.,---------
package body HENRY is
--produces the written data records from the parser inputs
-data is only written if it is something other than the
-- null settings
procedure WRITE HENRY DATA(ID)
                                  in DECLARED TYPE BLANK
              IN NAME on LEXEME TYPE DUMMY LEXEME.
                                         UNDEFINED.
              DEFINE in ACTION TYPE
              PARAM in PARAM CLASS
                                         NONE
              LINK
                    in POINTER) IS
put(result file "in write henry data") new line(result file)
 If ID
             BLANK then
  LINK IDENTITY ID
```

```
case ID is
      when LOCAL DECLARE - put(RESULT FILE, "Local declare").
     when GLOBAL DECLARE - put (RESULT FILE, "Global declare")
                      put(RESULT_FILE, "Undeclared").
      when others
    end case
 else put(RESULT_FILE, "NO DECLARATION")
 end if
 new line(RESULT FILE)
 If IN NAME(i)
                 NULL CHAR then
  LINK NOMEN(1 MAX LINE SIZE)
                                  IN NAME(I MAX LINE SIZE).
  PUT(RESULT FILE IN NAME).
 ELSE PUT(RESULT_FILE, "NO NAME").
 end if.
 new line(RESULT FILE).
 If DEFINE
              UNDEFINED then
  LINK TYPE DEFINE DEFINE:
case DEFINE is
when UNDEFINED
                        put(RESULT_FILE, "Undefined")
when HENRY HEAD NODE
                             put(RESULT FILE, "Henry Head Node")
when PACKAGE TYPE
                           put(RESULT FILE, "Package declaration")
when PROCEDURE TYPE
                            put(RESULT_FILE, "Procedure declaration")
when FUNCTION TYPE
                           put(RESULT_FILE, "Function declaration").
when PARAM TYPE
                         put(RESULT FILE, "Parameter declaration").
when ASSIGN TYPE
                         put(RESULT_FILE, "Assignment delimiter").
when IDENT TYPE
                       put(RESULT_FILE_"Identifier").
when DATA STRUCTURE
                            put(RESULT FILE, "Data structure descriptor").
when FUNCALL OR DS
                           put(RESULT_FILE, "Function or data descriptor")
                           put(RESULT_FILE, "Procedure or data descriptor")
when PROCALL OR DS
when END PARAM DECLARE put(RESULT FILE, "End parameter delimiter")
when END ACTUAL PARAM
                             put(RESULT FILE, "End actual parameter delimiter")
when END DECLARATIONS
                             putiRESULT FILE. "End declaration delimiter"
                           put(RESULT_FILE, "End assignment statement delimiter").
when END ASSIGN TYPE
when END PACKAGE DECLARE
                                 put(RESULT_FILE, "End package declaration delimiter").
when END PACKAGE TYPE
                           - put(RESULT FILE, "End package delimiter")
when END FUNCTION TYPE
                             put(RESULT_FILE, "End function delimiter").
when END PROCEDURE CALL - put(RESULT FILE "End procedure delimiter")
when others
                    put(RESULT_FILE, "Unknown")
end rase
new line(RESULT FILE)
end if
If PARAM
               NONE then
  LINK PARAM TYPE
                        PARAM
  CASE PARAMIS
                    PUT(RESULT FILE "IN PARAM")
  WHEN IN TYPE
                      PUT(RESULT FILE "OUT PARAM")
  WHEN OUT TYPE
  WHEN IN OUT TYPE
                         PUTARESULT FIFE "IN OUT PARAM"
  WHEN OTHERS PUTCRESULE FILE "NONE")
 FNDCASE
 end it
 new line RESULT FILE)
→ TWRITE HENRY DATA
```

```
--increments the line count for eventual inclusion into
-- the calculation of a particular procedures total length
-- the length number is used in the complexity calculation
procedure UPDATE LINE COUNT is
begin
put(result file, "in update line count"); new line(result file);
if not FORMAL PARAM DECLARE then
 HENRY LINE COUNT : - HENRY LINE COUNT + 1;
 end if:
end UPDATE LINE COUNT:
--produces the records to hold the line count information
-- the records are not initially tied to a particular procedure
-- but are a parallel data structure until in the Hen anal.pkg
--where they are linked to the procedure that they hold the
--data for
procedure WRITE LINE COUNT(IN NAME : in LEXEME TYPE: DUMMY LEXEME:
                 FIRST COUNT : in INTEGER := DUMMY9s;
                 LAST COUNT : in INTEGER :=: DUMMY9s;
                          in LINE POINTER) IS
begin
 put(HENRY FILE, "in WRITE LINE COUNT"); new line(HENRY FILE);
put(result file. "in write line count"); new line(result file);
 If IN NAME(1)
                    NULL CHAR then
           PTR ID NAME(1 MAX LINE SIZE)
                                                 IN NAME; end if.
                     DUMMY9s then PTR.START COUNT FIRST COUNT, end if.
 If FIRST COUNT
 If LAST COUNT
                     DUMMY9s then PTR.STOP COUNT : LAST COUNT; end if.
end WRITE LINE COUNT.
end HENRY.
               AN ADA SOFTWARE METRIC
  TITLE
... MODULE NAME - PACKAGE HENRY ANALYSIS
```

```
-- DATE CREATED 29 APR 87
-- LAST MODIFIED 29 MAY 87
-- AUTHOR LCDR PAUL M. HERZIG
-- DESCRIPTION: This package contains the analysis functions
-- required to identify each data flow in the Henry metric
```

with GLOBAL GLOBAL PARSER, BYPASS SUPPORT FUNCTIONS, HENRY GLOBAL TEXT TO use GLOBAL GLOBAL PARSER, BYPASS SUPPORT FUNCTIONS, HENRY GLOBAL TEXT TO:

package HENRY ANALYSIS is

package NEW INTEGER IO is new TEXT IO INTEGER IO(integer) use NEW INTEGER IO.

package REAL IO is new TEXT IO FLOAT IO(float) use REAL IO:

PROC FUNC COUNT INTEGER 0.
INDEX INTEGER:
NAME POINTER POINTER.

--PROC FUNC COUNT is the total number of procedures and functions in the --analyzed package

type SELECTOR TYPE is (PROCEDURE FIND, FUNCTION FIND VARIABLE FIND);
procedure CLEAN UP HENRY DATA(HEAD in POINTER),
procedure SET UP HENRY ARRAY(HEAD in POINTER
HEAD LINE in LINE POINTER).

procedure SPRUCE UP HENRY DATA; function LOCAL NAME(NAME POINTER : in POINTER; SELECTOR : in SELECTOR TYPE; INDEX : in INTEGER;

return BOOLEAN:

function CALCULATE LINE COUNT(WORK LINE LINE POINTER) return INTEGER.

function FIND STRING SIZE(IN STRING LEXEME TYPE) RETURN INTEGER function TRANSITIVITY IN(IN NAME LEXEME TYPE.

BEGIN LOOP, STOP LOOP, POINTER)

RETURN FLOAT.

function TRANSITIVITY OUT(IN NAME: LEXEME TYPE

TOP POINTER)
RETURN FLOAT

procedure CALCULATE METRIC(HEAD in POINTER HEAD LINE in LINE POINTER).

```
end HENRY ANALYSIS.
package body HENRY ANALYSIS is
--starts the process of setting up the raw Henry records into
--decipherable data, it also counts the numbers of procedures
--functions and fills in empty parameter type fields in the
-- Henry records
procedure CLEAN UP HENRY DATA(HEAD: IN POINTER) is
TEMP. TOP. BOTTOM: POINTER:
begin
 put(HENRY FILE, "in CLEAN UP HENRY"); new line(HENRY FILE);
 CLEARSCREEN:
 put("Processing Henry data records ... please wait");
 TOP : HEAD:
 BOTTOM : TOP.NEXT1:
-- move past package declarations
 LOOP
  EXIT WHEN TOP TYPE DEFINE = END PACKAGE DECLARE:
        : BOTTOM:
  BOTTOM : - TOP.NEXT1:
 END LOOP:
--count the number of procedures functions
 LOOP
   EXIT WHEN BOTTOM TYPE DEFINE = END PACKAGE TYPE;
   if (BOTTOM.TYPE DEFINE = PROCEDURE TYPE) or
    (BOTTOM.TYPE DEFINE = FUNCTION TYPE) then
    PROC FUNC COUNT := PROC FUNC COUNT + 1:
   end if:
   TEMP :≈ BOTTOM;
   BOTTOM = TEMP.NEXTI;
 end loop:
 BOTTOM: TOP:
--ensure all parameter records have a type defined
   FOR Lin L.PROC FUNC COUNT LOOP
    LOOP
       EXIT WHEN (TOP TYPE DEFINE - PROCEDURE TYPE) OR
            (TOP.TYPE DEFINE FUNCTION TYPE):
      TOP : BOTTOM NEXT1:
```

```
BOTTOM TOP.
    END LOOP.
    TEMP - TOP NEXT).
    If TEMP TYPE DEFINE PARAM TYPE AND
     TOP TYPE DEFINE FUNCTION TYPE then
     LOOP
       EXIT WHEN TEMP TYPE DEFINE - END PARAM DECLARE:
       of TEMP PARAM. TYPE: NOT IN FORMAL, PARAM, CLASS THEN
        EXIT WHEN (TEMP PARAM TYPE IN TYPF) OR
             (TEMP PARAM TYPE OUT TYPE) OR
             (TEMP PARAM TYPE IN OUT TYPE):
        BOTTOM TEMP
        TEMP
                 BOTTOM NEXT1.
       END LOOP.
       BOTTOM:
                 TEMP
       ТЕМР
               TOP NEXT1
       TOP
               TEMP.
       LOOP
        EXIT WHEN (TOP PARAM TYPE IN TYPE) OR
             (TOP PARAM TYPE OUT TYPE) OR
             ETOP PARAM TYPE IN OUT TYPE).
        TEMP PARAM TYPE BOTTOM PARAM TYPE
               TOP NEXT1
        TEMP
              TEMP
        TOP
       END LOOP
       PINE
        TOP TEMP
        BOTTOM TEMP
       end if
       TEMP TOP NEXTI
     END LOOP
-functions usually invoke the default in type parameter
consert this type if it is not defined.
    elsif TOP TYPE DEFINE FUNCTION TYPE THEN
     If TEMP TYPE DEFINE PARAM TYPE THEN
      TOOP
       EXIT WHEN TEMP TYPE DEFINE. END PARAM DECLARE
       TEMP PARAMITYPE IN TYPE
       TEMP
               BOTTOM NEXTI
       BOTTOM
                TEMP
      END LOOP
     and if.
    and d
    \Gamma \cap \mathcal{V}
          BOTTOM NEXT
   BOLLOM TOP
  ENDIOOP FORLOOP
AND CLEAN UP HENRY DATA
```

```
--sets up the Henry data records to mark the beginning of each
```

--function or procedure, it also ties the procedure line length

--records to its proper procedure function

procedure SET_UP_HENRY_ARRAY(HEAD in POINTER: HEAD_LINE |: in LINE_POINTER) is

WORK LINE TEMP LINE : LINE POINTER: TEMP, TOP, BOTTOM : POINTER:

begin

put(HENRY FILE, "in SET_UP_HENRY"); new_line(HENRY FILE); WORK_LINE: HEAD_LINE.NEXT_REC; TEMP_LINE: WORK_LINE: BOTTOM:= HEAD; TOP := BOTTOM.

-- GO PAST DECLARATIONS

LOOP

EXIT WHEN TOP.TYPE DEFINE END PACKAGE DECLARE: TOP BOTTOM.NEXT1;
BOTTOM TOP:
END LOOP:

--set up the Henry array records so that their pointers are at the --top of each procedure or function

FOR Lin 1 PROC FUNC COUNT LOOP LOOP

EXIT WHEN (TOP TYPE DEFINE PROCEDURE TYPE) OR (TOP.TYPE DEFINE FUNCTION TYPE).

TOP BOTTOM NEXT1

BOTTOM TOP.

END LOOP

HENRY ARRAY(I) NAME OF DATA(I MAX LINE SIZE) TOP NOMEN(I MAX LINE SIZE).

HENRY ARRAY(I) BEGIN POINTER TOP.

EXIT WHEN (BOTTOM TYPE DEFINE END FUNCTION TYPE) OR (BOTTOM TYPE DEFINE END PROCEDURE CALL)

TEMP BOTTOM NEXT1

BOTTOM TEMP

END LOOP

set the area outstresseds to their related procedure function

TOP BOTTOM NEXT BOTTOM TOP HENRY ARRAYDOTINE LENGTH POINTER WORK TINE WORK TINE TEMP TINE NEXT REC TEMP TINE WORK TINE

```
END LOOP: --FOR LOOP
end SET UP HENRY ARRAY:
-- this procedure calculates the length of each procedure function
-- the results are fed into line length records
function CALCULATE LINE COUNT(WORK LINE LINE POINTER)
                return INTEGER is
DIFFERENCE INTEGER: 0:
      INTEGER.
begin
 put(HENRY FILE, "in CALCULATE LINE COUNT"); new line(HENRY FILE).
 DIFFERENCE WORK LINE STOP COUNT - WORK LINE START COUNT:
 RETURN (DIFFERENCE);
end CALCULATE LINE COUNT:
......
-this function searches for local, within a procedure, and global-local,
-within a package, for variable name matches
--it is selectable for which name the search is conducted
function LOCAL NAME(NAME POINTER in POINTER:
          SELECTOR
                       in SELECTOR TYPE.
          INDEX
                       in INTEGER )
           return BOOLEAN is
NAME SOUGHT POINTER NAME LEXEME TYPE:
NAME SIZE POINTER SIZE INTEGER MAX LINE SIZE:
RESULT
                  BOOLEAN FALSE.
TEMP. TEMP1
                     POINTER.
               INTEGER 1:
begin
pot(HENRY FILE, "in LOCAL NAME"), new line(HENRY FILE);
NAME SOUGHT(I NAME SIZE). NAME POINTER NOMEN(I NAME SIZE).
CONVERT UPPER CASE(NAME SOUGHT, NAME SIZE);
if (SELECTOR PROCEDURE FIND) OR (SELECTOR FUNCTION FIND))
 AND PROCEUNG COUNT
 TOOP
   POINTER NAME(1 POINTER SIZE)
  HENRY ARRAY(I) NAME OF DATA(I POINTER SIZE).
  CONVERT UPPER CASE(POINTER NAME, POINTER SIZE)
           (NAME SOUGHT(1 NAME SIZE)
  RESULT
        POINTEL NAME(I POINTER SIZE)).
```

EXILWHEN (I. PROC FUNC COUNT) OR (RESULT):

- F - T - T END 1 O O P

```
1: - 1:
--if it is a variable name search first within the package
--declarations, next within the procedure declarations
elsif SELECTOR - VARIABLE FIND then
 TEMP := HEAD.NEXT1;
 LOOP
  EXIT WHEN (TEMP.TYPE DEFINE = END PACKAGE DECLARE) OR
        (RESULT):
  if TEMP. TYPE DEFINE = IDENT TYPE then
    POINTER NAME(1.POINTER SIZE):= TEMP.NOMEN(1.POINTER SIZE);
    CONVERT UPPER CASE(POINTER NAME, POINTER SIZE);
    RESULT := (NAME SOUGHT(1...NAME SIZE) =
         POINTER NAME(1..POINTER SIZE)):
  end if:
  TEMP1 := TEMP.NEXT1:
  TEMP := TEMP1:
 END LOOP;
--did not find the variable within the package declarations
--search the specified procedures declarations
 if NOT RESULT then
   TEMP: = HENRY ARRAY(INDEX).BEGIN POINTER:
  LOOP -- DID NOT FIND NAME IN PACKAGE DECLARATIONS
    EXIT WHEN (TEMP.TYPE DEFINE : END DECLARATIONS) OR
           (RESULT):
    if TEMP. TYPE DEFINE : IDENT TYPE then
      POINTER NAME(1. POINTER SIZE)
         TEMP.NOMEN(1..POINTER SIZE):
      CONVERT UPPER CASE(POINTER NAME, POINTER SIZE):
      RESULT := (NAME SOUGHT(1..NAME SIZE)
         POINTER NAME(1. POINTER SIZE)).
    end if:
    TEMP1 := TEMP.NEXT1:
    TEMP := TEMP1;
  END LOOP:
 end if:
end if:
RETURN (RESULT):
end LOCAL NAME;
```

- --finishes polishing the Henry records, the data can now be analyzed
- -- for local global data and starts the actual number crunching

procedure SPRUCE UP HENRY DATA is

```
TEMP, TEMP1, TEMP2: POINTER,
begin
 put(HENRY FILE, "in SPRUCE UP HENRY"); new line(HENRY FILE);
 FOR Lin 1.. PROC FUNC COUNT LOOP
   TEMP1 := HENRY ARRAY(I).BEGIN POINTER;
--loop past parameters
   LOOP
    EXIT WHEN TEMP1. TYPE DEFINE = END DECLARATIONS;
    TEMP2 := TEMP1.NEXT1;
    TEMP1 := TEMP2;
   END LOOP:
   TEMP := TEMP1.NEXT1:
--first analyze identifier types (variables) for local or global
--significance. Update the record if it is not local
   LOOP -- LOOK FOR IDENT TYPES
    EXIT WHEN (TEMP. TYPE DEFINE = END FUNCTION TYPE) OR
          (TEMP.TYPE DEFINE = END PROCEDURE CALL);
    if TEMP.TYPE DEFINE = IDENT TYPE then
      if TEMP.IDENTITY = BLANK then
        if LOCAL NAME(TEMP, VARIABLE FIND, I) then
          TEMP.IDENTITY := LOCAL DECLARE;
        else TEMP.IDENTITY := GLOBAL DECLARE:
        end if:
      end if:
     end if:
     TEMP1: TEMP.NEXT1:
     TEMP := TEMP1;
   END LOOP:
--now go through the Henry records looking for unresolved
--procedure or function calls update the Henry records
-- to reflect procedure types or function types or data structures
            HENRY ARRAY(I).BEGIN POINTER:
   TEMPL
   TEMP
            TEMP1 NEXT1.
-get past declarations
  LOOP
    EXIT WHEN TEMP TYPE DEFINE - END DECLARATIONS
```

TEMP NEXTL

TEMPT

TEMP1

TEMP END LOOP

```
LOOP
    EXIT WHEN (TEMP TYPE DEFINE - END FUNCTION TYPE) OR
         (TEMP.TYPE DEFINE = END PROCEDURE CALL);
    if TEMP.TYPE DEFINE - PROCALL OR DS then
     TEMP1 := TEMP;
     LOOP
                -- MOVE PAST THE PARAMETERS
       EXIT WHEN TEMP1.TYPE DEFINE = END ACTUAL PARAM;
       TEMP2 := TEMP1;
       TEMP1 := TEMP2.NEXT1:
     END LOOP:
     if (LOCAL NAME(TEMP, PROCEDURE FIND. I)) then
       TEMP.TYPE DEFINE := PROCEDURE TYPE;
       TEMP2 := TEMP1.NEXT1:
       if TEMP2.TYPE DEFINE = ASSIGN TYPE then
         TEMP. TYPE DEFINE := DATA STRUCTURE:
         --IF NOT IT IS A PROCEDURE CALL ONLY
        TEMP1 := TEMP2.NEXT1:
         LOOP
         EXIT WHEN TEMP1.TYPE DEFINE = END ASSIGN TYPE:
          if (TEMP1.TYPE DEFINE - FUNCALL OR DS) then
            if NOT LOCAL NAME(TEMP1, FUNCTION FIND, I)
            then
             TEMP1 TYPE DEFINE : DATA STRUCTURE:
            else TEMP1.TYPE DEFINE: FUNCTION TYPE:
            end if:
          end if;
          TEMP2: TEMP1:
          TEMP1: TEMP2.NEXT1:
         END LOOP:
       else TEMP TYPE DEFINE PROCEDURE TYPE:
       end if:
     end if.
--only function calls that cannot be resolved into a local name are
--specified as data structures
    elsif TEMP TYPE DEFINE FUNCALL OR DS then
     TEMPI
            TEMP.
     LOOP
           -- LOOKING FOR FUNCTIONS
     EXIT WHEN TEMP. TYPE DEFINE END ASSIGN TYPE:
       if TEMP TYPE DEFINE FUNCALL OR DS then
         if (LOCAL NAME(TEMP, FUNCTION FIND, I))
          TEMP TYPE DEFINE
                              FUNCTION TYPE:
        #ISO TEMP TYPE DEFINE DATA STRUCTURE:
        ent if
       en fif
       TEMP1 TEMP NEXT1
```

```
TEMP := TEMP1;
      END LOOP:
    end if:
    TEMP1 := TEMP:
    TEMP := TEMP1.NEXT1;
   END LOOP: --PROCALL OR DS LOOP
 END LOOP: -- FOR LOOP
end SPRUCE UP HENRY DATA:
--this function only works for Ada language strings that identify
--a variable
function FIND STRING SIZE(IN STRING: LEXEME TYPE) RETURN INTEGER
SIZE : INTEGER := 0;
BEGIN
PUT(HENRY FILE, "IN FIND STRING SIZE"); NEW LINE(HENRY FILE);
 FOR LIN LIMAX LINE SIZE LOOP
   IF IN STRING(I) = \overline{N}ULL CHAR THEN
    SIZE := SIZE + 1:
   END IF:
 END LOOP:
 RETURN SIZE:
END FIND STRING SIZE:
--transitivity is detected by searching the right hand side of
--assignment statements for a name match of the actual
--parameters from a function or procedure call
function TRANSITIVITY IN(IN NAME : LEXEME TYPE;
              BEGIN LOOP. STOP LOOP: POINTER)
              RETURN FLOAT is
ASSIGN MARK.
PROCEDURE MARK: BOOLEAN: FALSE;
TRANS COUNT : FLOAT := 0.0:
TEMP. TEMP1 : POINTER := BEGIN LOOP:
T1. T2
          : POINTER:
           : INTEGER . * MAX LINE SIZE:
MAX
BEGIN
--stop loop is determined by where in the parameter list you are
 LOOP
   EXIT WHEN TEMP STOP LOOP:
   if TEMP. TYPE DEFINE ASSIGN TYPE THEN
```

```
ASSIGN MARK := TRUE;
  elsif TEMP.TYPE DEFINE = END ASSIGN TYPE THEN
     ASSIGN MARK := FALSE:
  end if:
--mark whether you've passed an assignment
   if (TEMP.NOMEN(1..MAX) = IN.NAME(1..MAX)) AND
     (NOT ASSIGN MARK) THEN
    TRANS COUNT := TRANS COUNT + 1.0;
-if you have detected a name match count the number of assignment
--variables as transitive feed into the actual parameter
--note functions have already been calculated the same for
--data structures so skip these counts
    T1 := TEMP; T2 := T1.NEXT1;
    if (T1.TYPE DEFINE = IDENT TYPE) AND
      (T2.TYPE DEFINE - ASSIGN TYPE) then
      LOOP
       EXIT WHEN T2. TYPE DEFINE - END ASSIGN TYPE:
       If T2 TYPE DEFINE = IDENT TYPE THEN
           TRANS COUNT := TRANS COUNT + 1.0:
       end if:
       T1 := T2:
       T2 : T1 NEXT1:
      END LOOP:
     end if:
   end if:
   TEMP := TEMP1.NEXT1;
   TEMP1 := TEMP;
 END LOOP:
 RETURN(TRANS COUNT):
END TRANSITIVITY IN:
--if detect a name match on the right hand side of an assignment
--statement have a transitive relation on this variable but
-- there is no need to count the rest of the assignment
--variables because the most it can account for is 1
function TRANSITIVITY OUT(IN NAME: LEXEME TYPE;
               TOP : POINTER)
               RETURN FLOAT is
ASSIGN MARK : BOOLEAN : FALSE:
TRANS COUNT : FLOAT := 0.0:
TEMP, TEMP1: POINTER: - TOP:
          : INTEGER: MAX LINE SIZE:
MAX
```

```
BEGIN
LOOP
  EXIT WHEN (TEMP TYPE DEFINE END PROCEDURE CALL) OR
       (TEMP TYPE DEFINE - END FUNCTION TYPE).
  IF TEMP. TYPE DEFINE ASSIGN TYPE THEN
    ASSIGN MARK :- TRUE:
  ELSIF TEMP. TYPE DEFINE END ASSIGN TYPE THEN
    ASSIGN MARK := FALSE,
  END IF:
  IF (TEMP.NOMEN(1, MAX) IN NAME(1, MAX)) AND (ASSIGN MARK)
    TRANS COUNT: TRANS COUNT + 10:
  END IF:
  TEMP : TEMPLNEXTI:
  TEMP1 := TEMP;
 END LOOP:
 RETURN(TRANS COUNT);
END TRANSITIVITY OUT:
-- finishes polishing the data and with the transitivity functions calculates
-- the fan in fan out of data besides the global data structures
procedure CALCULATE METRIC (HEAD)
                                     in POINTER.
               HEAD LINE in LINE POINTER) is
TEMP LINE
                  LINE POINTER.
TEMP, TOP, TEMP1, TEMP2 POINTER
PROC PTR.
PARAM PTR
                    POINTER
FAN IN FAN OUT
                     FLOAT.
LENGTH
                 : INTEGER = 0.
                INTEGER
MAX
                           MAX LINE SIZE.
CODE EXPONENT
                     INTEGER
COMPLEXITY.
GLOBAL FLOW
GLOBAL READ.
GLOBAL WRITE.
                        FLOAT.
GLOBAL READ WRITE
```

- rigiobal flow represents the whole picture of global data flow
- -the equation is below and encompasses both read and write to
- -- global data structures
- --note, global data structures could be external function calls
- athere is no means to determine the difference

NEW NAME STRING(I MAX LINE SIZE)
NAME OF LEXEME TYPE
ASSIGN MARK

```
SIZE
                                      INTEGER MAX LINE SIZE
NEW SIZE
                                               INTEGER (c)
TEMP_NAME
                                                    STRING(L SIZE)
heein
  PUT(HENRY FILE "IN CALCUALTE METRIC". NEW TINE (HENRY FILE)
whest Herry, all boolean is so that the data can be reshown
 IF FIRST HENRY CALL then
  CLEAN UP HENRY DATA(HEAD)
  SET UP HENRY ARRAY HEAD HEAD LINE.
   SPRUCE UP HENRY DATA
   FOR The L. PROC. FUNC. COUNT LOOP.
       GLOBAL READ
       GLOBAL WRITE
       GLOBAL READ WRITE
       FAN IN
       FAN OUT
       COMPLEXITY
       GLOBAL FLOW
       LENGTH
       TEMP HENRY ARRAYA BE AN POINTER
       CLEAR HENRY LEXEME, LEMP NAME
       TEMP NAMED MAN LINE SIZE - HENRY ARRAY E NAME OF DAY A
       -17F
                         FIND STRING SIZE TEMP NAME
       CLEAR HENRY TEXEME-NEW NAME.
       CONVERT UPPER CASE TEMP NAME SIZE
  anothalize the variables time a trivial place of the floor of any order to the
  stational en-
       SCIEMPTYPE DEFINE PROCEDURE TYPE Own.
           OUT PUT DATAIL TYPE OF A PROCEDURE TYPE
            NEW SIZE SIZE IN The SECOND INCOME THE SECOND IN THE SECON
            NEW NAMED NEW SIZE TEMP NAME I SIZE
           OF PULL DATA I NAME OF UNEW SIZE IN NEW NAME ON WORK
           PULHENRY OUT "
            NEW LINEHIENRY OF LIZE
           PUT HENRY OUT NEW NAME
            NEW LINESHENRY OF IT
```

SET LEMP TYPE DEFINE OF FUNCTION TYPE 19 (1)

OUT PUT DATAGOTYPE (FOR FUNCTION TYPE

NEW SIZE SIZE (1)

NEW NAME (FOR SIZE OF FOR WORLD SIZE

OUT PUT DATAGONAME (FOR SIZE OF FOR WORLD SIZE

PUT HENRY OUT SEW NAME

PUT HENRY OUT NEW NAME

```
NEW TINE HENRY OF TH
ent of
normal sea precisioname for the data file.
TEMPT TEMP NEXT
HAMPI HAPI
to a toric this carraties increase global flow metric
1-1-1-1
  EXIL WHEN TEMPLITYPE DEFINE. END FUNCTION TYPE) OR
        TEMPLIAPE DEFINE END PROCEDURE CALL)
  STEMPS TYPE DEFINE ASSIGN TYPE then
   ASSIGN MARKER TRUE
    A TEMPLIAPE DEFINE FND ASSIGN TYPE then
   ASSIGN MARKER FAISE
   CHORAL MARKER - FALSE
   . , .
  SOLEMBOODEN HELY OF GLOBAL OF CLARE (AND (ASSIGN MARKER)
   STABAL PEAD SCIONAL READ STOR
    THE HAL MARKER ber
     CHORAL READ WRITE GLOBAL READ WRITE - 10.
    SOFT THE STATE OF GRAD DECLARED AND
     Cook Assich MARKER then
     TO BALLWRITE GODBAL WRITE LOS
     CORNE MARKER TRUE
   . . .
  TEMES TEMES NEXT (
  PEMP2 TEMP1
 rate of the
  and the second of the same to be to be formally acameters
 CLASS SEARCH NEXT C
  CANDAL CARE DEFINE PARAM TAPE then
    A X 11 WHEN TEMPT LYPE DEFINE FIND PARAM DECLARE
    A TEXABLE VARIANT LAPPE OF A TAPE THEY
     TANK BURELINGS
       CONTRACTOR AND TAPE COURT TAPE THEN
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      SCHOOL CARAMITAPE OF SOUT TAPETHEN
      AND THE FANTS IN
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     FAM: FAM: FAM: FAM:
            - AMP SEXTE
```

--look for procedure and function type actual parameters

```
TEMP
       TEMP1.
TEMPI
       TEMP NEXTI
TEMP
       TEMPT
LOOP
 EXIT WHEN (TEMP TYPE DEFINE FND FUNCTION TYPE OR
      (TEMP TYPE DEFINE END PROCEDURE CALL
 TEMPL
         TEMP
 TEMP
        TEMPI NEXTI
 If TEMP TYPE DEFINE ASSIGN TYPE then
   ASSIGN MARKER
                 TRUE
 elsif TEMP TYPE DEFINE END ASSIGN TYPE (nep
   ASSIGN MARKER - FALSE
  GLOBAL MARKER FALSE
 If TEMP TYPE DEFINE PROCEDURE TYPE Once
  TEMPI TEMP NEXTI
  LOOP
    EXIT WHEN TEMPLIAPE DEFINE LEND ACTUAL CARAM
    FAN OUT FAN OUT TO
    TEMP2
            TEMPL
    TEMPI
            TEMP2 NEXT1
  END LOOP
 elsif TEMP TYPE DEFINE FUNCTION TYPE THEN
```

-- Sount the function parameters

```
TEMPT
         TEMP NEXT1
  LOOP
    EXIT WHEN TEMPLIAPE DEFINE LEND ACTUAL EXHAM
    FAN OUT - FAN OUT - For
    TFMP2
           TEMPT
    TEMPT
           TEMP2 NEXT1
  ENDIOOP
  FANIN FANIN TO RETURN FROM FUNCTION
 MSOF (TEMP TYPE DEFINE DATA STRUCTURE AC
    (NOT ASSIGN MARK) THEN
  GLOBAL MARK TRUE
  GLOBAL WRITE GLOBAL WRITE
 HISTOTEMP TYPE DEFINE TRAIN STRUCTURE TO A TOTAL
    THEZ
   COLOBAL MARK THEN
    GLOBAL READ WRITE GLORAG RUA CAR OF
  CALNOT GLOBAL MARK THEN
   CORAL READ CONTRACTOR
  and the
1 M 1 - 10 P
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FOR A SHE NOR OF PIRTINGS OF END FENCTION TYPE OR

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```
GLOBAL FLOW := GLOBAL WRITE *
          (GLOBAL READ + GLOBAL READ WRITE) +
          GLOBAL READ WRITE *
          (GLOBAL READ + GLOBAL READ WRITE - 10).
   put(HENRY OUT, "NUMBER OF LINES = ");
   put(HENRY OUT, LENGTH):
   OUT PUT DATA(I).CODE LENGTH LENGTH:
   NEW LINE(HENRY OUT);
   put(HENRY OUT, "FAN IN =
   put(HENRY OUT, FAN IN);
   OUT PUT DATA(I). TYPE FAN IN := FAN IN:
   NEW LINE(HENRY OUT):
   put(HENRY OUT, "FAN OUT =
   put(HENRY OUT, FAN OUT);
  OUT PUT \overline{D}ATA(I). TYPE FAN OUT := FAN OUT;
  NEW LINE(HENRY OUT):
   put(HENRY OUT, "COMPLEXITY =
                                  "):
  put(HENRY OUT, COMPLEXITY):
  OUT PUT DATA(I). TYPE COMPLEXITY := COMPLEXITY.
  NEW LINE(HENRY OUT);
  put(HENRY OUT, "GLOBAL READ =
                                  ");
  put(HENRY OUT, GLOBAL READ):
  OUT PUT DATA(I). TYPE READ
                                := GLOBAL READ.
  NEW LINE(HENRY OUT):
  put(HENRY OUT, "GLOBAL WRITE =
                                   "):
  put(HENRY OUT, GLOBAL WRITE);
  OUT PUT DATA(I). TYPE WRITE
                                 : GLOBAL WRITE
  NEW LINE(HENRY OUT);
  put(HENRY OUT, "GLOBAL READ WRITE = ");
  put(HENRY OUT, GLOBAL READ WRITE):
  OUT PUT DATA(I). TYPE READ WRITE
                                     GLOBAL READ WRITE
  NEW LINE(HENRY OUT).
  put(HENRY OUT, "GLOBAL FLOW
  put(HENRY OUT, GLOBAL FLOW);
  OUT PUT DATA(I). TYPE FLOW
                                   GLOBAL FLOW
  NEW LINE(HENRY OUT, 2):
END LOOP:
PUT(HENRY OUT, "-----")
end if. -- FIRST HENRY CALL.
FIRST HENRY CALL: FALSE
END CALCULATE METRIC.
END HENRY ANALYSIS
```

TITLE AN ADA SOFTWARE METRIC

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RUNNING COUNT Insome INTEGER

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The state SET UP SCREEN(IN STRING IN ROW STRING TYPE

STRING SIZE IN INTEGER)

From the CENTER STRING(NAME in ROW STRING TYPE)

FN ROW WIDTH in lateger)

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 $\label{eq:constraints} N \mathbf{I} = (\mathbf{K}_1, \mathbf{M}_2, \mathbf{K}_3, \mathbf{M}_2, \mathbf{K}_3, \mathbf{M}_3, \mathbf{K}_3, \mathbf{K}_4, \mathbf{K$

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IN ROW, WIDTH in INTEGER) is

SCREEN WIDTH INTEGER 76: CENTER POS INTEGER 0. TEMP NAME ROW STRING TYPE: begin FOR LIN 1 30 LOOP TEMP NAME(I): NULL CHAR: END LOOP: TEMP NAME(1. WIDTH) NAME(1. WIDTH): CENTER POS - SCREEN WIDTH 2 - WIDTH 2: SET CURSOR POS(CENTER POS. IN ROW). PUT(TEMP_NAME): NEW LINE: end CENTER STRING. -Puts the header of each data screen up with an underline to set it --off from the data procedure SET_UP_SCREEN(IN_STRING_in_ROW_STRING_TYPE. STRING SIZE in INTEGER) is tagin. CLEARSCREEN. SET_REVERSE(ON) CENTER STRING(IN STRING 1 STRING SIZE) SET_REVERSE(OFF) NEW LINE(2). END SET UP SCREEN clisis the entire record stream of the Henry metric data procedure LIST HENRY DATA is TEMP NAME TEXEME TYPE SHOP! NAME STRING I SHORT NAME SUIT 4×1+c+R PUNCON COUNTY INTRODUCE 1.01 INTEGER 2 HEALTE STRING R. W. STRING CO. CO. 10 11 48 FACE

```
begin
 HEADER STRING(1 HEADER SIZE) - "TIST OF HENRY RECORDS"
 PUTCHENRY FILE "IN LIST HENRY DATA": NEW LINECHENRY FILE
 LOOP EXIT WHEN DONE
 SET UP SCREEN/HEADER STRING HEADER SIZE:
 LOOP
  pute"DECLARATION "
  State TEMP POINTER IDENTITY is
    when LOCAL DECLARE put;"I stall declare."
    when GLOBAL DECLARE put 1G1 has to lare "-
    when others par "I interpared":
   end tase
   new line
   par NAME
   FITEMP POINTER NUMENT - NULL SHAR For
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   TEMP NAME OMAN TONE SIZE
       TEMP POINTER NOMENT MAN LINE SIZE
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THE SRY FILE IN A LIT METRIC DATA. NEW TINE HENRY FILE OF VIEW STRING FIHE VIEW SIZE OF THE LIBIT HENRY METRIC VALLES OF FEN HENRY OF LITTURE SIZE.

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```
RESET HENRY OUT, IN FILE).
PLSE OPEN(HENRY OUT, IN FILE, HENRY OUT NAME):
er. I if
 SET UP SCREEN(HEADER STRING, HEADER SIZE);
 IN STRING(1 8) "PACKAGE ":
 (N STRING(9-49) INPUT FILE NAME(1..41):
 par IN STRING)
 NEW LINE(2)
 LOOP
  EXIT WHEN (END OF FILE(HENRY OUT) OR DONE);
  FOR JIN 1 49 LOOP
  IN STRING(J) NULL CHAR:
  END LOOP:
  GET LINE(HENRY OUT, IN STRING, NUMBER OF);
  PUT LINEAR STRING).
  PAUSE PRINT(STOP RUNNING COUNT, DONE);
  IF RUNNING COUNT 0) AND (NOT DONE) THEN
   RUNNING COUNT -- 1.
   SET UP SCREEN(HEADER STRING, HEADER SIZE);
  end if
 END LOOP
  IF NOT DONE; THEN
   STOP 1 RUNNING COUNT := 1;
   PAUSE PRINT(STOP, RUNNING COUNT, DONE):
  en Lif
 CLOSE(HENRY OUT):
er i 118T METRIC DATA
......
```

of ists the relative comparison metric data. This listing in pares each pricedure function analyzed with for example the of an innumbers. It also gives a verbal report for each function out to enture.

procedure WRITE RELATIVE DATA is

INDICATOR:

1.20

```
HEADER STRING(1..HEADER SIZE) := "THE RELATIVE PERFORMANCE DATA":
 SET UP SCREEN(HEADER STRING, HEADER SIZE);
 if PROC_FUNC_COUNT < 16 THEN STOP := PROC_FUNC_COUNT:
 else STOP := 16:
 PUT(HENRY FILE, "IN WRITE RELATIVE DATA"); NEW LINE(HENRY FILE);
-- name the outer loop so that can exit gracefully when the user
-- wants to quit
OUTER LOOP:
FOR J IN 1..7 LOOP
CASE J is
 when 1 \Rightarrow ROW STRING(1..6) := "FAN IN";
      SIZE := 6;
 when 2 = > ROW STRING(1..7) := "FAN OUT":
      SIZE := 7;
 when 3 \Rightarrow ROW STRING(1..10) := "COMPLEXITY";
      SIZE := 10;
 when 4 => ROW STRING(1..11) := "GLOBAL READ";
      SIZE := 11:
 when 5 => ROW STRING(1..12) := "GLOBAL WRITE";
      SIZE = 12:
 when 6 => ROW STRING(1..17) := "GLOBAL READ WRITE";
      SIZE : 17:
 when 7 => ROW STRING(1..11) := "GLOBAL FLOW";
      SIZE := 11;
 when others - null:
 end case.
 CENTER STRING(ROW STRING, 4, SIZE);
 FOR LIN 1. PROC. FUNC. COUNT LOOP.
 SET CURSOR POS(1, 1 - 5):
 REL ARRAY(I).NAME OF: OUT PUT DATA(I).NAME OF:
 PUT(REL_ARRAY(I).NAME_OF); SET_CURSOR_POS(42, I + 5); PUT(":");
--set up the names before write the data
 CASE J is
  when 1 - put(REL ARRAY(I) TYPE FAN IN).
  when 2 -- put(REL ARRAY(I) TYPE FAN OUT);
  when 3 -- put(REL_ARRAY(I) TYPE_COMPLEXITY);
  when 4
         - put(REL_ARRAY(I) TYPE_READ);
  when 5
         -- put(REL_ARRAY(I) TYPE_WRITE).
  when 6 - put(REL_ARRAY(I) TYPE_READ_WRITE);
  when 7 -- put(REL_ARRAY(I) TYPE_FLOW)
  when others
             - nu∐
 end case.
 NEW LINE.
 PAUSE PRINT(STOP RUNNING DONE)
```

--boolean done is set true by user answering the querry to quit

```
EXIT OUTER LOOP WHEN DONE:
 if (RUNNING = 0) AND (STOP = 16) THEN
    STOP : PROC FUNC COUNT - 17;
 elsif RUNNING = 0 THEN
    SET UP SCREEN(HEADER STRING, HEADER SIZE);
    RUNNING := 1;
   end if;
 end loop:
end loop OUTER LOOP:
--set up to loop again once have cycled through to first stop
--count. This means have filled the screen once
STOP := 1; RUNNING := 1;
PAUSE PRINT(STOP, RUNNING, DONE):
CLEARSCREEN;
PUT("The following are the maximums for each calculation: ");
put("-----");
new line:
put("Fan In
               : "); put(REL ARRAY(MAX FAN IN).NAME OF); new line:
put("Fan Out
               : "); put(REL ARRAY(MAX FAN OUT).NAME OF); new line;
put("Complexity : "): put(REL ARRAY(MAX COMPLEXITY).NAME OF): NEW LINE:
put("Global Read : "); put(REL ARRAY(MAX READ).NAME OF); NEW LINE;
PUT("Global Write : "): put(REL ARRAY(MAX WRITE).NAME OF); NEW LINE:
PUT("Global Read Write: "); put(REL_ARRAY(MAX_READ_WRITE), NAME_OF); NEW_LINE;
PUT("Global Flow : "); put(REL_ARRAY(MAX_FLOW).NAME_OF); NEW_LINE;
new line:
put("-----");
new line:
STOP := 1; RUNNING := 1;
PAUSE PRINT(STOP, RUNNING, DONE);
SET UP SCREEN(HEADER STRING, HEADER SIZE):
--calculate the indicator numbers so that can determine the relative
--performance of each procedure function within each category
FOR LIN L.PROC FUNC COUNT LOOP
 if REL ARRAY(1). TYPE FLOW = 0.0 THEN
 INDICATOR1 := REL ARRAY(I).TYPE COMPLEXITY
         REL ARRAY(I). TYPE FLOW:
 else INDICATOR1 := REL ARRAY(I).TYPE COMPLEXITY:
 end if:
 if REL_ARRAY(I).TYPE_FAN_OUT = 0.0 THEN
 INDICATOR2: REL ARRAY(I). TYPE FAN IN
         REL ARRAY(I). TYPE FAN OUT:
 else INDICATOR2 : REL ARRAY(I) TYPE FAN IN:
 if REL_ARRAY(I). TYPE_WRITE = 0.0 THEN
 INDICATOR3 . REL ARRAY(I).TYPE READ
         REL ARRAY(I) TYPE WRITE:
```

```
else INDICATOR3 REL ARRAY(I) TYPE READ
PUT(REL ARRAY(L NAME OF) put " ")
new line.
--put out the results of the indicator analysis
 JE INDICATOR I UPPER LIMIT THEN
  PUT(" - Has significant complexity compared to gli bal data flow ")
   new line
   EINDICATOR2 - UPPER LIMIT THEN
    put(" - This implies poor internal code structure. ")
    put(" - Consider remodularization ").
    new line
  elsif INDICATOR2 LOWER LIMIT THEN
    PUT(" - This implies an extremely complex interface ")
    new line
  end if
 ELSIF INDICATOR1 - LOWER LIMIT THEN
  PUT(" - Has significant global data flow compared to complexity ").
  new line.
  if INDICATOR3 UPPER LIMIT THEN
   put(" - This implies an overworked data structure ")
   put(" - or a considerable number of function calls ").
   put(" - Consider redistributing the data structure into this module")
   new line
  elsif INDICATOR3 - LOWER LIMIT THEN
   PUT(" - This implies a program stress point").
   put(" - or a critical data flow point").
   put(" - Consider reorganizing the data structure.");
   new line
  end if.
 ELSE
   TEMP HOLDER(1/10) = REL ARRAY(I).NAME OF(1/10);
   put(" - Is a fairly well balanced "); put(TEMP_HOLDER).
   put(" - This implies good modularization. "):
   new line
 END IF,
STOP 1: RUNNING : 1:
PAUSE PRINT(STOP, RUNNING, DONE):
EXIT WHEN DONE.
end toop:
if NOT DONE THEN
STOP
       1. RUNNING 1 1:
PAUSE PRINT(STOP, RUNNING, DONE):
end if
```

and WRITE RELATIVE DATA

produces a Carolhart of lactometric labellation to be come in open con-

CONTROL OF APPEARING DATA A

LOOP CAT INTEGER

ROW STRING ROW STRING TYPE

HEADER STRING ROW STRING TYPE

SIZE INTEGER

STOP RUNNING INTEGER I

DONE BOOLEAN FAISE

SCALE INTEGER

ALM LOOP CAT

REM CAT INTEGER

HEADER SIZE INTEGER

1 4816

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The possibility the number of screens need to display or emainder of unit is the partial screen that is left, yer

HEADER STRING(1.30) THE GRAPHICAL PERFORMANCE DATA of NUM LOOP CNT 1. THEN STOP 5. STOP REM CNT end of PUT(HENRY FILE "IN WRITE RELATIVE DATA") NEW LINE(HENRY FILE) SET UP SCREEN(HEADER STRING HEADER SIZE).

asset up to exit gracefully when the user wants to quit

GRAPH LOOP FOR JIN 1-7 LOOP CASE J is when I -- ROW STRING(1.6) "FAN IN" SIZE 6. when 2 -- ROW STRING(1.7) "FAN OUT". SIZE 7 when 3 -- ROW STRING(1 10) "COMPLEXITY" SIZE 10. when 4 - ROW STRING(1-11) -"GLOBAL READ". SIZE 11 when 5 -- ROW STRING(1.12) -- "GLOBAL WRITE"; SIZE 12. when 6 -- ROW STRING(1...17) : "GLOBAL READ WRITE". SIZE: 17. when 7 → ROW STRING(1..11) : "GLOBAL FLOW"; SIZE : 11;

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RUNNING OF AND (STOP) IN THEN
a PROCEEUNG COUNT - 5 THEN
  STOP
en faf
  SET UP SCREEN(HEADER STRING HEADER SIZE)
  RUNNING I
elsif (RUNNING = 0) AND (STOP = 5) THEN
 if PROCEEUNG COUNT 5 THEN
  SET UP SCREEN(HEADER STRING HEADER SIZE)
  RUNNING
 elsif PROCEFUNC COUNT 5 THEN
  NUM LOOP CNT - NUM LOOP CNT - 1.
  if NUM LOOP CNT > 1 THEN STOP = 5, RUNNING = 1
  elsif REM_CNT = 0 THEN
    STOP REM CNT RUNNING 1.
  else SET UP SCREEN(HEADER STRING, HEADER SIZE
     RUNNING L
  end if.
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end if:

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FOUL PUT DATA(AX READ) TYPE READ OF THEN REL ARRAY(BIYPE READ NORMALIZER *
FOUL PUT DATA(BIYPE READ
OUT PUT DATA(MAX READ) TYPE READ)
FISH REL ARRAY(BIYPE READ 0.0)
FISH REL ARRAY(BIYPE READ 0.0)

of OUT PUT DATA(MAX WRITE) TYPE WRITE 00 THEN REL ARRAY(I) TYPE WRITE NORMALIZER*
(OUT PUT DATA(I) TYPE WRITE
OUT PUT DATA(MAX WRITE) TYPE WRITE).

else REL ARRAY(I) TYPE WRITE: 00
end of

of OUT PUT DATA(MAX READ WRITE) TYPE READ WRITE 0.0 THEN REL ARRAY(I) TYPE READ WRITE. NORMALIZER * (OUT PUT DATA(I).TYPE READ WRITE

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LUXBURGE BYPASS FUNCTION IS

25 s tunction compares the lexeme of the arrent token with the

it ken currently being sought by the parser. If the current token

Type is identifier, then a test is conducted to ensure it is not.

a reserved word

the strong BYPASS(TOKEN ARRAY ENTRY CODE integer) return to alean is

CONSUME boolean FALSE

TEXTIME string(LTINESIZE)

~I/F natural

HENRY LEXEME. string(L MAX EINE SIZE)

tiegan

GET CURRENT TOKEN RECORD(CURRENT TOKEN RECORD LEXEME LENGTH)
LEXEME CURRENT TOKEN RECORD LEXEME.

SIZE CURRENT TOKEN RECORD LEXEME SIZE - 1

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                                          . .
                                                             1.3.1.3.11
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           SONATE IN THER SASE HENRY OF ALMOST A
            SAME SEPERAL ASE LEXEME SIZE
            MENER WRITE ENABLE For
   whether the body was not for the common on the Physical Police
           WRITE HEARY DATA BEANN HEARY JEANNE DEAT 1966 A SOCIETY
           CEEVIE NODENEXT HEN LAST RECORD
          HENRY WRITE EXABLE FAISE
        \epsilon \simeq 4 \cdot 1
        OPERAND METER HEAD NODE CURRENT TOKEN FROM RECORD OF A VEHICLER
        THE LARE TAPE - VARIABLE DECLARE
     ent tit
Asian TOKEN NUMERIC LITERAL
      J. C. PRENT TOKEN RECORD TOKEN TYPE - NUMBER THE SOL
       CONSEMI TRUE
       DECLARE TYPE CONSTANT DECLARE
       OPERAND METRICHEAD NODE CURRENT TOKEN RECORD DECLARE INTE
        DECLARE TYPE VARIABLE DECLARE.
        if HENRY WRITE ENABLE then
            WRITE HENRY DATA(LOCAL DECLARE HENRY TEXEME HENRY LEVEL 1999)
                                  NONE NEXT HEN).
           CREATE NODE(NEXT HEN LAST RECORD)
           HENRY WRITE ENABLE FALSE.
         end if
    end if
when TOKEN CHARACTER LITERAL
    of (CURRENT TOKEN RECORD TOKEN TYPE CHARACTER 111) then
       if HENRY WRITE ENABLE then
          WRITE HENRY DATA(LOCAL DECLARE, HENRY TEXEME, IDENT TYPE
                                NONE, NEXT, HEN);
           CREATE NODE(NEXT HEN. LAST RECORD)
           HENRY WRITE ENABLE FALSE.
        end if.
       CONSUME: TRUE.
    end if.
```

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 $\begin{array}{lll} & & & \text{In } F(F, X) = X \circ F \\ & & & \text{Normal } F(X) = X \circ F(X) \circ F(X)$

 $K = \{ (X, \theta, \theta) : X \in M(\theta) : x$

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* KENTELL

V STATEXEMETEXEME SIZE ("for" then
NSTAIL TRUE

Wise Light Note (THERS)

NOTE A THE NEW FOREST THEN SELECTION OF THE SELEC

when TOKEN RETURN

f ADJUST LEXEMF(LEXEME SIZE) = "return") then

constMF = TRUE;

end if

when TOKEN EXIT

if (ADJUST LEXEME(LEXEME, SIZE) = "exit") then
CONSUME = TRUE,
end if.

when TOKEN PROCEDURE

if (ADJUST LEXEME(LEXEME, SIZE) = "procedure") then
CONSUME TRUE;
end if.

when TOKEN FUNCTION ...

if (ADJUST LEXEME(LEXEME, SIZE) - "function") then
CONSUME - TRUE,
end if:

when TOKEN WITH --if (ADJUST LEXEME(LEXEME, SIZE) = "with") then
CONSUME: TRUE;
end if.

when TOKEN USE if (ADJUST LEXEME(LEXEME, SIZE) = "use") then CONSUME = TRUE; end if:

```
when TOKEN PACKAGE ---
  if (ADJUST LEXEME(LEXEME, SIZE) "package") then
   CONSUME TRUE.
  end if.
when TOKEN BODY =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "body") then
   CONSUME : TRUE;
  end if:
when TOKEN RANGE =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "range") then
   CONSUME := TRUE;
  end if:
when TOKEN IN = >
  if (ADJUST LEXEME(LEXEME, SIZE) = "in") then
   CONSUME := TRUE;
  end if:
when TOKEN OUT =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "out") then
   CONSUME := TRUE;
  end if:
when TOKEN SUBTYPE
  if (ADJUST LEXEME(LEXEME, SIZE) - "subtype") then
   CONSUME := TRUE:
  end if:
when TOKEN TYPE = -
  if (ADJUST LEXEME(LEXEME, SIZE) "type") then
  CONSUME: TRUE:
  end if:
when TOKEN IS
  if (ADJUST LEXEME(LEXEME SIZE) "is") then
  CONSUMÉ TRUE.
  end if:
when TOKEN NULL
  if (ADJUST LEXEME(LEXEME SIZE) - Fig. 1) ther
  CONSUME TRUE
  end if.
when TOKEN ACCESS
  if (ADJUST LEXEME) LEXEME SIZE
  CONSUME. TRUE
  end if
```

when TOKEN ARRAY

if (ADJUST LEXEME LEXEME SIZE

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where the KNN content A is a North-Half-ME size of table them.

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 $(\kappa + \gamma_{k}) = (\gamma_{k} + \gamma_{k}) \cdot (\gamma_{k} + \gamma_{k})$

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 $s = -s^{2} + (SEME)^{2}EXEME, SDE = (1) (x) experts <math display="inline">n^{2}$), then

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The state of the s

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CONTRACTOR

Car No. 1884

KEN AMETER

- 1-1 ANEMESTEME SIZE) "limited") then

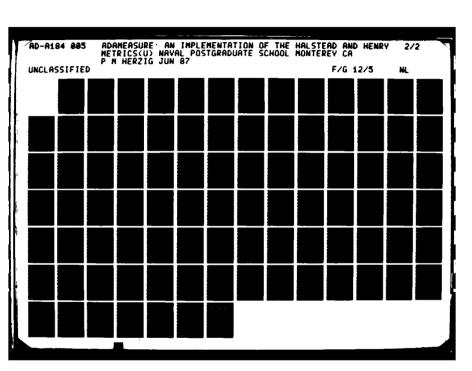
- NO TRUE

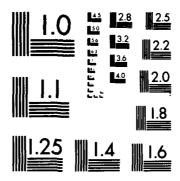
KIN INSK

* MISUST TEXEME(LEXEME, SIZE) = "task") then

ONSUME TRUE

. . .





MICROCOPY RESOLUTION TEST CHART
NATIONAL BUREAU OF STANDARDS-1963-A

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when TOKEN ENTRY =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "entry") then
   CONSUME := TRUE;
  end if;
when TOKEN ACCEPT =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "accept") then
   CONSUME := TRUE;
  end if:
when TOKEN DELAY =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "delay") then
   CONSUME := TRUE;
  end if;
when TOKEN SELECT =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "select") then
   CONSUME := TRUE;
  end if:
when TOKEN TERMINATE =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "terminate") then
   CONSUME := TRUE;
  end if:
when TOKEN ABORT =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "abort") then
   CONSUME := TRUE;
  end if;
when TOKEN SEPARATE =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "separate") then
   CONSUME := TRUE;
  end if:
when TOKEN RAISE =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "raise") then
   CONSUME := TRUE:
  end if:
when TOKEN GENERIC =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "generic") then
   CONSUME := TRUE;
  end if:
when TOKEN AT =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "at") then
   CONSUME := TRUE;
  end if;
when TOKEN REVERSE =>
```

```
if (ADJUST LEXEME(LEXEME, SIZE) = "reverse") then
   CONSUME := TRUE;
  end if:
when TOKEN DO =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "do") then
   CONSUME := TRUE:
  end if;
when TOKEN GOTO =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "goto") then
   CONSUME := TRUE;
  end if;
when TOKEN OF =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "of") then
   CONSUME := TRUE;
  end if;
when TOKEN ALL =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "all") then
   CONSUME := TRUE;
  end if;
when TOKEN PRAGMA =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "pragma") then
   CONSUME := TRUE;
  end if;
when TOKEN AND =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "and") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN AND. CONSUME, RESERVE WORD TEST);
when TOKEN OR =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "or") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN OR, CONSUME, RESERVE WORD TEST);
when TOKEN NOT =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "not") then
   CONSUME := TRUE;
  end if:
  OPERATOR METRIC(TOKEN NOT, CONSUME, RESERVE WORD TEST):
when TOKEN XOR =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "xor") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN XOR. CONSUME, RESERVE WORD TEST);
```

```
when TOKEN MOD =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "mod") then
   CONSUME := TRUE;
  end if;
  OPERATOR METRIC(TOKEN MOD, CONSUME, RESERVE WORD TEST);
when TOKEN REM =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "rem") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN REM, CONSUME, RESERVE WORD TEST);
when TOKEN ABSOLUTE =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "abs") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN ABSOLUTE, CONSUME, RESERVE WORD TEST);
when TOKEN ASTERISK =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "*") then
   CONSUME := TRUE;
  end if;
  OPERATOR METRIC(TOKEN ASTERISK, CONSUME, RESERVE WORD TEST);
when TOKEN SLASH =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "/") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN SLASH, CONSUME, RESERVE WORD TEST);
when TOKEN EXPONENT =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "**") then
   CONSUME := TRUE;
  end if:
  OPERATOR METRIC(TOKEN EXPONENT, CONSUME, RESERVE WORD TEST);
when TOKEN PLUS =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "+") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN PLUS, CONSUME, RESERVE WORD TEST):
when TOKEN MINUS ->
  if (ADJUST LEXEME(LEXEME, SIZE) = "-") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN MINUS, CONSUME, RESERVE WORD TEST):
when TOKEN AMPERSAND ->
  if (ADJUST LEXEME(LEXEME, SIZE) = "&") then
   CONSUME :- TRUE;
```

```
end if:
  OPERATOR METRIC(TOKEN AMPERSAND, CONSUME, RESERVE WORD TEST);
when TOKEN EQUALS =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "=") then
   CONSUME := TRUE;
  end if;
  OPERATOR METRIC(TOKEN EQUALS, CONSUME, RESERVE WORD TEST);
when TOKEN NOT EQUALS =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "/=") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN NOT EQUALS, CONSUME, RESERVE WORD TEST);
when TOKEN LESS THAN =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "<") then
   CONSUME := TRUE;
  end if:
  OPERATOR_METRIC(TOKEN LESS THAN, CONSUME, RESERVE WORD TEST);
when TOKEN LESS THAN EQUALS =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "<=") then
   CONSUME := TRUE;
  end if:
  OPERATOR METRIC(TOKEN LESS THAN EQUALS, CONSUME, RESERVE WORD TEST)
when TOKEN GREATER THAN =>
  if (ADJUST LEXEME(LEXEME, SIZE) = ">") then
   CONSUME := TRUE;
  OPERATOR METRIC(TOKEN GREATER THAN, CONSUME, RESERVE WORD TEST):
when TOKEN GREATER THAN EQUALS =>
  if (ADJUST LEXEME(LEXEME, SIZE) = ">=") then
  CONSUME := TRUE;
  end if:
  OPERATOR METRIC(TOKEN GREATER THAN EQUALS, CONSUME, RESERVE WORD 1
when TOKEN ASSIGNMENT =>
  if (ADJUST LEXEME(LEXEME, SIZE) = ":=") then
   CONSUME := TRUE;
   OPERATOR METRIC(TOKEN ASSIGNMENT, CONSUME, RESERVE WORD TEST);
  end if;
when TOKEN COMMA =>
  if (ADJUST LEXEME(LEXEME. SIZE) = ",") then
   CONSUME := TRUE;
  end if:
when TOKEN SEMICOLON =>
  if (ADJUST LEXEME(LEXEME, SIZE) = ";") then
```

```
UPDATE LINE COUNT;
   CONSUME := TRUE;
  end if:
when TOKEN PERIOD =>
  if (ADJUST LEXEME(LEXEME, SIZE) = ".") then
   CONSUME := TRUE:
  end if;
when TOKEN LEFT PAREN =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "(") then
   CONSUME := TRUE:
  end if:
when TOKEN RIGHT PAREN =>
  if (ADJUST LEXEME(LEXEME, SIZE) = ")") then
   CONSUME := TRUE:
  end if;
when TOKEN COLON =>
  if (ADJUST LEXEME(LEXEME, SIZE) = ":") then
   CONSUME := TRUE;
  end if;
when TOKEN APOSTROPHE =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "'") then
   CONSUME := TRUE;
  end if;
when TOKEN RANGE DOTS =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "..") then
   CONSUME := TRUE:
  end if:
when TOKEN ARROW =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "=>") then
   CONSUME := TRUE;
  end if:
when TOKEN BAR =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "|") then
   CONSUME := TRUE;
  end if;
when TOKEN BRACKETS =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "<>") then
   CONSUME := TRUE;
  end if;
when TOKEN LEFT BRACKET =>
  if (ADJUST LEXEME(LEXEME, SIZE) = "<<") then
   CONSUME := TRUE;
```

```
end if;
  when TOKEN RIGHT BRACKET =>
    if (ADJUST LEXEME(LEXEME, SIZE) = ">>") then
     CONSUME := TRUE;
    end if:
  when others => null;
 end case;
 ADJUST TOKEN BUFFER(CONSUME, RESERVE WORD TEST);
 return (CONSUME);
end BYPASS:
 -- this procedure tests all identifiers to verify they are not reserved
 -- words. The most common reserved words are tested first and the process
 -- halts when a match is made or the test fails.
procedure CONDUCT RESERVE WORD TEST(CONSUME : in out boolean) is
begin
 RESERVE WORD TEST := TRUE;
 for RESERVE WORD INDEX in TOKEN END. TOKEN ABSOLUTE loop
  if (BYPASS(RESERVE WORD INDEX)) then
   CONSUME := FALSE;
  end if:
  exit when not CONSUME;
 end loop;
 RESERVE WORD TEST := FALSE;
end CONDUCT RESERVE WORD TEST;
end BYPASS FUNCTION;
              AN ADA SOFTWARE METRIC
-- TITLE:
-- MODULE NAME: PACKAGE PARSER 0
-- DATE CREATED: 09 OCT 86
-- LAST MODIFIED: 30 MAY 87
-- AUTHORS:
                LCDR JEFFREY L. NIEDER
           LT KARL S. FAIRBANKS, JR.
           LCDR PAUL M. HERZIG
-- DESCRIPTION: This package contains eight functions that
     make up the highest level productions for our top-down.
     recursive descent parser.
```

```
with PARSER 1, PARSER 2, PARSER 3, HENRY GLOBAL, HENRY, BYPASS FUNCTION,
   HALSTEAD METRIC, GLOBAL PARSER, GLOBAL, TEXT 10:
use PARSER 1, PARSER 2, PARSER 3, HENRY GLOBAL, HENRY, BYPASS FUNCTION,
  HALSTEAD METRIC, GLOBAL PARSER, GLOBAL, TEXT 10;
package PARSER 0 is
 function COMPILATION return boolean;
 function COMPILATION UNIT return boolean;
 function CONTEXT CLAUSE return boolean;
 function BASIC UNIT return boolean;
 function LIBRARY UNIT return boolean;
 function SECONDARY UNIT return boolean;
 function LIBRARY UNIT BODY return boolean;
 function SUBUNIT return boolean:
end PARSER_0;
package body PARSER 0 is
 -- COMPILATION --> COMPILATION UNIT +
function COMPILATION return boolean is
begin
put("In compilation "); new line;
put(RESULT FILE. "In compilation "); new line(RESULT FILE);
 if (COMPILATION UNIT) then
  while (COMPILATION UNIT) loop
   null;
  end loop:
  return (TRUE):
  return (FALSE);
 end if:
end COMPILATION;
 -- COMPILATION UNIT --> CONTEXT CLAUSE BASIC UNIT
function COMPILATION UNIT return boolean is
put(RESULT FILE. "In compilation unit "); new line(RESULT FILE);
 if (CONTEXT CLAUSE) then
   if (BASIC UNIT) then
    return (TRUE);
    return (FALSE);
   end if:
 else
  return (FALSE);
 end if:
end COMPILATION UNIT:
```

```
-- CONTEXT CLAUSE --> [with WITH OR USE CLAUSE | use WITH OR USE CLAUSE | * | *
function CONTEXT CLAUSE return boolean is
put(RESULT FILE, "In context clause"); new line(RESULT FILE);
 while (BYPASS(TOKEN WITH)) loop
  if not (WITH OR USE CLAUSE) then
   SYNTAX ERROR("Context clause");
  end if:
  while (BYPASS(TOKEN USE)) loop
   if not (WITH OR USE CLAUSE) then
    SYNTAX ERROR("Context clause");
   end if;
  end loop;
                           -- inner while loop
 end loop:
                                  -- outer while loop
 return (TRUE);
end CONTEXT CLAUSE;
 -- BASIC UNIT --> LIBRARY UNIT
             --> SECONDARY UNIT
function BASIC UNIT return boolean is
put(RESULT FILE, "In basic unit"); new line(RESULT FILE);
 if (LIBRARY UNIT) then
  return (TRUE);
 elsif (SECONDARY UNIT) then
  return (TRUE);
 else
  return (FALSE);
 end if:
end BASIC UNIT;
 -- LIBRARY UNIT --> procedure PROCEDURE UNIT
              --> function FUNCTION UNIT
              --> package PACKAGE DECLARATION
              --> generic GENERIC DECLARATION
function LIBRARY UNIT return boolean is
put(RESULT FILE. "In library unit "); new line(RESULT FILE);
 if (BYPASS(TOKEN PROCEDURE)) then
  DECLARE TYPE := PROCEDURE DECLARE;
  if (PROCEDURE UNIT) then
   return (TRUE);
  else
   SYNTAX ERROR("Library unit"):
                                  -- if procedure unit statement
 elsif (BYPASS(TOKEN FUNCTION)) then
```

```
DECLARE TYPE := FUNCTION DECLARE;
  if (FUNCTION UNIT) then
   return (TRUE);
  else
   SYNTAX ERROR("Library unit");
                                 -- if function unit statement
 elsif (BYPASS(TOKEN PACKAGE)) then
  DECLARE TYPE := PACKAGE DECLARE;
  if (PACKAGE DECLARATION) then
   return (TRUE);
   SYNTAX ERROR("Library unit");
                                 -- if package declaration
 elsif (BYPASS(TOKEN GENERIC)) then
  if (GENERIC DECLARATION) then
   return (TRUE);
  else
   SYNTAX ERROR("Library unit");
  end if:
                                  -- if generic declaration
 else
  return (FALSE):
 end if:
end LIBRARY UNIT;
 -- SECONDARY UNIT --> LIBRARY UNIT BODY
               --> SUBUNIT
function SECONDARY UNIT return boolean is
put(RESULT FILE, "In secondary unit"); new line(RESULT FILE);
 if (LIBRARY UNIT BODY) then
  return (TRUE);
 elsif (SUBUNIT) then
  return (TRUE);
 else
  return (FALSE):
 end if:
end SECONDARY UNIT;
 -- LIBRARY UNIT BODY --> procedure PROCEDURE UNIT
                 --> function FUNCTION UNIT
                 --> package PACKAGE DECLARATION
                 --> generic GENERIC DECLARATION
function LIBRARY UNIT BODY return boolean is
put(RESULT FILE, "In library unit body "); new line(RESULT FILE);
 if (BYPASS(TOKEN PROCEDURE)) then
  DECLARE TYPE := PROCEDURE DECLARE:
  if (PROCEDURE UNIT) then
```

```
return (TRUE);
  else
  SYNTAX ERROR("Library unit body");
                                    -- if procedure unit statement
 elsif (BYPASS(TOKEN FUNCTION)) then
  DECLARE TYPE := FUNCTION DECLARE;
  if (FUNCTION UNIT) then
    return (TRUE);
  else
    SYNTAX ERROR("Library unit body");
                                    -- if function unit statement
 elsif (BYPASS(TOKEN PACKAGE)) then
  DECLARE TYPE := \overline{P}ACKAGE DECLARE;
  HENRY WRITE ENABLE := TRUE;
  put(result file, "true"); new line(result file);
  if (PACKAGE DECLARATION) then
   return (TRUE);
  else
   SYNTAX ERROR("Library unit body");
                                   -- if package_declaration
  return (FALSE);
 end if;
                            -- if bypass(token procedure)
end LIBRARY UNIT BODY;
 -- SUBUNIT --> separate (NAME) PROPER BODY
function SUBUNIT return boolean is
put(RESULT FILE, "In subunit "); new line(RESULT FILE);
 if (BYPASS(TOKEN SEPARATE)) then
  if (BYPASS(TOKEN_LEFT_PAREN)) then
   if (NAME) then
    if (BYPASS(TOKEN RIGHT PAREN)) then
      if (PROPER BODY) then
       return (TRUE);
     else
       SYNTAX ERROR("Subunit");
                            -- if proper body statement
     SYNTAX ERROR("Subunit");
    end if;
                                   -- if bypass(token right paren)
   else
    SYNTAX ERROR("Subunit");
   end if:
                                   -- if name statement
 else
   SYNTAX ERROR("Subunit");
 end if;
                                   -- if bypass(token left paren)
 return (FALSE);
end if;
                            -- if bypass(token separate)
```

```
end SUBUNIT;
end PARSER 0;
  TITLE:
              AN ADA SOFTWARE METRIC
-- MODULE NAME: PACKAGE PARSER 1
-- DATE CREATED: 17 JUL 86
-- LAST MODIFIED: 30 MAY 87
                LCDR JEFFREY L. NIEDER
-- AUTHORS:
           LT KARL S. FAIRBANKS, JR.
           LCDR PAUL M. HERZIG
-- DESCRIPTION: This package contains thirty-six functions
     that make up the top level productions for our top-down,
     recursive descent parser. Each function is preceded
     by the grammar productions they are implementing.
with PARSER 2, PARSER 3, HENRY GLOBAL, HENRY, BYPASS FUNCTION.
   HALSTEAD METRIC, GLOBAL PARSER, GLOBAL, TEXT 10;
use PARSER 2. PARSER 3. HENRY GLOBAL, HENRY, BYPASS FUNCTION,
  HALSTEAD METRIC. GLOBAL PARSER, GLOBAL, TEXT 10;
package PARSER 1 is
 function GENERIC DECLARATION return boolean;
 function GENERIC PARAMETER DECLARATION return boolean;
 function GENERIC FORMAL PART return boolean:
 function PROCEDURE UNIT return boolean;
 function SUBPROGRAM BODY return boolean:
 function FUNCTION UNIT return boolean;
 function FUNCTION UNIT TAIL return boolean:
 function FUNCTION BODY return boolean;
 function FUNCTION BODY TAIL return boolean;
 function TASK DECLARATION return boolean;
 function TASK BODY return boolean;
 function TASK BODY TAIL return boolean:
 function PACKAGE DECLARATION return boolean;
 function PACKAGE UNIT return boolean;
 function PACKAGE BODY return boolean;
 function PACKAGE BODY TAIL return boolean;
 function PACKAGE TAIL END return boolean;
 function DECLARATIVE PART return boolean:
 function BASIC DECLARATIVE ITEM return boolean;
 function BASIC DECLARATION return boolean;
 function LATER DECLARATIVE ITEM return boolean:
 function PROPER BODY return boolean:
```

```
function SEQUENCE OF STATEMENTS return boolean;
function STATEMENT return boolean;
function COMPOUND STATEMENT return boolean;
function BLOCK STATEMENT return boolean:
function IF STATEMENT return boolean:
function CASE STATEMENT return boolean:
function CASE STATEMENT ALTERNATIVE return boolean;
function LOOP STATEMENT return boolean;
function EXCEPTION HANDLER return boolean;
function ACCEPT STATEMENT return boolean;
function SELECT STATEMENT return boolean;
function SELECT STATEMENT TAIL return boolean;
function SELECT ALTERNATIVE return boolean;
function SELECT ENTRY CALL return boolean;
end PARSER 1:
package body PARSER 1 is
-- GENERIC DECLARATION --> GENERIC PARAMETER DECLARATION?
                           GENERIC FORMAL PART
function GENERIC DECLARATION return boolean is
put(RESULT FILE, "In generic declaration"); new line(RESULT FILE);
 if (GENERIC PARAMETER DECLARATION) then
  null:
 end if:
 if (GENERIC FORMAL PART) then
  return(TRUE);
 else
  return (FALSE);
 end if;
end GENERIC DECLARATION;
 -- GENERIC PARAMETER DECLARATION --> IDENTIFIER LIST: MODE? NAME
                                  := EXPRESSION ?!;
                            --> type private DISCRIMINANT PART?
                                  is PRIVATE TYPE DECLARATION;
                                type private DISCRIMINANT PART?
                                  is GENERIC TYPE DEFINITION;
                            --> with procedure PROCEDURE UNIT
                            --> with function FUNCTION UNIT
function GENERIC PARAMETER DECLARATION return boolean is
put(RESULT FILE, "In generic parameter declaration"); new line(RESULT FILE);
 if (IDENTIFIER LIST) then
  if (BYPASS(TOKEN COLON)) then
   if (MODE) then
```

```
null;
  end if;
                                   -- if mode statement
  if (NAME) then
                                   -- check for type mark
   if (BYPASS(TOKEN ASSIGNMENT)) then
    if (EXPRESSION) then
      null;
    else
      SYNTAX ERROR("Generic parameter declaration");
    end if:
                         -- if expression statement
   end if:
                         -- if bypass(token assignment)
   if (BYPASS(TOKEN SEMICOLON)) then
    return (TRUE);
   else
    SYNTAX ERROR("Generic parameter declaration");
                         -- if bypass(token semicolon)
   end if:
  else
   SYNTAX ERROR("Generic parameter declaration");
  end if;
                         -- if type mark statement
 else
  SYNTAX ERROR("Generic parameter declaration");
                         -- if bypass(token colon)
elsif (BYPASS(TOKEN TYPE)) then
 if (BYPASS(TOKEN IDENTIFIER)) then
  if (DISCRIMINANT PART) then
   null;
  end if:
                                   -- if discriminant part
  if (BYPASS(TOKEN IS)) then
   if (PRIVATE TYPE DECLARATION) then
     if (BYPASS(TOKEN SEMICOLON)) then
      return (TRUE);
     else
      SYNTAX ERROR("Generic parameter declaration");
                            -- if bypass(token semicolon)
     end if;
   elsif (GENERIC TYPE DEFINITION) then
     if (BYPASS(TOKEN SEMICOLON)) then
      return (TRUE);
      SYNTAX ERROR("Generic parameter declaration");
                            -- if bypass(token semicolon)
    SYNTAX ERROR("Generic parameter declaration");
   end if;
                                   -- if private type declaration
  else
   SYNTAX ERROR("Generic parameter declaration");
  end if:
                                   -- if bypass(token is)
 else
  SYNTAX ERROR("Generic parameter declaration");
                                   -- if bypass(token identifier)
elsif (BYPASS(TOKEN WITH)) then
 if (BYPASS(TOKEN PROCEDURE)) then
  DECLARE TYPE := PROCEDURE DECLARE;
  if (PROCEDURE UNIT) then
```

```
return (TRUE);
    else
     SYNTAX ERROR("Generic parameter declaration");
                                   -- if procedure unit statement
   elsif (BYPASS(TOKEN FUNCTION)) then
    DECLARE TYPE := FUNCTION DECLARE;
    if (FUNCTION UNIT) then
     return (TRUE);
    else
     SYNTAX ERROR("Generic parameter declaration");
    end if:
                                   -- if function unit statement
    SYNTAX ERROR("Generic parameter declaration");
   end if:
                                   - if bypass(token_procedure)
 else
  return (FALSE);
                         -- if identifier list
end GENERIC PARAMETER DECLARATION;
 -- GENERIC FORMAL PART --> procedure PROCEDURE UNIT
                     --> function FUNCTION UNIT
                     --> package PACKAGE DECLARATION
function GENERIC FORMAL PART return boolean is
put(RESULT_FILE, "In generic formal part"); new_line(RESULT_FILE);
 if (BYPASS(TOKEN PROCEDURE)) then
  DECLARE TYPE := PROCEDURE DECLARE:
  if (PROCEDURE UNIT) then
   return (TRUE);
  else
    SYNTAX ERROR("Generic formal part");
                                   -- if procedure unit statement
 elsif (BYPASS(TOKEN FUNCTION)) then
  DECLARE TYPE := \overline{FUNCTION} DECLARE;
  if (FUNCTION UNIT) then
   return (TRUE);
  else
   SYNTAX ERROR("Generic formal part");
                                   -- if function unit statement
 elsif (BYPASS(TOKEN PACKAGE)) then
  DECLARE TYPE := \overline{P}ACKAGE DECLARE:
  if (PACKAGE DECLARATION) then
   return (TRUE);
   SYNTAX ERROR("Generic formal part");
  end if:
                                  -- if package declaration
 else
  return (FALSE);
 end if;
end GENERIC FORMAL PART:
```

```
-- PROCEDURE UNIT --> identifier [FORMAL PART ?] is SUBPROGRAM BODY
              --> identifier FORMAL PART ? ;
               --> identifier [FORMAL PART ?] renames NAME;
function PROCEDURE UNIT return boolean is
put(RESULT FILE, "In procedure unit"); new line(RESULT FILE);
 DECLARATION := TRUE;
  HENRY WRITE ENABLE := TRUE;
 if (BYPASS(TOKEN IDENTIFIER)) then
  if PACKAGE BODY DECLARE then
   WRITE HENRY DATA(LOCAL DECLARE, DUMMY LEXEME,
            PROCEDURE TYPE, NONE, LAST RECORD);
  end if:
  SCOPE LEVEL := SCOPE LEVEL + 1;
  if (FORMAL PART) then
   null:
  end if:
                         -- if formal part statement
  if (BYPASS(TOKEN IS)) then
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, END PARAM DECLARE,
            NONE, NEXT HEN);
   CREATE NODE(NEXT HEN, LAST RECORD);
   WRITE LINE COUNT(LAST RECORD.NOMEN, HENRY LINE COUNT.
            DUMMY9s, NEXT LINE);
   if (SUBPROGRAM BODY) then
    WRITE HENRY DATA(BLANK, DUMMY LEXEME, END PROCEDURE CALL,
              NONE. NEXT HEN);
    CREATE NODE(NEXT HEN, LAST RECORD);
    WRITE LINE COUNT(DUMMY LEXEME, DUMMY9s, HENRY LINE COUNT.
              NEXT LINE);
    CREATE LINE COUNT NODE(NEXT LINE, LAST LINE);
    SCOPE LEVEL := SCOPE LEVEL - 1;
    return (TRUE);
   else
    SYNTAX ERROR("Procedure unit"):
                         -- if subprogram body statement
  elsif (BYPASS(TOKEN SEMICOLON)) then
   SCOPE LEVEL := SCOPE LEVEL - 1;
   return (TRUE);
  elsif (BYPASS(TOKEN RENAMES)) then
   if (NAME) then
    if (BYPASS(TOKEN SEMICOLON)) then
     SCOPE LEVEL := SCOPE LEVEL - 1;
     return (TRUE);
      SYNTAX ERROR("Procedure unit");
    end if:
                         -- if bypass(token semicolon)
   else
    SYNTAX ERROR("Procedure unit");
                         -- if name statement
   end if;
```

```
end if:
                          -- if bypass(token is)
 else
  return (FALSE);
 end if:
                             -- if bypass(token identifier)
end PROCEDURE UNIT;
 -- SUBPROGRAM BODY --> new NAME [GENERIC ACTUAL PART?];
                --> separate;
                --> <>;
                --> DECLARATIVE PART? begin SEQUENCE OF STATEMENTS
                    [exception [EXCEPTION HANDLER] + ? end [DESIGNATOR ?];
                --> NAME ;
function SUBPROGRAM BODY return boolean is
NAME POINTER: POINTER;
begin
put(RESULT FILE. "In subprogram body "); new line(RESULT FILE);
NAME POINTER := NEXT HEN;
 DECLARATION := TRUE;
 if (BYPASS(TOKEN NEW))then
  HENRY WRITE ENABLE := FALSE;
  if (NAME) then
   if (GENERIC ACTUAL PART) then
    null:
   end if;
                         -- if generic actual part
   if (BYPASS(TOKEN SEMICOLON)) then
    return (TRUE);
   else
    SYNTAX ERROR("Subprogram body");
   end if:
                         -- if bypass(token semicolon)
  else
   SYNTAX ERROR("Subprogram body");
                         -- if name statement
 elsif (BYPASS(TOKEN SEPARATE)) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
  else
   SYNTAX ERROR("Subprogram body");
                         -- if bypass(token semicolon)
 elsif (BYPASS(TOKEN BRACKETS)) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
   SYNTAX ERROR("Subprogram body"):
  end if:
                         -- if bypass(token semicolon)
 elsif (DECLARATIVE PART) then
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, END DECLARATIONS,
            NONE, NEXT HEN):
  CREATE NODE(NEXT HEN, LAST RECORD);
```

```
if (BYPASS(TOKEN BEGIN)) then
  DECLARATION := FALSE:
  if (SEQUENCE OF STATEMENTS) then
   if (BYPASS(TOKEN EXCEPTION)) then
     if (EXCEPTION HANDLER) then
      while (EXCEPTION HANDLER) loop
       null:
      end loop;
    else
      SYNTAX ERROR("Subprogram body");
                           -- if exception handler statement
   end if:
                                  -- if bypass(token exception)
   if (BYPASS(TOKEN END)) then
      HENRY WRITE ENABLE := FALSE;
    if (DESIGNATOR) then
      null;
    end if:
                         -- if designator statement
    if (BYPASS(TOKEN SEMICOLON)) then
      DECLARATION := TRUE:
      return (TRUE);
    else
      SYNTAX ERROR("Subprogram body");
    end if;
                         -- if bypass(token semicolon)
    SYNTAX ERROR("Subprogram body");
   end if:
                         -- if bypass(token end)
  else
   SYNTAX ERROR("Subprogram body");
  end if;
                        -- if sequence of statements
 else
  SYNTAX ERROR("Subprogram body");
                         -- if bypass(token begin)
elsif (BYPASS(TOKEN BEGIN)) then
 DECLARATION := \overline{FALSE};
 WRITE HENRY DATA(BLANK, DUMMY LEXEME, END DECLARATIONS,
           NONE. NEXT HEN);
CREATE NODE(NEXT HEN, LAST RECORD);
if (SEQUENCE OF STATEMENTS) then
  if (BYPASS(TOKEN EXCEPTION)) then
   if (EXCEPTION HANDLER) then
    while (EXCEPTION HANDLER) loop
     null:
    end loop;
    SYNTAX ERROR("Subprogram body");
   end if:
                                  -- if exception handler statement
  end if:
                                  -- if bypass(token exception)
  if (BYPASS(TOKEN END)) then
   HENRY WRITE ENABLE := FALSE;
   if (DESIGNATOR) then
    null:
   end if:
                        -- if designator statement
```

```
if (BYPASS(TOKEN SEMICOLON)) then
      DECLARATION := TRUE;
      return (TRUE):
     else
      SYNTAX ERROR("Subprogram body"):
                          -- if bypass(token semicolon)
    else
     SYNTAX ERROR("Subprogram body");
                          -- if bypass(token end)
   else
    SYNTAX ERROR("Subprogram body");
   end if:
                          -- if sequence of statements
  elsif (NAME) then
   if (BYPASS(TOKEN SEMICOLON)) then
    return (TRUE);
  else
    SYNTAX ERROR("Subprogram body");
  end if:
                          -- if bypass(token semicolon)
 else
  return (FALSE);
 end if:
                          -- if bypass(token new)
end SUBPROGRAM BODY;
 -- FUNCTION UNIT --> DESIGNATOR FUNCTION UNIT TAIL
function FUNCTION UNIT return boolean is
put(RESULT FILE. "In function unit "); new line(RESULT FILE);
 DECLARATION := TRUE:
 HENRY WRITE ENABLE := TRUE;
 if (DESIGNATOR) then
  if PACKAGE BODY DECLARE then
    WRITE HENRY DATA(LOCAL DECLARE, DUMMY LEXEME, FUNCTION TYPE,
              NONE, LAST RECORD):
    WRITE LINE COUNT(LAST RECORD.NOMEN, HENRY LINE COUNT.
              DUMMY9s, NEXT LINE);
  end if:
  SCOPE LEVEL := SCOPE LEVEL + 1;
  if (FUNCTION UNIT TAIL) then
   SCOPE LEVEL := SCOPE LEVEL - 1;
   return (TRUE):
  else
   SYNTAX ERROR("Function unit");
  end if;
 else
  return (FALSE):
 end if:
end FUNCTION UNIT:
```

```
-- FUNCTION UNIT TAIL --> is new NAME GENERIC ACTUAL PART?];
                      --> FORMAL PART ? return NAME FUNCTION BODY
function FUNCTION UNIT TAIL return boolean is
begin
put(RESULT FILE. "In function unit tail "); new line(RESULT FILE);
 if (BYPASS(TOKEN IS)) then
     FUNCTION PARAM DECLARE := TRUE;
   if (BYPASS(TOKEN NEW)) then
    if (NAME) then
     if (GENERIC ACTUAL PART) then
      null;
     end if:
                           -- if generic actual part
     if (BYPASS(TOKEN SEMICOLON)) then
      return (TRUE);
     else
      SYNTAX ERROR("Function unit tail");
     end if;
                           -- if bypass(token semicolon)
    else
     SYNTAX ERROR("Function unit tail");
                           -- if name statement
    end if;
   else
    SYNTAX ERROR("Function unit tail");
  end if;
                           -- if bypass(token new)
 elsif (FORMAL PART) then
   FUNCTION PARAM DECLARE := FALSE:
  if (BYPASS(TOKEN RETURN)) then
   if (NAME) then
                                     -- check for type mark
     if (FUNCTION BODY) then
      return (TRUE);
      SYNTAX ERROR("Function unit tail");
     end if;
                           -- if function body statement
   else
     SYNTAX ERROR("Function unit tail");
   end if;
                           -- if type mark statement
  else
   SYNTAX ERROR("Function unit tail");
                           -- if bypass(token return)
 elsif (BYPASS(TOKEN RETURN)) then
  if (NAME) then
                                     -- check for type mark
   if (FUNCTION BODY) then
     return (TRUE);
     SYNTAX ERROR("Function unit tail");
   end if:
                           -- if function body statement
   SYNTAX ERROR("Function unit tail");
  end if;
                           -- if type mark statement
 else
  return (FALSE):
                           -- if bypass(token is)
end FUNCTION UNIT TAIL:
```

```
-- FUNCTION BODY --> is FUNCTION BODY TAIL?
function FUNCTION BODY return boolean is
put(RESULT_FILE, "In function body"); new line(RESULT_FILE);
 if (BYPASS(TOKEN IS)) then
  WRITE HENRY DATA(BLANK, DUMMY LEXEME, END PARAM DECLARE, NONE, NEXT I
  CREATE NODE(NEXT HEN, LAST RECORD);
  if (FUNCTION BODY TAIL) then
   WRITE LINE COUNT(DUMMY LEXEME, DUMMY9s, HENRY LINE COUNT.
            NEXT LINE);
   CREATE LINE COUNT NODE(NEXT LINE, LAST LINE):
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, END FUNCTION TYPE,
             NONE, NEXT HEN):
   CREATE NODE(NEXT HEN, LAST RECORD);
  end if:
  return (TRUE):
 elsif (BYPASS(TOKEN SEMICOLON)) then
  return (TRUE):
 else
  return (FALSE):
 end if:
end FUNCTION BODY:
 -- FUNCTION BODY TAIL --> separate;
                   -- > < > :
                   --> SUBPROGRAM BODY
                   --> NAME:
function FUNCTION BODY TAIL return boolean is
put(RESULT_FILE, "In function body tail"); new line(RESULT_FILE);
 if (BYPASS(TOKEN SEPARATE)) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE):
  else
   SYNTAX ERROR("Function body tail");
                       -- if bypass(token semicolon)
 elsif (BYPASS(TOKEN BRACKETS)) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE):
  else
   SYNTAX ERROR("Function body tail");
                      -- if bypass(token semicolon)
 elsif (SUBPROGRAM BODY) then
  return (TRUE):
 elsif (NAME) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE):
```

```
else
   SYNTAX ERROR("Function body tail");
                        -- if bypass(token semicolon)
  end if:
  return (FALSE);
                        -- if bypass(token separate)
 end if:
end FUNCTION BODY TAIL;
 -- TASK DECLARATION --> body TASK BODY;
                 --> [type ?] identifier [is ENTRY DECLARATION *
                         REPRESENTATION CLAUSE * end identifier ? ? ? ;
function TASK DECLARATION return boolean is
begin
put(RESULT FILE, "In task declaration"); new line(RESULT FILE);
 DECLARATION := TRUE;
 if (BYPASS(TOKEN TYPE)) then
  null:
 end if:
                            -- if bypass(token type)
 if (BYPASS(TOKEN BODY)) then
  if (TASK BODY) then
   if (BYPASS(TOKEN SEMICOLON)) then
    return (TRUE);
    SYNTAX ERROR("Task declaration");
   end if:
  else
   SYNTAX ERROR("Task declaration"):
                                   -- if task body statement
 elsif (BYPASS(TOKEN IDENTIFIER)) then
  SCOPE LEVEL: SCOPE LEVEL - 1:
  if (BYPASS(TOKEN IS)) then
   while (ENTRY DECLARATION) loop
    null:
   end loop;
   while (REPRESENTATION CLAUSE) loop
    null:
   end loop:
   if (BYPASS(TOKEN_END)) then
     if (BYPASS(TOKEN IDENTIFIER)) then
      null:
    end if;
                                   -- if bypass(token identifier)
     if (BYPASS(TOKEN SEMICOLON)) then
      SCOPE LEVEL - SCOPE LEVEL - 1;
      return (TRUE);
    else
      SYNTAX ERROR("Task declaration").
    end if:
                                   - if bypass(token semicolon)
   else
    SYNTAX ERROR("Task declaration").
   end if.
                                   - if by pass (token end)
```

```
elsif (BYPASS(TOKEN SEMICOLON)) then
   SCOPE LEVEL := SCOPE LEVEL - 1;
   return (TRUE);
  else
   SYNTAX ERROR("Task declaration");
                                  -- if bypass(token is)
  end if;
 else
  return (FALSE);
 end if;
                           -- if bypass(token body)
end TASK DECLARATION;
 -- TASK BODY --> identifier is TASK BODY TAIL
function TASK BODY return boolean is
begin
put(RESULT FILE, "In task body"); new line(RESULT FILE);
 if (BYPASS(TOKEN IDENTIFIER)) then
  SCOPE LEVEL := SCOPE LEVEL + 1;
  if (BYPASS(TOKEN IS)) then
   if (TASK BODY TAIL) then
    SCOPE LEVEL := SCOPE LEVEL - 1:
    return (TRUE);
   else
    SYNTAX ERROR("Task body");
   end if:
                                  -- if task body tail statement
  else
   SYNTAX ERROR("Task body");
  end if;
                                  -- if bypass(token is)
  return (FALSE);
                           -- if bypass(token_identifier)
 end if;
end TASK BODY;
 -- TASK BODY TAIL --> separate
                --> DECLARATIVE PART? begin SEQUENCE OF STATEMENTS
                     exception EXCEPTION HANDLER + ? end identifier?
function TASK BODY TAIL return boolean is
put(RESULT FILE, "In task body tail"); new line(RESULT FILE);
 DECLARATION := TRUE;
 if (BYPASS(TOKEN SEPARATE)) then
  return (TRUE);
 elsif (DECLARATIVE PART) then
  if (BYPASS(TOKEN BEGIN)) then
   DECLARATION := FALSE;
    if (SEQUENCE OF STATEMENTS) then
     if (BYPASS(TOKEN EXCEPTION)) then
      if (EXCEPTION HANDLER) then
       while (EXCEPTION HANDLER) loop
```

```
null:
       end loop;
      else
       SYNTAX ERROR("Task body tail");
                             -- if exception handler statement
     end if:
                                     -- if bypass(token exception)
     if (BYPASS(TOKEN END)) then
      if (BYPASS(TOKEN IDENTIFIER)) then
       null;
      end if:
                             -- if bypass(token identifier)
      DECLARATION := TRUE;
      return (TRUE);
      SYNTAX ERROR("Task body tail");
     end if:
                             -- if bypass(token end)
   else
     SYNTAX ERROR("Task body tail");
   end if:
                             -- if sequence of statements
  else
    SYNTAX ERROR("Task body tail");
  end if:
                             -- if bypass(token begin)
 elsif (BYPASS(TOKEN BEGIN)) then
  DECLARATION := FALSE;
  if (SEQUENCE OF STATEMENTS) then
   if (BYPASS(TOKEN EXCEPTION)) then
     if (EXCEPTION HANDLER) then
      while (EXCEPTION HANDLER) loop
       null;
      end loop;
     else
      SYNTAX ERROR("Task body tail");
                                     -- if exception handler statement
     end if:
   end if:
                                     -- if bypass(token exception)
   if (BYPASS(TOKEN END)) then
     if (BYPASS(TOKEN IDENTIFIER)) then
      null;
     end if:
                             -- if bypass(token identifier)
     DECLARATION := TRUE;
     return (TRUE);
     SYNTAX ERROR("Task body tail");
   end if:
                             -- if bypass(token end)
  else
   SYNTAX ERROR("Task body tail");
  end if:
                             -- if sequence of statements
 else
  return (FALSE):
 end if:
                             -- if bypass(token separate)
end TASK BODY TAIL;
```

```
-- PACKAGE DECLARATION --> body PACKAGE BODY
                    --> identifier PACKAGE UNIT
function PACKAGE DECLARATION return boolean is
put(RESULT FILE, "In package declaration"); new line(RESULT FILE);
 DECLARATION := TRUE;
 HENRY WRITE ENABLE := TRUE;
 if (BYPASS(TOKEN BODY)) then
  PACKAGE BODY DECLARE := TRUE;
  HENRY WRITE ENABLE := FALSE;
  if (PACKAGE BODY) then
   return (TRUE):
   SYNTAX ERROR("Package declaration");
                          -- if package unit statement
  end if:
 elsif (BYPASS(TOKEN IDENTIFIER)) then
  WRITE HENRY DATA(LOCAL DECLARE, DUMMY LEXEME, PACKAGE TYPE,
            NONE, LAST RECORD):
  SCOPE LEVEL := SCOPE LEVEL + 1;
  if (PACKAGE UNIT) then
   SCOPE LEVEL := SCOPE LEVEL + 1;
   return (TRUE):
  else
   SYNTAX ERROR("Package declaration");
  end if;
                                 -- if package unit tail statement
 else
  return (FALSE);
 end if:
                          -- if bypass(token package)
end PACKAGE DECLARATION:
 -- PACKAGE BODY --> identifier is PACKAGE BODY TAIL
function PACKAGE BODY return boolean is
begin
put(RESULT FILE, "In package body "); new line(RESULT FILE);
 if (BYPASS(TOKEN IDENTIFIER)) then
  SCOPE LEVEL := SCOPE LEVEL + 1;
  if (BYPASS(TOKEN IS)) then
   if (PACKAGE BODY TAIL) then
     WRITE HENRY DATA(BLANK, DUMMY LEXEME, END PACKAGE TYPE,
              NONE, NEXT HEN):
    SCOPE LEVEL := SCOPE LEVEL - 1;
    return (TRUE):
   else
    SYNTAX ERROR("Package body"):
                    -- if package body tail statement
   end if;
   SYNTAX ERROR("Package body"):
  end if:
                    -- if bypass(token is)
 else
```

```
return (FALSE);
                    -- if bypass(token identifier)
 end if:
end PACKAGE BODY;
 -- PACKAGE BODY TAIL --> separate;
                 --> DECLARATIVE PART? [begin SEQUENCE OF STATEMENTS
                      exception EXCEPTION HANDLER - ? ?
                        end identifier?;
function PACKAGE_BODY_TAIL return boolean is
put(RESULT FILE, "In package body tail"); new line(RESULT FILE):
 DECLARATION := TRUE:
 if (BYPASS(TOKEN SEPARATE)) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
  else
   SYNTAX ERROR("Package body tail");
                           -- if bypass(token_semicolon)
  end if;
 elsif (DECLARATIVE PART) then
  DECLARATION := FALSE;
  if (BYPASS(TOKEN BEGIN)) then
   if (SEQUENCE OF STATEMENTS) then
    if (BYPASS(TOKEN EXCEPTION)) then
      if (EXCEPTION HANDLER) then
       while (EXCEPTION HANDLER) loop
        null;
       end loop;
      else
       SYNTAX ERROR("Package body tail");
      end if:
                           -- if exception handler statement
    end if;
                                  -- if bypass(token exception)
    if (BYPASS(TOKEN END)) then
       HENRY WRITE ENABLE := FALSE;
       if (BYPASS(TOKEN IDENTIFIER)) then
         null;
                           -- if bypass(token identifier)
      if (BYPASS(TOKEN SEMICOLON)) then
       DECLARATION := TRUE:
       return (TRUE);
       SYNTAX ERROR("Package body tail"):
      end if;
                           -- if bypass(token semicolon)
      SYNTAX ERROR("Package body tail");
    end if:
                           -- if bypass(token end)
   else
    SYNTAX ERROR("Package body tail"):
                            -- if sequence of statements
  elsif (BYPASS(TOKEN END)) then
    HENRY WRITE ENABLE: * FALSE:
```

```
if (BYPASS(TOKEN IDENTIFIER)) then
   null:
  end if:
                           -- if bypass(token identifier)
  if (BYPASS(TOKEN SEMICOLON)) then
   DECLARATION := TRUE;
   return (TRUE);
  else
    SYNTAX ERROR("Package body tail");
                           -- if bypass(token_semicolon)
  end if;
 else
  SYNTAX ERROR("Package body tail");
                           -- if bypass(token begin)
elsif (BYPASS(TOKEN BEGIN)) then
 DECLARATION := FALSE;
 if (SEQUENCE OF STATEMENTS) then
  if (BYPASS(TOKEN EXCEPTION)) then
   if (EXCEPTION HANDLER) then
    while (EXCEPTION HANDLER) loop
      null;
    end loop;
     SYNTAX ERROR("Package body tail");
   end if:
                                  -- if exception handler statement
  end if:
                                  -- if bypass(token exception)
  if (BYPASS(TOKEN END)) then
    HENRY WRITE ENABLE := FALSE;
   if (BYPASS(TOKEN IDENTIFIER)) then
     null:
   end if:
                           -- if bypass(token identifier)
   if (BYPASS(TOKEN SEMICOLON)) then
    DECLARATION := TRUE;
    return (TRUE);
     SYNTAX ERROR("Package body tail");
                           -- if bypass(token semicolon)
   end if;
  else
   SYNTAX ERROR("Package body tail"):
                           -- if bypass(token end)
  end if:
 else
  SYNTAX ERROR("Package body tail");
 end if:
                           -- if sequence of statements
elsif (BYPASS(TOKEN END)) then
  HENRY WRITE ENABLE := FALSE;
 if (BYPASS(TOKEN IDENTIFIER)) then
  null;
                           -- if bypass(token identifier)
 if (BYPASS(TOKEN SEMICOLON)) then
  return (TRUE):
  SYNTAX ERROR("Package body tail"):
 end if:
                           -- if bypass(token semicolon)
else
```

```
return (FALSE);
                           -- if bypass(token separate)
 end if:
end PACKAGE BODY TAIL;
 -- PACKAGE UNIT --> is PACKAGE TAIL END
              --> renames NAME;
function PACKAGE UNIT return boolean is
put(RESULT FILE, "In package unit"); new line(RESULT FILE);
 if (BYPASS(TOKEN IS)) then
  if (PACKAGE TAIL END) then
   return (TRUE);
  else
    SYNTAX ERROR("Package unit");
  end if;
 elsif (BYPASS(TOKEN RENAMES)) then
  if (NAME) then
    if (BYPASS(TOKEN SEMICOLON)) then
     return (TRUE);
    else
     SYNTAX ERROR("Package unit");
                                   -- if bypass(token semicolon)
    end if:
    SYNTAX ERROR("Package unit");
                            -- if name statement
   end if;
  else
   return (FALSE):
                            -- if bypass(token is)
  end if:
end PACKAGE UNIT:
 -- PACKAGE TAIL END --> new NAME [GENERIC ACTUAL PART?];
                  --> BASIC DECLARATIVE ITEM |* private
                         [BASIC DECLARATIVE ITEM]* ?] end [identifier ?] :
 function PACKAGE_TAIL_END return boolean is
 put(RESULT_FILE,"In package tail_end "); new_line(RESULT_FILE);
  if (BYPASS(TOKEN NEW)) then
   if (NAME) then
    if (GENERIC ACTUAL PART) then
      null:
                             -- if generic actual part statement
    if (BYPASS(TOKEN SEMICOLON)) then
      return (TRUE);
    else
      SYNTAX ERROR("Package tail end"):
                             - if bypass(token semicolon)
    end if:
   else
    SYNTAX ERROR("Package tail end"):
```

```
end if:
                         -- if name statement
elsif (BASIC DECLARATIVE ITEM) then
 while (BASIC DECLARATIVE ITEM) loop
  null;
 end loop;
 if (BYPASS(TOKEN PRIVATE)) then
  while (BASIC DECLARATIVE ITEM) loop
   null;
  end loop;
 end if:
                                -- if bypass(token private)
 if (BYPASS(TOKEN END)) then
   HENRY WRITE ENABLE := FALSE;
  if (BYPASS(TOKEN IDENTIFIER)) then
   null:
  end if:
  if (BYPASS(TOKEN SEMICOLON)) then
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, END PACKAGE DECLARE.
             NONE, NEXT HEN);
   CREATE NODE(NEXT HEN, LAST RECORD);
   return (TRUE);
  else
   SYNTAX ERROR("Package tail end");
                          -- if bypass(token semicolon)
  end if:
 else
  SYNTAX ERROR("Package tail end");
                          -- if bypass(token end)
elsif (BYPASS(TOKEN PRIVATE)) then
 while (BASIC DECLARATIVE ITEM) loop
  null:
 end loop:
 if (BYPASS(TOKEN END)) then
   HENRY WRITE ENABLE:= FALSE;
  if (BYPASS(TOKEN IDENTIFIER)) then
    null:
  end if:
  if (BYPASS(TOKEN SEMICOLON)) then
    WRITE HENRY DATA(BLANK, DUMMY LEXEME, END PACKAGE DECLARE.
             NONE, NEXT HEN);
    CREATE NODE(NEXT HEN, LAST RECORD):
    return (TRUE);
  else
    SYNTAX ERROR("Package tail end"):
                          -- if bypass(token semicolon)
  end if;
  SYNTAX ERROR("Package tail end"):
 end if:
                          -- if bypass(token end)
elsif (BYPASS(TOKEN END)) then
  HENRY WRITE ENABLE := FALSE;
 if (BYPASS(TOKEN IDENTIFIER)) then
   null:
 if (BYPASS(TOKEN SEMICOLON)) then
```

```
WRITE HENRY DATA(BLANK, DUMMY LEXEME, END PACKAGE DECLARE,
             NONE, NEXT HEN);
   CREATE NODE(NEXT HEN, LAST RECORD):
   return (TRUE);
   SYNTAX ERROR("Package tail end");
                          -- if bypass(token semicolon)
  end if;
 else
  return (FALSE);
                          -- if bypass(token new)
 end if:
end PACKAGE TAIL END;
 -- BASIC DECLARATIVE ITEM --> BASIC DECLARATIVE
                      --> REPRESENTATION CLAUSE
                      --> use WITH OR USE CLAUSE
function BASIC DECLARATIVE ITEM return boolean is
put(RESULT FILE, "In basic declarative item"); new line(RESULT FILE);
HENRY WRITE ENABLE := TRUE;
 if (BASIC DECLARATION) then
  HENRY WRITE ENABLE := FALSE;
  return (TRUE);
 elsif (REPRESENTATION_CLAUSE) then
  return (TRUE);
 elsif (BYPASS(TOKEN USE)) then
  if (WITH OR USE CLAUSE) then
   return (TRUE);
  else
   SYNTAX ERROR("Basic declarative item");
  end if;
 else
  return (FALSE);
 end if:
end BASIC_DECLARATIVE_ITEM;
 -- DECLARATIVE PART--> BASIC DECLARATIVE ITEM* LATER DECLARATIVE ITEM*
function DECLARATIVE PART return boolean is
put(RESULT_FILE, "In declarative_part"); new line(RESULT_FILE):
 while (BASIC DECLARATIVE ITEM) loop
  null;
 end loop:
 while (LATER DECLARATIVE ITEM) loop
 end loop:
 return (TRUE):
end DECLARATIVE PART:
```

```
-- BASIC DECLARATION --> type TYPE DECLARATION
                 --> subtype SUBTYPE DECLARATION
                 --> procedure PROCEDURE UNIT
                 --> function FUNCTION UNIT
                 --> package PACKAGE DECLARATION
                 --> generic GENERIC DECLARATION
                 --> IDENTIFIER DECLARATION
                 --> task TASK DECLARATION
function BASIC DECLARATION return boolean is
put(RESULT FILE. "In basic declaration"); new line(RESULT FILE);
 if (BYPASS(TOKEN TYPE)) then
  if (TYPE DECLARATION) then
   return (TRUE);
  elsc
   SYNTAX ERROR("Basic declaration");
 elsif (BYPASS(TOKEN SUBTYPE)) then
  if (SUBTYPE DECLARATION) then
   return (TRUE):
  else
   SYNTAX ERROR("Basic declaration");
  end if:
 elsif (BYPASS(TOKEN PROCEDURE)) then
  DECLARE TYPE := PROCEDURE DECLARE:
  if (PROCEDURE UNIT) then
  HENRY WRITE ENABLE : = FALSE;
   return (TRUE):
  else
   SYNTAX ERROR("Basic declaration");
                                 -- if procedure unit statement
 elsif (BYPASS(TOKEN FUNCTION)) then
  DECLARE TYPE := FUNCTION DECLARE;
  if (FUNCTION UNIT) then
  HENRY WRITE ENABLE := FALSE;
   return (TRUE);
  else
   SYNTAX ERROR("Basic declaration");
                                 -- if function unit statement
  end if;
 elsif (BYPASS(TOKEN PACKAGE)) then
  DECLARE TYPE := PACKAGE DECLARE:
  if (PACKAGE DECLARATION) then
   return (TRUE):
  else
   SYNTAX ERROR("Basic declaration");
                                 -- if package declaration
 elsif (BYPASS(TOKEN GENERIC)) then
  if (GENERIC DECLARATION) then
   return (TRUE);
  else
```

```
SYNTAX ERROR("Basic declaration"):
  end if:
                                 -- if generic declaration
 elsif (IDENTIFIER DECLARATION) then
  HENRY WRITE ENABLE := FALSE;
  return (TRUE);
 elsif (BYPASS(TOKEN TASK)) then
  DECLARE TYPE := TASK DECLARE;
  if (TASK DECLARATION) then
   return (TRUE);
  else
   SYNTAX ERROR("Basic declaration");
  end if:
 else
  return (FALSE);
 end if:
end BASIC DECLARATION;
      .....
 -- LATER DECLARATIVE ITEM --> PROPER BODY
                      --> generic GENERIC DECLARATION
                      -- > use WITH OR USE CLAUSE
function LATER DECLARATIVE ITEM return boolean is
put(RESULT FILE, "In later declarative item "); new line(RESULT FILE);
                                       -- check for body declaration
if (PROPER BODY) then
  return (TRUE);
 elsif (BYPASS(TOKEN GENERIC)) then
  if (GENERIC DECLARATION) then
   return (TRUE);
   SYNTAX ERROR("Later declarative item"):
                                 -- if generic declaration
 elsif (BYPASS(TOKEN USE)) then
  if (WITH OR USE CLAUSE) then
   return (TRUE);
  else
   SYNTAX ERROR("Later declarative item"):
                                -- if with or use clause
  end if:
 else
  return (FALSE);
end LATER DECLARATIVE ITEM:
 -- PROPER BODY --> procedure PROCEDURE UNIT
             -- · function FUNCTION UNIT
             -- package PACKAGE DECLARATION
             -- task TASK DECLARATION
function PROPER BODY return boolean is
begin
```

```
put(RESULT_FILE, "In proper_body"); new_line(RESULT_FILE);
 if (BYPASS(TOKEN PROCEDURE)) then
  DECLARE TYPE: PROCEDURE DECLARE:
  if (PROCEDURE UNIT) then
   return (TRUE):
  else
   SYNTAX ERROR("Proper body"):
  end if:
                                  -- if procedure unit statement
 elsif (BYPASS(TOKEN FUNCTION)) then
  DECLARE TYPE: FUNCTION DECLARE:
  if (FUNCTION UNIT) then
   return (TRUE):
  else
   SYNTAX ERROR("Proper body");
  end if:
                                  -- if function unit statement
 elsif (BYPASS(TOKEN PACKAGE)) then
  DECLARE TYPE := PACKAGE DECLARE:
  if (PACKAGE DECLARATION) then
   return (TRUE):
  else
   SYNTAX ERROR("Proper body");
                                  -- if package declaration
 elsif (BYPASS(TOKEN TASK)) then
  DECLARE TYPE := TASK DECLARE;
  if (TASK DECLARATION) then
   return (TRUE):
  else
   SYNTAX ERROR("Proper body");
  end if:
 else
  return (FALSE);
 end if:
                          -- if bypass(token procedure)
end PROPER BODY:
 -- SEQUENCE OF STATEMENTS --> STATEMENT -
function SEQUENCE OF STATEMENTS return boolean is
begin
put(RESULT FILE, "In sequence of statements"), new line(RESULT FILE);
 if (STATEMENT) then
  while (STATEMENT) loop
   null:
  end loop:
  return (TRUE):
 else
  return (FALSE);
 end if:
end SEQUENCE OF STATEMENTS.
```

```
-- STATEMENT -- LABEL ? SIMPLE STATEMENT
          --> LABEL? COMPOUND STATEMENT
function STATEMENT return boolean is
put(RESULT FILE. "In statement "); new line(RESULT FILE);
 if (LABEL) then
  null:
 end if:
 if (SIMPLE STATEMENT) then
  return (TRUE).
 elsif (COMPOUND STATEMENT) then
  return (TRUE):
  return (FALSE):
 end if:
end STATEMENT:
  .....
 -- COMPOUND STATEMENT --> if IF STATEMENT
                   -- case CASE STATEMENT
                   --> LOOP STATEMENT
                   -- BLOCK STATEMENT
                   -- · accept ACCEPT STATEMENT
                   -- select SELECT STATEMENT
function COMPOUND STATEMENT return boolean is
put(RESULT_FILE, "In compound_statement"); new_line(RESULT_FILE);
 if (BYPASS(TOKEN IF)) then
  NESTING METRIC(IF CONSTRUCT):
  if (IF STATEMENT) then
   return (TRUE):
   SYNTAX ERROR("Compound statement");
                                -- if if statement
  end if:
 elsif (BYPASS(TOKEN CASE)) then
  NESTING METRIC(CASE CONSTRUCT);
  if (CASE STATEMENT) then
   return (TRUE):
   SYNTAX ERROR("Compound statement");
                                -- if case statement
 elsif (LOOP STATEMENT) then
  return (TRUE)
 elsif (BLOCK STATEMENT) then
  return (TRUE):
 elsif (BYPASS(TOKEN ACCEPT)) then
  if (ACCEPT STATEMENT) then
   return (TRUE):
   SYNTAX ERROR("Compound statement");
  end if:
```

```
elsif (BYPASS(TOKEN SELECT)) then
  if (SELECT STATEMENT) then
   return (TRUE):
  else
   SYNTAX ERROR("Compound statement"):
  end if:
 else
  return (FALSE):
 end if:
end COMPOUND STATEMENT:
 -- BLOCK STATEMENT -- identifier : ? declare DECLARATIVE PART ?
                      begin SEQUENCE OF STATEMENTS exception
                       EXCEPTION HANDLER + ? ? end identifier? :
function BLOCK STATEMENT return boolean is
 DECLARE STATUS: boolean;
put(RESULT FILE. "In block statement"); new line(RESULT FILE);
 if (DECLARATION) then
  DECLARE STATUS := TRUE;
 else
  DECLARATION : = TRUE:
  DECLARE STATUS: FALSE:
 end if:
 DECLARE TYPE := BLOCK DECLARE;
 if (BYPASS(TOKEN IDENTIFIER)) then
  SCOPE LEVEL := SCOPE LEVEL + 1;
  if (BYPASS(TOKEN COLON)) then
   SCOPE LEVEL := SCOPE LEVEL - 1;
  else
   SYNTAX ERROR("Block statement");
                                 -- if bypass(token_colon)
  end if:
 else
  DECLARE TYPE := VARIABLE DECLARE;
                           -- if bypass(token identifier)
 end if:
 if (BYPASS(TOKEN DECLARE)) then
  SCOPE LEVEL := SCOPE LEVEL + 1;
  if (DECLARATIVE PART) then
   null:
   SYNTAX ERROR("Block statement");
                                 -- if declarative part statement
  end if:
 end if:
                           -- if bypass(token declare)
 if (BYPASS(TOKEN BEGIN)) then
  DECLARATION : = FALSE;
  if (SEQUENCE OF STATEMENTS) then
   if (BYPASS(TOKEN EXCEPTION)) then
    if (EXCEPTION HANDLER) then
      while (EXCEPTION HANDLER) loop
       null:
```

```
end loop.
    else
     SYNTAX ERROR("Block statement");
                                   -- if exception handler statement
    end if.
                                  -- if bypass(token exception)
   end if
   if (BYPASS(TOKEN_END)) then
    if (BYPASS(TOKEN IDENTIFIER)) then
      null:
    end if.
                                   -- if bypass(token identifier)
    if (BYPASS(TOKEN SEMICOLON)) then
      SCOPE LEVEL: SCOPE LEVEL - 1:
     DECLARATION: TRUE:
      return (TRUE).
    else
      SYNTAX ERROR("Block statement");
                                   -- if bypass(token semicolon)
   else
    SYNTAX ERROR("Block statement");
   end if
                                   -- if bypass(token end)
  else
   SYNTAX ERROR("Block statement"):
  end if:
                                   -- if sequence of statements
 else
  if not (DECLARE STATUS) then
   DECLARATION FALSE:
  end if:
  return (FALSE):
 end if:
                           -- if bypass(token begin)
end BLOCK STATEMENT.
 -- IF STATEMENT -- EXPRESSION then SEQUENCE OF STATEMENTS
                     elsif EXPRESSION then SEQUENCE OF STATEMENTS *
                     else SEQUENCE OF STATEMENTS? end if :
function IF STATEMENT return boolean is
put(RESULT FILE, "In if statement"); new line(RESULT FILE);
 if (EXPRESSION) then
  if (BYPASS(TOKEN THEN)) then
   if (SEQUENCE OF STATEMENTS) then
    while (BYPASS(TOKEN ELSIF)) loop
      if (EXPRESSION) then
       if (BYPASS(TOKEN THEN)) then
        if not (SEQUENCE OF STATEMENTS) then
         SYNTAX ERROR("If statement"):
                           -- if not sequence of statements
       else
        SYNTAX ERROR("If statement");
                            -- if bypass(token then)
       end if:
       SYNTAX ERROR("If statement");
```

```
end if.
                            -- if expression statement
     end loop:
     if (BYPASS(TOKEN ELSE)) then
      if (SEQUENCE OF STATEMENTS) then
       null:
      PISP
       SYNTAX ERROR("If statement"):
                            -- if sequence of statements
     end if:
                            -- if bypass(token else)
     if (BYPASS(TOKEN END)) then
      if (BYPASS(TOKEN IF)) then
       if (BYPASS(TOKEN SEMICOLON)) then
         NESTING METRIC(IF END):
        return (TRUE):
       else
        SYNTAX ERROR("If statement");
       end if:
                            -- if bypass(token semicolon)
       SYNTAX ERROR("If statement"):
      end if.
                            -- if bypass(token if)
      SYNTAX ERROR("If statement");
                            -- if bypass(token end)
     SYNTAX ERROR("If statement");
   end if:
                            -- if sequence of statements
    SYNTAX ERROR("If statement"):
  end if.
                            -- if bypass(token then)
 else
  return (FALSE).
 end if
                            -- if expression statement
end IF STATEMENT.
-- CASE STATEMENT -- EXPRESSION is CASE STATEMENT ALTERNATIVE - end case.
function CASE STATEMENT return boolean is
put(RESULT_FILE, "In case statement"), new line(RESULT_FILE);
 if (EXPRESSION) then
  if (BYPASS(TOKEN IS)) then
   if (CASE STATEMENT ALTERNATIVE) then
     while (CASE STATEMENT ALTERNATIVE) loop
      null
    end loop.
    if (BYPASS(TOKEN END)) then
      if (BYPASS(TOKEN CASE)) then
       if (BYPASS(TOKEN SEMICOLON)) then
        NESTING METRIC(CASE END).
        return (TRUE):
       0 50
```

```
SYNTAX ERROR("Case statement"):
                            -- if bypass(token semicolon)
       end if:
      else
       SYNTAX ERROR("Case statement"):
                            -- if bypass(token case)
    else
      SYNTAX ERROR("Case statement");
    end if:
                            -- if bypass(token end)
    SYNTAX ERROR("Case statement");
   end if:
                            -- if case statement alternative
  else
   SYNTAX ERROR("Case statement");
                            -- if bypass(token is)
  end if;
 else
  return (FALSE);
 end if;
                            -- if expression statement
end CASE STATEMENT;
 -- CASE STATEMENT ALTERNATIVE --> when CHOICE ( CHOICE * =>
                                    SEQUENCE OF STATEMENTS
function CASE STATEMENT ALTERNATIVE return boolean is
put(RESULT FILE, "In case statement alternative"); new line(RESULT FILE);
 if (BYPASS(TOKEN WHEN)) then
  if (CHOICE) then
   while (BYPASS(TOKEN BAR)) loop
     if not (CHOICE) then
      SYNTAX ERROR("Case statement alternative");
                             -- if not choice statement
   end loop;
    if (BYPASS(TOKEN ARROW)) then
     if (SEQUENCE OF STATEMENTS) then
      return (TRUE):
     else
      SYNTAX ERROR("Case statement alternative");
                             -- if sequence of statements
     SYNTAX ERROR("Case statement alternative"):
                             -- if bypass(token arrow)
   end if:
  else
    SYNTAX ERROR("Case statement alternative"):
  end if:
                             -- if choice statement
 else
  return (FALSE):
                             -- if bypass(token when)
end CASE STATEMENT ALTERNATIVE:
```

```
-- LOOP STATEMENT --> identifier: ? ITERATION SCHEME? loop
                      SEQUENCE OF STATEMENTS end loop identifier ?!;
function LOOP STATEMENT return boolean is
put(RESULT FILE, "In loop statement"); new line(RESULT FILE);
 if (BYPASS(TOKEN IDENTIFIER)) then
  if (BYPASS(TOKEN COLON)) then
   null:
  else
   SYNTAX ERROR("Loop statement");
  end if;
                           -- if bypass(token colon)
 end if;
                           -- if bypass(token identifier)
 if (ITERATION SCHEME) then
  NO ITERATION := FALSE;
                            -- if iteration scheme statement
 if (BYPASS(TOKEN LOOP)) then
  if (NO ITERATION) then
   NESTING METRIC(LOOP CONSTRUCT);
   NO ITERATION := TRUE;
  end if:
  if (SEQUENCE OF STATEMENTS) then
   if (BYPASS(TOKEN END)) then
    if (BYPASS(TOKEN LOOP)) then
      if (BYPASS(TOKEN IDENTIFIER)) then
       null:
      end if:
                           -- if bypass(token identifier)
      if (BYPASS(TOKEN SEMICOLON)) then
       NESTING METRIC(LOOP END):
       return (TRUE);
      else
       SYNTAX ERROR("Loop statement");
      end if:
                           -- if bypass(token semicolon)
     else
      SYNTAX ERROR("Loop statement"):
                           -- if bypass(token loop)
   else
    SYNTAX ERROR("Loop statement");
   end if:
                           -- if bypass(token end)
   SYNTAX ERROR("Loop statement");
  end if:
                           -- if sequence of statements
 else
  return (FALSE);
 end if:
                           -- if bypass(token loop)
end LOOP STATEMENT;
 -- EXCEPTION HANDLER -- · when EXCEPTION CHOICE - EXCEPTION CHOICE *
                        SEQUENCE OF STATEMENTS
function EXCEPTION HANDLER return boolean is
```

```
begin
put(RESULT FILE, "In exception handler"); new line(RESULT FILE);
 if (BYPASS(TOKEN WHEN)) then
  if (EXCEPTION CHOICE) then
   while (BYPASS(TOKEN BAR)) loop
    if not (EXCEPTION CHOICE) then
      SYNTAX ERROR("Exception handler");
    end if:
                                   -- if not exception choice
   end loop;
   if (BYPASS(TOKEN ARROW)) then
    if (SEQUENCE OF STATEMENTS) then
      return (TRUE);
    else
      SYNTAX ERROR("Exception handler");
                                   -- if sequence of statements
   else
     SYNTAX ERROR("Exception handler");
                                   -- if bypass(token arrow)
   SYNTAX ERROR("Exception handler");
  end if:
                                   -- if exception choice statement
 else
  return (FALSE);
                            -- if bypass(token-when)
 end if:
end EXCEPTION HANDLER;
 -- ACCEPT STATEMENT --> identifier (EXPRESSION) ? FORMAL PART ?
                         do SEQUENCE OF STATEMENTS end identifier??;
function ACCEPT STATEMENT return boolean is
begin
put(RESULT FILE, "In accept statement"); new line(RESULT FILE);
 if (BYPASS(TOKEN IDENTIFIER)) then
  if (BYPASS(TOKEN LEFT PAREN)) then
    if (EXPRESSION) then
     if (BYPASS(TOKEN RIGHT PAREN)) then
      null;
     else
      SYNTAX ERROR("Accept statement"):
     end if;
                                   -- if bypass(token right paren)
    else
     SYNTAX ERROR("Accept statement"):
    end if:
                                   -- if expression statement
  end if:
                                   -- if bypass(token left paren)
  if (FORMAL PART) then
    null:
   end if.
                                   -- if formal part statement
   if (BYPASS(TOKEN DO)) then
    if (SEQUENCE OF STATEMENTS) then
     if (BYPASS(TOKEN END)) then
      if (BYPASS(TOKEN IDENTIFIER)) then
```

```
null:
     end if:
                            -- if bypass(token identifier)
    else
     SYNTAX ERROR("Accept statement");
                            -- if bypass(token end)
     SYNTAX ERROR("Accept statement"):
   end if:
                                   -- if sequence of statements
  end if:
                                   -- if bypass(token do)
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
  else
   SYNTAX ERROR("Accept statement"):
  end if:
                                   -- if bypass(token semicolon)
 else
  return (FALSE):
 end if:
                            -- if bypass(token identifier)
end ACCEPT STATEMENT:
 -- SELECT STATEMENT --> SELECT STATEMENT TAIL SELECT ENTRY CALL end select :
function SELECT STATEMENT return boolean is
begin
put(RESULT FILE, "In select statement"); new line(RESULT FILE);
 if (SELECT STATEMENT TAIL) then
  if (SELECT ENTRY CALL) then
   if (BYPASS(TOKEN END)) then
     if (BYPASS(TOKEN SELECT)) then
      if (BYPASS(TOKEN SEMICOLON)) then
       return (TRUE);
      else
       SYNTAX ERROR("Select statement");
                            -- if bypass(token semicolon)
     else
      SYNTAX ERROR("Select statement");
     end if:
                                   -- if bypass(token select)
     SYNTAX ERROR("Select statement");
   end if:
                                   -- if bypass(token end)
   SYNTAX ERROR("Select statement");
  end if:
                                   -- if select entry call statement
 else
  return (FALSE):
                            -- if select statement tail
 end if:
end SELECT STATEMENT:
 -- SELECT STATEMENT TAIL --> SELECT ALTERNATIVE or SELECT ALTERNATIVE *
                       --- NAME: SEQUENCE OF STATEMENTS?
```

Contraction (Contractor)

```
function SELECT STATEMENT TAIL return boolean is
put(RESULT FILE, "In select statement tail"); new line(RESULT FILE);
 if (SELECT ALTERNATIVE) then
  while (BYPASS(TOKEN OR)) loop
   if not (SELECT ALTERNATIVE) then
    SYNTAX ERROR("Select statement tail");
   end if:
  end loop:
  return (TRUE);
 elsif (NAME) then
                                   -- check for entry call statement
  if (BYPASS(TOKEN SEMICOLON)) then
   if (SEQUENCE OF STATEMENTS) then
   end if:
                                   -- if sequence of statements
   return (TRUE);
  else
   SYNTAX ERROR("Select statement tail");
                                   -- if bypass(token semicolon)
 else
  return (FALSE);
                            -- if select alternative statement
 end if:
end SELECT STATEMENT TAIL;
 -- SELECT ALTERNATIVE --> [when EXPRESSION => ?] accept ACCEPT STATEMENT
                            SEQUENCE OF STATEMENTS?
                     --> when EXPRESSION => ?' delay DELAY STATEMENT
                            SEQUENCE OF STATEMENTS?
                     --> when EXPRESSION => ? terminate;
function SELECT ALTERNATIVE return boolean is
begin
put(RESULT FILE, "In select alternative"); new line(RESULT FILE);
 if (BYPASS(TOKEN WHEN)) then
  if (EXPRESSION) then
    if (BYPASS(TOKEN ARROW)) then
     null:
    else
     SYNTAX ERROR("Select alternative");
                                   -- if bypass(token arrow)
  else
    SYNTAX ERROR("Select alternative");
  end if;
                                   -- if expression statement
 end if;
                            -- if bypass(token when)
 if (BYPASS(TOKEN ACCEPT)) then
  if (ACCEPT STATEMENT) then
    if (SEQUENCE OF STATEMENTS) then
     null:
    end if:
                                   -- if sequence of statements
    return (TRUE);
   else
```

```
SYNTAX ERROR("Select alternative");
                                   -- if accept statement
  end if;
 elsif (BYPASS(TOKEN DELAY)) then
  if (DELAY STATEMENT) then
   if (SEQUENCE OF STATEMENTS) then
    null:
   end if:
                                   -- if sequence of statements
   return (TRUE):
   SYNTAX ERROR("Select alternative");
                                   -- if delay statement
 elsif (BYPASS(TOKEN TERMINATE)) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
   SYNTAX ERROR("Select alternative"):
  end if;
                                   -- if bypass(token semicolon)
 else
  return (FALSE):
 end if:
                            -- if bypass(token accept)
end SELECT ALTERNATIVE:
 -- SELECT ENTRY CALL --> else SEQUENCE OF STATEMENTS
                  --> or delay DELAY STATEMENT SEQUENCE OF STATEMENTS?
function SELECT ENTRY CALL return boolean is
begin
put(RESULT_FILE. "In select_entry_call"); new_line(RESULT_FILE):
 if (BYPASS(TOKEN ELSE)) then
  if (SEQUENCE OF STATEMENTS) then
   return (TRUE):
  else
   SYNTAX ERROR("Select entry call");
                                   -- if sequence of statements
 elsif (BYPASS(TOKEN OR)) then
  if (BYPASS(TOKEN DELAY)) then
   if (DELAY STATEMENT) then
     if (SEQUENCE OF STATEMENTS) then
     end if:
                                   -- if sequence of statements
     return (TRUE):
     SYNTAX ERROR("Select entry call"):
                                   -- if delay statement
   SYNTAX ERROR("Select entry call"):
                                    -- if bypass(token delay)
  end if;
 else
  return (FALSE):
 end if;
                            -- if bypass(token else)
end SELECT ENTRY CALL:
```

```
AN ADA SOFTWARE METRIC
-- TITLE:
-- MODULE NAME: PACKAGE PARSER 2
-- DATE CREATED: 18 JUL 86
-- LAST MODIFIED: 30 MAY 87
                LCDR JEFFREY L. NIEDER
 AUTHORS:
           LT KARL S. FAIRBANKS, JR.
           LCDR PAUL M. HERZIG
-- DESCRIPTION: This package contains thirty-three functions
     that are the middle level productions for our top-down.
     recursive descent parser. Each function is preceded
     by the grammaar productions they are implementing.
with PARSER 3. PARSER 4. HENRY GLOBAL, HENRY, BYPASS FUNCTION,
   BYPASS SUPPORT FUNCTIONS, GLOBAL PARSER, GLOBAL, TEXT 10;
use PARSER 3. PARSER 4. HENRY GLOBAL, HENRY. BYPASS FUNCTION,
  BYPASS SUPPORT FUNCTIONS, GLOBAL PARSER, GLOBAL, TEXT 10;
package PARSER 2 is
 IDENT DECLARE: BOOLEAN:= FALSE:
 function GENERIC ACTUAL PART return boolean;
 function GENERIC ASSOCIATION return boolean;
 function GENERIC FORMAL PARAMETER return boolean;
 function GENERIC TYPE DEFINITION return boolean:
 function PRIVATE TYPE DECLARATION return boolean;
 function TYPE DECLARATION return boolean;
 function SUBTYPE DECLARATION return boolean;
 function DISCRIMINANT PART return boolean;
 function DISCRIMINANT SPECIFICATION return boolean;
 function TYPE DEFINITION return boolean;
 function RECORD TYPE DEFINITION return boolean;
 function COMPONENT LIST return boolean;
 function COMPONENT DECLARATION return boolean:
 function VARIANT PART return boolean;
 function VARIANT return boolean:
 function WITH OR USE CLAUSE return boolean:
 function FORMAL PART return boolean:
 function IDENTIFIER DECLARATION return boolean:
 function IDENTIFIER DECLARATION TAIL return boolean:
 function EXCEPTION TAIL return boolean:
 function EXCEPTION CHOICE return boolean:
 function CONSTANT TERM return boolean;
```

```
function IDENTIFIER TAIL return boolean:
 function PARAMETER SPECIFICATION return boolean:
 function IDENTIFIER LIST return boolean:
 function MODE return boolean:
 function DESIGNATOR return boolean;
 function SIMPLE STATEMENT return boolean;
 function ASSIGNMENT OR PROCEDURE CALL return boolean;
 function LABEL return boolean:
 function ENTRY DECLARATION return boolean;
 function REPRESENTATION CLAUSE return boolean:
 function RECORD REPRESENTATION CLAUSE return boolean;
end PARSER 2:
package body PARSER 2 is
 -- GENERIC ACTUAL PART --> (GENERIC ASSOCIATION , GENERIC ASSOCIATION )
function GENERIC ACTUAL PART return boolean is
begin
 if (BYPASS(TOKEN LEFT PAREN)) then
  if (GENERIC ASSOCIATION) then
    while (BYPASS(TOKEN COMMA)) loop
     if not (GENERIC ASSOCIATION) then
      SYNTAX ERROR("Generic actual part"):
    end if:
                                   -- if not generic association
    end loop:
   if (BYPASS(TOKEN RIGHT PAREN)) then
    return (TRUE):
    SYNTAX ERROR("Generic actual part"):
                                  -- if bypass(token right paren)
   end if:
  else
   SYNTAX ERROR("Generic actual part");
  end if:
                                  -- if generic association statement
 else
  return(FALSE);
 end if:
                           -- if bypass(token left paren)
end GENERIC ACTUAL PART:
 -- GENERIC ASSOCIATION -- - GENERIC FORMAL PARAMETER? EXPRESSION
function GENERIC ASSOCIATION return boolean is
 if (GENERIC FORMAL PARAMETER) then
  null:
 end if:
                           -- if generic formal parameter statement
 if (EXPRESSION) then
                                         -- check for generic actual parameter
  return (TRUE):
 else
```

```
return (FALSE);
                           -- if expression
 end if:
end GENERIC ASSOCIATION:
 -- GENERIC FORMAL PARAMETER --> identifier =>
                        --> string literal =>
function GENERIC FORMAL PARAMETER return boolean is
 LOOK AHEAD TOKEN := TOKEN RECORD BUFFER(TOKEN ARRAY INDEX + 1):
 if (ADJUST LEXEME(LOOK AHEAD TOKEN.LEXEME,
                LOOK AHEAD TOKEN.LEXEME SIZE - 1) = "= >") then
  if (BYPASS(TOKEN IDENTIFIER)) then
   if (BYPASS(TOKEN ARROW)) then
    return (TRUE):
   else
    SYNTAX ERROR("Generic formal parameter"):
                                  -- if bypass(token arrow)
  elsif (BYPASS(TOKEN STRING LITERAL)) then
   if (BYPASS(TOKEN ARROW)) then
    return (TRUE):
   else
    SYNTAX ERROR("Generic formal parameter");
   end if:
                                  -- if bypass(token arrow)
   SYNTAX ERROR("Generic formal parameter");
  end if;
                                  -- if bypass(token identifier)
 else
  return (FALSE);
                           -- if adjust lexeme(lookahead token)
end GENERIC FORMAL PARAMETER;
 -- GENERIC TYPE DEFINITION --> ( <> )
                       --> range <>
                       --> digits <>
                       --> delta <>
                       --> array ARRAY TYPE DEFINITION
                       --> access SUBTYPE INDICATION
function GENERIC TYPE DEFINITION return boolean is
begin
 if (BYPASS(TOKEN LEFT PAREN)) then
  if (BYPASS(TOKEN BRACKETS)) then
   if (BYPASS(TOKEN RIGHT PAREN)) then
    return (TRUE);
    SYNTAX ERROR("Generic type definition"):
   end if:
                                  -- if bypass(token right paren)
   SYNTAX ERROR("Generic type definition"):
```

```
end if:
                                  -- if bypass(token brackets)
 elsif (BYPASS(TOKEN RANGE)) or else (BYPASS(TOKEN DIGITS))
  or else (BYPASS(TOKEN DELTA)) then
  if (BYPASS(TOKEN BRACKETS)) then
   return (TRUE):
  else
   SYNTAX ERROR("Generic type definition");
                                   - if bypass(token brackets)
 elsif (BYPASS(TOKEN ARRAY)) then
  if (ARRAY TYPE DEFINITION) then
   return (TRUE):
   SYNTAX ERROR("Generic type definition");
                                   -- if array type definition
 elsif (BYPASS(TOKEN ACCESS)) then
  if (SUBTYPE INDICATION) then
   return (TRUE);
  else
   SYNTAX ERROR("Generic type definition");
  end if:
                                   -- if subtype indication
 else
  return (FALSE):
                           -- if bypass(token left paren)
end GENERIC TYPE DEFINITION:
 -- PRIVATE TYPE DECLARATION --> limited? private
function PRIVATE TYPE DECLARATION return boolean is
 if (BYPASS(TOKEN LIMITED)) then
  null:
 end if:
 if (BYPASS(TOKEN PRIVATE)) then
  return (TRUE);
 else
  return (FALSE):
 end if:
end PRIVATE TYPE DECLARATION:
 -- SUBTYPE DECLARATION --> identifier is SUBTYPE INDICATION :
function SUBTYPE DECLARATION return boolean is
begin
 if (BYPASS(TOKEN IDENTIFIER)) then
  if (BYPASS(TOKEN IS)) then
   if (SUBTYPE INDICATION) then
    if (BYPASS(TOKEN SEMICOLON)) then
     return (TRUE):
     SYNTAX ERROR("Subtype declaration");
```

```
end if:
                                   -- if bypass(token semicolon)
   else
    SYNTAX ERROR("Subtype declaration"):
   end if;
                            -- if subtype indication statement
   SYNTAX ERROR("Subtype declaration"):
  end if:
                            -- if bypass(token is)
 else
  return (FALSE);
 end if:
                            -- if bypass(token identifier)
end SUBTYPE DECLARATION:
 -- TYPE DECLARATION --> identifier DISCRIMINANT PART?
                        is SUBTYPE INDICATION:
function TYPE DECLARATION return boolean is
 if (BYPASS(TOKEN IDENTIFIER)) then
  if (DISCRIMINANT PART) then
   null;
  end if:
                            -- if discriminant part statement
  if (BYPASS(TOKEN IS)) then
                                   -- declaration is full type if 'is'
   if (PRIVATE TYPE DECLARATION) then
   elsif (TYPE DEFINITION) then -- present, otherwise incomplete type
     null:
   else
     SYNTAX ERROR("Type declaration"):
   end if:
                            -- if type definition statement
  end if:
                            -- if bypass(token is)
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
   SYNTAX ERROR("Type declaration"):
  end if:
                            -- if bypass(token semicolon)
 else
  return (FALSE);
                            -- if bypass(token identifier)
 end if:
end TYPE DECLARATION:
 -- DISCRIMINANT PART --> (DISCRIMINANT SPECIFICATION
                         : DISCRIMINANT SPECIFICATION * )
function DISCRIMINANT PART return boolean is
begin
 if (BYPASS(TOKEN LEFT PAREN)) then
  if (DISCRIMINANT SPECIFICATION) then
   while (BYPASS(TOKEN SEMICOLON)) loop
     if not (DISCRIMINANT SPECIFICATION) then
      SYNTAX ERROR("Discriminant part"):
```

```
end if;
                       -- if not discriminant specification
   end loop:
   if (BYPASS(TOKEN RIGHT PAREN)) then
    return (TRUE);
   else
     SYNTAX ERROR("Discriminant part");
                       -- if bypass(token right paren)
   end if:
  else
   SYNTAX ERROR("Discriminant part"):
                       -- if discriminant specification
 else
  return (FALSE);
 end if:
                       -- if bypass(token left paren)
end DISCRIMINANT PART;
 -- DISCRIMINANT SPECIFICATION --> IDENTIFIER LIST: NAME := EXPRESSION?
function DISCRIMINANT SPECIFICATION return boolean is
 if (IDENTIFIER LIST) then
  if (BYPASS(TOKEN COLON)) then
   if (NAME) then
                                   -- check for type mark
     if (BYPASS(TOKEN ASSIGNMENT)) then
      if (EXPRESSION) then
       null;
      else
       SYNTAX ERROR("Discriminant specification");
      end if:
                            -- if expression statement
     end if;
                            -- if bypass(token assignment)
     return (TRUE);
   else
     SYNTAX ERROR("Discriminant specification");
   end if:
                            -- if name statement
  else
   SYNTAX ERROR("Discriminant specification");
  end if:
                            -- if bypass(token colon)
 else
  return (FALSE);
                            -- if identifier list statement
end DISCRIMINANT SPECIFICATION;
 -- TYPE DEFINITION --> ENUMERATION TYPE DEFINITION
                 --> INTEGER TYPE DEFINITION
                 --> digits FLOATING OR FIXED POINT CONSTRAINT
                 --> delta FLOATING OR FIXED POINT CONSTRAINT
                 -- array ARRAY TYPE DEFINITION
                 --> record RECORD TYPE DEFINITION
                 -- - access SUBTYPE INDICATION
                 -- new SUBTYPE INDICATION
```

```
function TYPE DEFINITION return boolean is
begin
 if (ENUMERATION TYPE DEFINITION) then
  return (TRUE):
 elsif (INTEGER TYPE DEFINITION) then
  return (TRUE):
 elsif (BYPASS(TOKEN DIGITS)) or else (BYPASS(TOKEN DELTA)) then
  if (FLOATING OR FIXED POINT CONSTRAINT) then
   return (TRUE):
  else
   SYNTAX ERROR("Type definition"):
                                   -- floating or fixed point constraint
  end if;
 elsif (BYPASS(TOKEN ARRAY)) then
  if (ARRAY TYPE DEFINITION) then
   return (TRUE):
  else
   SYNTAX ERROR("Type definition");
                                   -- if array type definition
 elsif (BYPASS(TOKEN RECORD STRUCTURE)) then
  if (RECORD TYPE LEFINITION) then
   return (TRUE):
  else
   SYNTAX ERROR("Type definition"):
  end if;
                                   -- if record type definition
 elsif (BYPASS(TOKEN ACCESS)) or else (BYPASS(TOKEN NEW)) then
  if (SUBTYPE INDICATION) then
   return (TRUE):
  else
   SYNTAX ERROR("Type definition");
  end if:
                                   -- if subtype indication
 else
  return (FALSE);
 end if;
end TYPE DEFINITION:
 -- RECORD TYPE DEFINITION --> COMPONENT LIST end record
function RECORD TYPE DEFINITION return boolean is
begin
 if (COMPONENT LIST) then
  if (BYPASS(TOKEN END)) then
   if (BYPASS(TOKEN RECORD STRUCTURE)) then
    return (TRUE):
    SYNTAX ERROR("Record type definition");
   end if;
                                   -- if bypass(token record-structure)
  else
   SYNTAX ERROR("Record type definition");
  end if:
                                   -- if bypass(token end)
 else
  return (FALSE);
```

```
- if component list statement
end if.
and RECORD TYPE DEFINITION.
      ____
 .. COMPONENT LIST ... COMPONENT DECLARATION * VAARIANT PART ?
               -- · null:
function COMPONENT LIST return boolean is
 while (COMPONENT DECLARATION) loop
  null.
 end loop.
 if (VARIANT PART) then
  null.
 elsif (BYPASS(TOKEN NULL)) then
  if (BYPASS(TOKEN SEMICOLON)) then
   null
  end if
 end if:
 return (TRUE).
and COMPONENT LIST.
 -- COMPONENT DECLARATION -- IDENTIFIER LIST SUBTYPE INDICATION
                               EXPRESSION?
function COMPONENT DECLARATION return boolean is
begin
 if (IDENTIFIER LIST) then
  if (BYPASS(TOKEN COLON)) then
   if (SUBTYPE INDICATION) then
    if (BYPASS(TOKEN ASSIGNMENT)) then
      if (EXPRESSION) then
       if (BYPASS(TOKEN SEMICOLON)) then
        return (TRUE):
       else
        SYNTAX ERROR("Component declaration").
                          -- if bypass(token semicolon)
       end if:
       SYNTAX ERROR("Component declaration");
      end if:
                          -- if expression statement
     end if:
                                 -- if bypass(token assignment)
     if (BYPASS(TOKEN SEMICOLON)) then
      return (TRUE);
     else
      SYNTAX ERROR("Component declaration");
     end if:
                                 -- if bypass(token semicolon)
     SYNTAX ERROR("Component declaration"):
                                 -- if subtype indication statement
   end if:
   SYNTAX ERROR("Component declaration"):
```

```
end if:
                                    -- if bypass(token colon)
else
  return (FALSE):
end if:
                            -- if identifier list statement
end COMPONENT DECLARATION:
 -- VARIANT PART -- > case identifier is VARIANT + end case :
function VARIANT PART return boolean is
 if (BYPASS(TOKEN CASE)) then
  if (BYPASS(TOKEN IDENTIFIER)) then
   if (BYPASS(TOKEN IS)) then
    if (VARIANT) then
     while (VARIANT) loop
       null:
      end loop:
      if (BYPASS(TOKEN END)) then
       if (BYPASS(TOKEN CASE)) then
        if (BYPASS(TOKEN SEMICOLON)) then
         return (TRUE);
        else
         SYNTAX ERROR("Variant part");
        end if:
                            -- if bypass(token semicolon)
       else
        SYNTAX ERROR("Variant part");
       end if:
                            -- if bypass(token case)
      else
       SYNTAX ERROR("Variant part");
      end if:
                            -- if bypass(token end)
    else
     SYNTAX ERROR("Variant part"):
                                    -- if variant statement
    SYNTAX ERROR("Variant part");
   end if;
                                    -- if bypass(token is)
   SYNTAX ERROR("Variant part");
  end if:
                                    -- if bypass(token identifier)
 else
  return (FALSE):
 end if:
                            -- if bypass(token case)
end VARIANT PART;
-- VARIANT -- > when CHOICE CHOICE * = > COMPONENT LIST
function VARIANT return boolean is
if (BYPASS(TOKEN WHEN)) then
  if (CHOICE) then
```

```
while (BYPASS(TOKEN BAR)) loop
    if not (CHOICE) then
      SYNTAX ERROR("Variant"):
     end if:
                                   -- if not choice statement
   end loop:
   if (BYPASS(TOKEN ARROW)) then
    if (COMPONENT LIST) then
      return (TRUE);
     else
      SYNTAX ERROR("Variant");
     end if:
                                   -- if component list statement
   else
    SYNTAX ERROR("Variant");
   end if:
                                   -- if bypass(token arrow)
   SYNTAX ERROR("Variant");
  end if:
                                   -- if choice statement
 else
  return (FALSE);
 end if:
                            -- if bypass(token when)
end VARIANT:
 -- WITH OR USE CLAUSE --> identifier i. identifier *;
function WITH OR USE CLAUSE return boolean is
 if (BYPASS(TOKEN IDENTIF'ER)) then
  while (BYPASS(TOKEN COMMA)) loop
   if not (BYPASS(TOKEN IDENTIFIER)) then
     SYNTAX ERROR("With or use clause"):
   end if:
  end loop:
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
   SYNTAX ERROR("With or use clause");
  end if:
                            -- if bypass(token semicolon)
 else
  return (FALSE):
                            -- if bypass(token identifier)
end WITH OR USE CLAUSE:
 -- FORMAL PART -- (PARAMETER SPECIFICATION : PARAMETER SPECIFICATION *)
function FORMAL PART return boolean is
 if (BYPASS(TOKEN LEFT PAREN)) then
    FORMAL PARAM DECLARE:
   if (PARAMETER SPECIFICATION) then
    while (BYPASS(TOKEN SEMICOLON)) loop
```

```
if not (PARAMETER SPECIFICATION) then
     SYNTAX ERROR("Formal part");
    end if;
                    -- if not parameter specification statement
   end loop;
   if (BYPASS(TOKEN RIGHT PAREN)) then
    if PACKAGE BODY DECLARE then
    WRITE HENRY DATA(BLANK, DUMMY LEXEME, END PARAM DECLARE,
              NONE, NEXT HEN);
    CREATE NODE(NEXT HEN, LAST RECORD);
    FORMAL PARAM DECLARE := FALSE;
    return (TRUE):
    SYNTAX ERROR("Formal part");
   end if:
                    -- if bypass(token right paren) statement
  else
   SYNTAX ERROR("Formal part");
  end if:
                    -- if parameter specification statement
 else
  return (FALSE):
                    -- if bypass(token left paren) statement
 end if:
end FORMAL PART:
 -- IDENTIFIER DECLARATION --> IDENTIFIER LIST: IDENTIFIER DECLARATION TAIL
function IDENTIFIER DECLARATION return boolean is
begin
put(RESULT_FILE, "IN IDENTIFIER DECLARATION"); NEW_LINE(RESULT_FILE);
 HENRY WRITE ENABLE := TRUE;
 IDENT DECLARE := TRUE;
 if (IDENTIFIER LIST) then
  if (BYPASS(TOKEN COLON)) then
   if (IDENTIFIER DECLARATION TAIL) then
    HENRY WRITE ENABLE := FALSE:
    return (TRUE):
    SYNTAX ERROR("Identifier declaration");
                     -- if identifier list statement
   end if:
  else
   SYNTAX ERROR("Identifier declaration"):
  end if:
                     -- if bypass(token colon)
 PISP
  return(FALSE);
                     -- if identifier list statement
 end if
end IDENTIFIER DECLARATION:
 -- IDENTIFIER DECLARATION TAIL -- exception EXCEPTION TAIL
                           -- constant CONSTANT TERM
```

```
array ARRAY TYPE DEFINITION
                                     EXPRESSION ? ;
                             --> NAME IDENTIFIER TAIL
function IDENTIFIER DECLARATION TAIL return boolean is
put(RESULT_FILE, "IN IDENTIFIER DECLARATION TAIL"): NEW_LINE(RESULT_FILE);
 if (BYPASS(TOKEN EXCEPTION)) then
  if (EXCEPTION TAIL) then
   return (TRUE):
  else
   SYNTAX ERROR("Identifier declaration tail");
                         -- if exception tail statement
 elsif (BYPASS(TOKEN CONSTANT)) then
  if (CONSTANT TERM) then
   return (TRUE):
  else
   SYNTAX ERROR("Identifier declaration tail");
                                    -- if constant term statement
 elsif (BYPASS(TOKEN ARRAY)) then
  if (ARRAY TYPE DEFINITION) then
   if (BYPASS(TOKEN ASSIGNMENT)) then
     if (EXPRESSION) then
      null:
     else
      SYNTAX ERROR("Identifier declaration tail"):
     end if:
                                    -- if expression statement
   end if:
                                    -- if bypass(token assignment)
  else
   SYNTAX ERROR("Identifier declaration tail");
                                    -- if array type definition
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
   SYNTAX ERROR("Identifier declaration tail"):
  end if:
                                   -- if bypass(token semicolon)
 elsif (NAME) then
  if (IDENTIFIER TAIL) then
   return (TRUE);
   SYNTAX ERROR("Identifier declaration tail"):
  end if:
                                   -- if identifier tail
 else
  return (FALSE);
                            -- if bypass(token exception)
 end if:
end IDENTIFIER DECLARATION TAIL:
 -- EXCEPTION TAIL --> :
                -- renames NAME:
function EXCEPTION TAIL return boolean is
```

begin

```
if (BYPASS(TOKEN SEMICOLON)) then
  return (TRUE):
 elsif (BYPASS(TOKEN RENAMES)) then
  if (NAME) then
    if (BYPASS(TOKEN SEMICOLON)) then
     return (TRUE);
    else
     SYNTAX ERROR("Exception tail");
    end if:
                           -- if bypass(token semicolon)
    SYNTAX ERROR("Exception tail");
  end if;
                           -- if name statement
 else
  return (FALSE);
 end if;
                           -- if bypass(token_semicolon)
end EXCEPTION TAIL:
 -- EXCEPTION CHOICE --> identifier
                 --> others
function EXCEPTION CHOICE return boolean is
begin
 if (BYPASS(TOKEN IDENTIFIER)) then
  return (TRUE):
 elsif (BYPASS(TOKEN OTHERS)) then
  return (TRUE):
 else
  return (FALSE);
 end if:
end EXCEPTION CHOICE:
 -- CONSTANT_TERM --> array ARRAY TYPE DEFINITION := EXPRESSION? :
               --> := EXPRESSION;
               --> NAME IDENTIFIER TAIL
function CONSTANT_TERM return boolean is
begin
 if (BYPASS(TOKEN ARRAY)) then
  if (ARRAY TYPE DEFINITION) then
   if (BYPASS(TOKEN ASSIGNMENT)) then
    if (EXPRESSION) then
     null;
    else
     SYNTAX ERROR("Constant term");
    end if:
                                  -- if expression statement
   end if:
                                  -- if bypass(token assignment)
  else
   SYNTAX ERROR("Constant term"):
                                  -- if array type definition
  if (BYPASS(TOKEN SEMICOLON)) then
```

```
return (TRUE);
    SYNTAX ERROR("Constant term"):
   end if:
                                     -- if bypass(token semicolon)
 elsif (BYPASS(TOKEN ASSIGNMENT)) then
   if (EXPRESSION) then
    if (BYPASS(TOKEN SEMICOLON)) then
     return (TRUE):
    else
     SYNTAX ERROR("Constant term");
    end if:
                                     -- if bypass(token semicolon)
   else
    SYNTAX ERROR("Constant term");
   end if;
                                     -- if expression statement
 elsif (NAME) then
  if (IDENTIFIER TAIL) then
    return (TRUE);
    SYNTAX ERROR("Constant term"):
  end if;
                                     -- if identifier tail statement
 else
  return (FALSE);
 end if:
                             -- if bypass(token array)
end CONSTANT TERM;
 -- IDENTIFIER TAIL --> CONSTRAINT? := EXPRESSION?;
                 --> renames NAME?;
function IDENTIFIER TAIL return boolean is
put(RESULT_FILE, "IN IDENTIFIER TAIL"); NEW_LINE(RESULT_FILE);
 if (CONSTRAINT) then
  null:
 end if:
                             -- if constraint statement
 if (BYPASS(TOKEN RENAMES)) then
  if (NAME) then
   null:
  else
   SYNTAX ERROR("Identifier tail"):
  end if:
                                    -- if name statement
 end if:
                             -- if bypass(token renames)
 if (BYPASS(TOKEN ASSIGNMENT)) then
  if (EXPRESSION) then
   null:
  else
   SYNTAX ERROR("Identifier tail"):
  end if:
                                    -- if expression statement
                             -- if bypass(token assignment)
if (BYPASS(TOKEN SEMICOLON)) then
  return (TRUE):
else
```

```
return (FALSE):
                          -- if bypass(token semicolon)
 end if:
end IDENTIFIER TAIL:
 -- PARAMETER SPECIFICATION --> IDENTIFIER LIST: MODE NAME (= EXPRESSION ?)
function PARAMETER SPECIFICATION return boolean is
put(RESULT_FILE, "IN PARAMETER SPECIFICATION"); NEW_LINE(RESULT_FILE);
 HENRY WRITE ENABLE := TRUE; --to capture first parameter
 if (IDENTIFIER LIST) then
  if (BYPASS(TOKEN COLON)) then
   if (MODE) then
    if (NAME) then
                                  -- check for type mark
      if (BYPASS(TOKEN ASSIGNMENT)) then
       if (EXPRESSION) then
        null:
       else
        SYNTAX ERROR("Parameter specification"):
       end if:
                          -- if expression statement
      end if:
                           -- if bypass(token assignment)
      return (TRUE);
      SYNTAX ERROR("Parameter specification");
                          -- if name statement
    SYNTAX ERROR("Parameter specification");
                           -- if mode statement
   end if:
  else
   SYNTAX ERROR("Parameter specification");
  end if:
                           -- if bypass(token colon)
  return (FALSE):
                           -- if identifier list statement
end PARAMETER SPECIFICATION:
-- IDENTIFIER LIST -- identifier identifier *
function IDENTIFIER LIST return boolean is
begin
put(RESULT_FILE, "IN IDENTIFIER LIST"); NEW_LINE(RESULT_FILE);
 if (BYPASS(TOKEN IDENTIFIER)) then
  if FORMAL PARAM DECLARE AND PACKAGE BODY DECLARE then
    WRITE HENRY DATA(BLANK, DUMMY LEXEME, PARAM TYPE, NONE, LAST RECORD);
  elsif (NOT PACKAGE BODY DECLARE) then
   WRITE HENRY DATA(LOCAL DECLARE, DUMMY LEXEME, IDENT TYPE.
              NONE, LAST RECORD):
  end if;
  while (BYPASS(TOKEN COMMA)) loop
   if (IDENT DECLARE) OR (FORMAL PARAM DECLARE AND PACKAGE BODY DECLARE)
```

```
then
    HENRY WRITE ENABLE := TRUE:
   if FORMA: PARAM DECLARE AND PACKAGE BODY DECLARE then
    WRITE HENRY DATA(BLANK, DUMMY LEXEME, PARAM TYPE,
              NONE, NEXT HEN);
   elsif (NOT FORMAL PARAM DECLARE) then
    WRITE HENRY DATA(LOCAL DECLARE, DUMMY LEXEME, IDENT TYPE,
             NONE, NEXT HEN);
   end if:
   if not (BYPASS(TOKEN IDENTIFIER)) then
    SYNTAX ERROR("Identifier list");
   end if:
                  -- if not bypass(token identifer) statement
  end loop;
  return (TRUE);
else
  return (FALSE);
end if:
                   -- if bypass(token identifier) statement
end IDENTIFIER LIST;
-- MODE --> in ?
      --> in out
       --> out
function MODE return boolean is
put(RESULT_FILE, "IN PARAMETER MODE"); NEW LINE(RESULT_FILE);
 if (BYPASS(TOKEN IN)) then
  if PACKAGE BODY DECLARE THEN
  WRITE HENRY DATA(BLANK, DUMMY LEXEME, PARAM TYPE, IN TYPE, LAST RECORD
  if (BYPASS(TOKEN OUT)) then
  if PACKAGE BODY DECLARE then
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, PARAM TYPE,
                IN OUT TYPE, LAST RECORD):
   end if:
  end if;
 elsif (BYPASS(TOKEN OUT)) then
 if PACKAGE BODY DECLARE then
  WRITE HENRY DATA(BLANK, DUMMY LEXEME, PARAM TYPE, OUT TYPE, LAST RECOR
 end if;
 end if:
 if (LAST_RECORD.TYPE_DEFINE = PARAM_TYPE)
  AND (LAST RECORD.PARAM TYPE = NONE) THEN
  WRITE HENRY DATA(BLANK, DUMMY LEXEME, PARAM TYPE, IN TYPE, LAST RECORD)
end if:
return (TRUE);
end MODE:
```

```
-- DESIGNATOR --> identifier
             --> string literal
function DESIGNATOR return boolean is
begin
 if (BYPASS(TOKEN IDENTIFIER)) then
  return (TRUE);
 elsif (BYPASS(TOKEN_STRING_LITERAL)) then
  return (TRUE);
 else
  return (FALSE);
 end if:
end DESIGNATOR;
 -- SIMPLE STATEMENT --> null;
                 --> ASSIGNMENT OR PROCEDURE CALL
                 --> exit EXIT STATEMENT
                 --> return RETURN STATEMENT
                 --> goto GOTO STATEMENT
                 --> delay DELAY STATEMENT
                 --> abort ABORT STATEMENT
                 --> raise RAISE STATEMENT
function SIMPLE STATEMENT return boolean is
begin
 if (BYPASS(TOKEN NULL)) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE):
  else
   SYNTAX ERROR("Simple statement"):
 elsif (ASSIGNMENT OR PROCEDURE CALL) then -- includes a check for a
  return (TRUE);
                                        -- code statement and an
                                           entry call statement.
 elsif (BYPASS(TOKEN EXIT)) then
  if (EXIT STATEMENT) then
   return (TRUE);
  else
   SYNTAX ERROR("Simple statement");
  end if;
 elsif (BYPASS(TOKEN RETURN)) then
  if (RETURN STATEMENT) then
   return (TRUE);
  else
   SYNTAX ERROR("Simple statement"):
 elsif (BYPASS(TOKEN GOTO)) then
  if (GOTO STATEMENT) then
   return (TRUE):
   SYNTAX ERROR("Simple statement");
  end if:
```

```
elsif (BYPASS(TOKEN DELAY)) then
  if (DELAY STATEMENT) then
   return (TRUE):
  else
   SYNTAX ERROR("Simple statement"):
  end if:
 elsif (BYPASS(TOKEN ABORT)) then
  if (ABORT STATEMENT) then
   return (TRUE);
  else
   SYNTAX ERROR("Simple statement");
  end if:
 elsif (BYPASS(TOKEN RAISE)) then
  if (RAISE STATEMENT) then
   return (TRUE):
  else
   SYNTAX ERROR("Simple statement");
  end if;
 else
  return (FALSE):
 end if:
end SIMPLE STATEMENT;
 -- ASSIGNMENT OR PROCEDURE CALL --> NAME := EXPRESSION :
                          --> NAME;
function ASSIGNMENT OR PROCEDURE CALL return boolean is
ASSIGN POINTER, FUNCALL POINTER:
begin
put(result file, "in assign or procedure call"); new line(result file);
 HENRY WRITE ENABLE := TRUE;
 ASSIGN POINTER := NEXT HEN:
 if (NAME) then
  if (BYPASS(TOKEN ASSIGNMENT)) then
   ASSIGN STATEMENT := TRUE;
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, ASSIGN TYPE,
            NONE, NEXT HEN);
   CREATE NODE(NEXT HEN, LAST RECORD):
   if NAME TAIL SET then
    WRITE HENRY DATA(BLANK, DUMMY LEXEME, PROCALL OR DS,
              NONE, ASSIGN POINTER):
  end if:
  FUNCALL POINTER : " NEXT HEN:
  HENRY WRITE ENABLE := TRUE.
   if (EXPRESSION) then
    if (BYPASS(TOKEN SEMICOLON)) then
      NAME TAIL SET : FALSE;
     ASSIGN STATEMENT : FALSE:
     WRITE HENRY DATA(BLANK, DUMMY LEXEME, END ASSIGN TYPE.
```

```
NONE, NEXT HEN);
      CREATE NODE(NEXT HEN, LAST RECORD);
      HENRY WRITE ENABLE := FALSE:
      return (TRUE);
                            -- parsed an assignment statement
      SYNTAX ERROR("Assignment or procedure call");
     end if:
                            -- if bypass(token semicolon)
     SYNTAX ERROR("Assignment or procedure call");
                            -- if expression statement
   elsif (BYPASS(TOKEN SEMICOLON)) then
    WRITE HENRY DATA(BLANK, DUMMY LEXEME, PROCALL OR DS,
              NONE, ASSIGN POINTER);
    CREATE NODE(NEXT HEN, LAST RECORD):
    return (TRUE);
                                   -- parsed a procedure call statement
    SYNTAX ERROR("Assignment or procedure call");
   end if:
                           -- if bypass(token assignment)
 else
   return (FALSE):
                            -- if name statement
end ASSIGNMENT OR PROCEDURE CALL:
 -- LABEL --> << identifier >>
function LABEL return boolean is
 if (BYPASS(TOKEN LEFT BRACKET)) then
  if (BYPASS(TOKEN IDENTIFIER)) then
    if (BYPASS(TOKEN RIGHT BRACKET)) then
     return (TRUE);
   else
     SYNTAX ERROR("Label");
                                  -- if bypass(token right_bracket)
  else
   SYNTAX ERROR("Label"):
  end if:
                                  -- if bypasss(token identifier)
 else
  return (FALSE);
 end if:
                           -- if bypass(token left bracket)
end LABEL;
 -- ENTRY DECLARATION --> entry identifier (DISCRETE RANGE)?
                         FORMAL PART?;
function ENTRY PECLARATION return boolean is
 if (BYPASS(TOKEN ENTRY)) then
  if (BYPASS(TOKEN IDENTIFIER)) then
   if (BYPASS(TOKEN LEFT PAREN)) then
```

```
if (DISCRETE RANGE) then
      if (BYPASS(TOKEN RIGHT PAREN)) then
       null;
      else
       SYNTAX ERROR("Entry declaration");
      end if;
                            -- if bypass(token right paren)
    else
      SYNTAX ERROR("Entry declaration");
    end if:
                                   -- if discrete range statement
   end if:
                                   -- if bypass(token left paren)
   if (FORMAL PART) then
    null;
   end if:
                                   -- if formal part statement
   if (BYPASS(TOKEN SEMICOLON)) then
    return (TRUE);
    SYNTAX ERROR("Entry declaration");
                                   -- if bypass(token_semicolon)
   end if:
  else
   SYNTAX ERROR("Entry declaration");
  end if:
                                    -- if bypass(token identifier)
 else
  return (FALSE):
 end if:
                            -- if bypass(token entry)
end ENTRY DECLARATION;
 -- REPRESENTATION CLAUSE --> for NAME use record RECORD REPRESENTATION CLAUSE
                       --> for NAME use at ? SIMPLE EXPRESSION;
function REPRESENTATION CLAUSE return boolean is
 if (BYPASS(TOKEN FOR)) then
  if (NAME) then
    if (BYPASS(TOKEN USE)) then
     if (BYPASS(TOKEN RECORD STRUCTURE)) then
      if (RECORD REPRESENTATION CLAUSE) then
       return (TRUE);
      else
       SYNTAX ERROR("Representation clause");
                            -- if record representation clause
     elsif (BYPASS(TOKEN AT)) then
      if (SIMPLE EXPRESSION) then
       if (BYPASS(TOKEN SEMICOLON)) then
        return (TRUE);
        SYNTAX ERROR("Representation clause"):
       end if:
                            -- if bypass(token semicolon)
      else
       SYNTAX ERROR("Representation clause");
                             -- if simple expression statement
     elsif (SIMPLE EXPRESSION) then
```

```
if (BYPASS(TOKEN SEMICOLON)) then
       return (TRUE);
      else
       SYNTAX ERROR("Representation clause"):
                            -- if bypass(token semicolon)
      end if:
    else
      SYNTAX ERROR("Representation clause");
    end if:
                                    -- if bypass(token record)
    SYNTAX ERROR("Representation clause");
   end if;
                                    -- if bypass(token use)
   SYNTAX ERROR("Representation clause");
  end if;
                                    -- if name statement
 else
  return (FALSE);
 end if:
                            -- if bypass(token for)
end REPRESENTATION CLAUSE:
 -- RECORD REPRESENTATION CLAUSE --> at mod SIMPLE EXPRESSION?
                                  NAME at SIMPLE EXPRESSION range RANGES *
                                 end record;
function RECORD REPRESENTATION CLAUSE return boolean is
 if (BYPASS(TOKEN AT)) then
  if (BYPASS(TOKEN MOD)) then
   if (SIMPLE EXPRESSION) then
     null;
   else
     SYNTAX ERROR("Record representation clause");
   end if;
                                    -- if simple expression
   SYNTAX ERROR("Record representation clause");
  end if;
                                    -- if bypass(token mod)
 end if:
                            -- if bypass(token at)
 while (NAME) loop
  if (BYPASS(TOKEN AT)) then
   if (SIMPLE EXPRESSION) then
     if (BYPASS(TOKEN RANGE)) then
      if (RANGES) then
       null:
      else
       SYNTAX ERROR("Record representation clause"):
                            -- if ranges statement
      SYNTAX ERROR("Record representation clause"):
     end if:
                                    -- if bypass(token range)
     SYNTAX ERROR("Record representation clause"):
                                    -- if simple expression
    end if:
```

```
else
   SYNTAX ERROR("Record representation clause"):
                                 -- if bypass(token at)
  end if:
 end loop:
 if (BYPASS(TOKEN END)) then
  if (BYPASS(TOKEN RECORD STRUCTURE)) then
   if (BYPASS(TOKEN SEMICOLON)) then
    return (TRUE);
    SYNTAX ERROR("Record representation clause");
                                  -- if bypass(token semicolon)
   end if:
  else
   SYNTAX ERROR("Record representation clause"):
  end if:
                                  -- if bypass(token record structure)
 else
  return (FALSE);
 end if:
                                  -- if bypass(token end)
end RECORD REPRESENTATION CLAUSE;
end PARSER 2;
  TITLE:
              AN ADA SOFTWARE METRIC
-- MODULE NAME: PACKAGE PARSER 3
-- DATE CREATED: 22 JUL 86
  LAST MODIFIED: 30 MAY 87
-- AUTHORS:
                 LCDR JEFFREY L. NIEDER
           LT KARL S. FAIRBANKS, JR.
            LCDR PAUL M. HERZIG
-- DESCRIPTION: This package contains thirty-five functions
     that make up the baseline productions for our top-down,
     recursive descent parser. Each function is preceded
     by the grammar productions they are implementing.
with PARSER 4. HENRY GLOBAL, HENRY, BYPASS FUNCTION, HALSTEAD METRIC,
   GLOBAL PARSER, GLOBAL, TEXT IO:
use PARSER 4. HENRY GLOBAL, HENRY, BYPASS FUNCTION, HALSTEAD METRIC.
  GLOBAL PARSER, GLOBAL, TEXT 10:
package PARSER 3 is
 function SUBTYPE INDICATION return boolean:
 function ARRAY TYPE DEFINITION return boolean:
 function CHOICE return boolean:
 function ITERATION SCHEME return boolean:
 function LOOP PARAMETER SPECIFICATION return boolean;
```

```
function EXPRESSION return boolean:
 function RELATION return boolean:
 function RELATION TAIL return boolean;
 function SIMPLE EXPRESSION return boolean:
 function SIMPLE EXPRESSION TAIL return boolean;
 function TERM return boolean;
 function FACTOR return boolean;
 function PRIMARY return boolean:
 function CONSTRAINT return boolean;
 function FLOATING OR FIXED POINT CONSTRAINT return boolean;
 function INDEX CONSTRAINT return boolean:
 function RANGES return boolean:
 function AGGREGATE return boolean;
 function COMPONENT ASSOCIATION return boolean;
 function ALLOCATOR return boolean;
 function NAME return boolean;
 function NAME TAIL return boolean;
 function LEFT PAREN NAME TAIL return boolean;
 function ATTRIBUTE DESIGNATOR return boolean:
 function INTEGER TYPE DEFINITION return boolean:
 function DISCRETE RANGE return boolean;
 function EXIT STATEMENT return boolean:
 function RETURN STATEMENT return boolean:
 function GOTO STATEMENT return boolean;
 function DELAY STATEMENT return boolean:
 function ABORT STATEMENT return boolean:
 function RAISE STATEMENT return boolean;
end PARSER 3;
  -----
package body PARSER 3 is
 -- SUBTYPE INDICATION --> NAME CONSTRAINT?
function SUBTYPE INDICATION return boolean is
begin
 if (NAME) then
                                 -- check for type mark
  if (CONSTRAINT) then
   null:
  end if:
  return (TRUE):
  return (FALSE);
 end if:
end SUBTYPE INDICATION:
 -- ARRAY TYPE DEFINITION -- (INDEX CONSTRAINT of SUBTYPE INDICATION
 -- this function parses both constrained and unconstrained arrays
function ARRAY TYPE DEFINITION return boolean is
```

```
begin
 if (BYPASS(TOKEN LEFT PAREN)) then
   if (INDEX CONSTRAINT) then
    if (BYPASS(TOKEN OF)) then
     if (SUBTYPE INDICATION) then
      return (TRUE);
     else
      SYNTAX ERROR("Array definition"):
                          -- if subtype indication
    else
     SYNTAX ERROR("Array definition");
                          -- if bypass(token of)
    SYNTAX ERROR("Array definition");
   end if:
                          -- if index constraint statement
 else
   return (FALSE);
 end if.
                             -- if bypass(token left paren)
end ARRAY TYPE DEFINITION:
 -- CHOICE --> EXPRESSION .. SIMPLE EXPRESSION?
          --> EXPRESSION CONSTRAINT?
         --> others
function CHOICE return boolean is
begin
 if (EXPRESSION) then
  if (BYPASS(TOKEN RANGE DOTS)) then -- check for discrete range
   if (SIMPLE EXPRESSION) then
     null:
   else
     SYNTAX ERROR("Choice");
                                   -- if simple expression statement
  elsif (CONSTRAINT) then
   null:
  end if:
                                   -- if bypass token range dots
  return (TRUE);
 elsif (BYPASS(TOKEN OTHERS)) then
  return (TRUE);
 else
  return (FALSE):
 end if:
end CHOICE:
 -- ITERATION SCHEME --> while EXPRESSION
                 -- · for LOOP PARAMETER SPECIFICATION
function ITERATION SCHEME return boolean is
 if (BYPASS(TOKEN WHILE)) then
```

```
NESTING METRIC(WHILE CONSTRUCT):
  if (EXPRESSION) then
   return (TRUE):
   SYNTAX ERROR("Iteration scheme"):
  end if:
 elsif (BYPASS(TOKEN FOR)) then
  NESTING METRIC(FOR CONSTRUCT):
  if (LOOP PARAMETER SPECIFICATION) then
   return (TRUE);
   SYNTAX ERROR("Iteration scheme");
  end if:
 else
  return (FALSE):
 end if:
end ITERATION SCHEME:
 -- LOOP PARAMETER SPECIFICATION --> identifier in reverse? DISCRETE RANGE
function LOOP PARAMETER SPECIFICATION return boolean is
 if (BYPASS(TOKEN IDENTIFIER)) then
  if (BYPASS(TOKEN IN)) then
   if (BYPASS(TOKEN REVERSE)) then
    null:
   end if:
                        -- if bypass(token reverse)
   if (DISCRETE RANGE) then
    return (TRUE);
    SYNTAX ERROR("Loop parameter specification");
                         -- if discrete range statement
   end if:
   SYNTAX ERROR("Loop parameter specification"):
  end if:
                        -- if bypass(token in)
 else
  return (FALSE):
                        -- if bypass(token identifier)
end LOOP PARAMETER SPECIFICATION.
 -- EXPRESSION --> RELATION RELATION TAIL?
function EXPRESSION return boolean is
begin
 if (RELATION) then
  if (RELATION TAIL) then
   null:
                       -- if relation tail statement
  end if:
  return (TRUE);
 else
```

```
return (FALSE):
 end if:
                        -- if relation statement
end EXPRESSION:
 -- RELATION -- - SIMPLE EXPRESSION SIMPLE EXPRESSION TAIL ?
function RELATION return boolean is
 if (SIMPLE EXPRESSION) then
  if (SIMPLE EXPRESSION TAIL) then
   null:
  end if:
                        -- if simple expression tail statement
  return (TRUE);
 else
  return (FALSE):
 end if:
                        -- if simple expression statement
end RELATION:
 -- RELATION TAIL --> [and [then ?] RELATION]*
                --> or else? RELATION*
                --> xor RELATION*
function RELATION TAIL return boolean is
begin
 while (BYPASS(TOKEN AND)) loop
  if (BYPASS(TOKEN THEN)) then
   null:
  end if:
                                    -- if bypass(token then)
  if not (RELATION) then
   SYNTAX ERROR("Relation tail");
  end if:
                                    -- if not relation statement
 end loop;
 while (BYPAS3(TOKEN OR)) loop
  if (BYPASS(TOKEN ELSE)) then
   null;
  end if:
                                    -- if bypass(token else)
  if not (RELATION) then
   SYNTAX ERROR("Relation tail");
  end if:
                                    -- if not relation statement
 end loop;
 while (BYPASS(TOKEN XOR)) loop
  if not (RELATION) than
   SYNTAX ERROR("Relation tail"):
  end if:
                                    -- if not relation statement
 end loop;
 return (TRUE):
end RELATION TAIL:
```

```
-- SIMPLE EXPRESSION --> - ? TERM BINARY ADDING OPERATOR TERM *
                  --> -? TERM BINARY ADDING OPERATOR TERM *
function SIMPLE EXPRESSION return boolean is
 if (BYPASS(TOKEN PLUS) or BYPASS(TOKEN MINUS)) then
  if (TERM) then
   while (BINARY ADDING OPERATOR) loop
    if not (TERM) then
      SYNTAX ERROR("Simple expression"):
                        -- if not term statement
    end if;
   end loop;
   return (TRUE):
   SYNTAX ERROR("Simple expression");
  end if:
                        -- if term statement
 elsif (TERM) then
  while (BINARY ADDING OPERATOR) loop
   if not (TERM) then
    SYNTAX ERROR("Simple expression");
   end if:
                        -- if not term statement
  end loop:
  return (TRUE):
  return (FALSE):
 end if:
                        -- if bypass(token plus) et al statement
end SIMPLE EXPRESSION:
 -- SIMPLE EXPRESSION TAIL --> RELATIONAL OPERATOR SIMPLE EXPRESSION
                       --> [not ?] in RANGES
                       --> not? in NAME
function SIMPLE EXPRESSION TAIL return boolean is
 if (RELATIONAL OPERATOR) then
  if (SIMPLE EXPRESSION) then
   return (TRUE):
   SYNTAX ERROR("Simple expression tail");
                            -- if simple expression statement
 elsif (BYPASS(TOKEN NOT)) then
  if (BYPASS(TOKEN IN)) then
   if (RANGES) then
     return (TRUE):
   elsif (NAME) then
                                   -- check for type mark
    return (TRUE):
    SYNTAX ERROR("Simple expression tail"):
                            -- if ranges statement
   end if:
  else
   SYNTAX ERROR("Simple expression tail"):
  end if:
                            -- if bypass(token in) statement
```

```
elsif (BYPASS(TOKEN IN)) then
  if (RANGES) then
   return (TRUE);
  elsif (NAME) then
                                    -- check for type mark
   return (TRUE);
  else
   SYNTAX ERROR("Simple expression tail");
                            -- if ranges statement
 else
  return (FALSE):
 end if:
                            -- if relational operator statement
end SIMPLE EXPRESSION TAIL:
 -- TERM --> FACTOR MULTIPLYING OPERATOR FACTOR *
function TERM return boolean is
begin
 if (FACTOR) then
  while (MULTIPLYING OPERATOR) loop
   if not (FACTOR) then
    SYNTAX ERROR("Term");
   end if:
                       -- if not factor statement
  end loop:
  return (TRUE):
 else
  return (FALSE):
 end if:
                       -- if factor statement
end TERM:
 -- FACTOR --> PRIMARY ** PRIMARY ?
         -- > abs PRIMARY
         --> not PRIMARY
function FACTOR return boolean is
begin
 if (PRIMARY) then
  if (BYPASS(TOKEN EXPONENT)) then
   if (PRIMARY) then
     null:
   else
     SYNTAX ERROR("Factor"):
   end if:
                       -- if primary statement
  end if:
                       -- if bypass(token exponent) statement
  return (TRUE):
 elsif (BYPASS(TOKEN ABSOLUTE)) then
  if (PRIMARY) then
   return (TRUE);
   SYNTAX ERROR("Factor"):
  end if:
                       -- if primary(abs) statement
```

```
elsif (BYPASS(TOKEN NOT)) then
  if (PRIMARY) then
   return (TRUE);
   SYNTAX ERROR("Factor"):
  end if:
                      -- if primary(not) statement
 else
  return (FALSE):
 end if:
                      -- if primary statement
end FACTOR:
 -- PRIMARY --> numeric literal
         --> null
         --> string literal
         --> new ALLOCATOR
         --> NAME
       --> AGGREGATE
function PRIMARY return boolean is
begin
 HENRY WRITE ENABLE := TRUE;
 if (BYPASS(TOKEN NUMERIC LITERAL)) then
  WRITE HENRY DATA(BLANK, DUMMY LEXEME, IDENT TYPE, NONE, LAST RECORD):
  return (TRUE):
 elsif (BYPASS(TOKEN NULL)) then
  return (TRUE):
 elsif (BYPASS(TOKEN STRING LITERAL)) then
  WRITE HENRY DATA(BLANK, DUMMY LEXEME, IDENT TYPE, NONE, LAST RECORD):
  return (TRUE):
 elsif (BYPASS(TOKEN NEW)) then
  if (ALLOCATOR) then
   return (TRUE);
  else
   SYNTAX ERROR("Primary");
  end if:
                                 -- if allocator statement
 elsif (NAME) then
  return (TRUE):
 elsif (AGGREGATE) then
  return (TRUE):
 else
  return (FALSE):
 end if:
                        -- if bypass(token left paren)
end PRIMARY:
 -- CONSTRAINT --> range RANGES
             --- digits FLOATING OR FIXED POINT CONSTRAINT
             -- > delta FLOATING OR FIXED POINT CONSTRAINT
             -- (INDEX CONSTRAINT
```

```
function CONSTRAINT return boolean is
begin
 if (BYPASS(TOKEN RANGE)) then
  if (RANGES) then
   return (TRUE);
  elsif (BYPASS(TOKEN BRACKETS)) then
                                              -- check for <> when parsing
   return (TRUE);
                                        -- an unconstrained array
  else
   SYNTAX ERROR("Constraint"):
  end if.
                                        -- if ranges statement
 elsif (BYPASS(TOKEN DIGITS)) or else (BYPASS(TOKEN DELTA)) then
  if (FLOATING OR FIXED POINT CONSTRAINT) then
   return (TRUE);
  else
   SYNTAX ERROR("Constraint"):
  end if:
 elsif (BYPASS(TOKEN LEFT PAREN)) then
  if (INDEX CONSTRAINT) then
   return (TRUE);
  else
   SYNTAX ERROR("Constraint"):
  end if:
 else
  return (FALSE):
 end if:
end CONSTRAINT:
       -- FLOATING OR FIXED POINT CONSTRAINT --> SIMPLE EXPRESSION range RANGES?
function FLOATING OR FIXED POINT CONSTRAINT return boolean is
begin
 if (SIMPLE EXPRESSION) then
  if (BYPASS(TOKEN RANGE)) then
   if (RANGES) then
    null;
    SYNTAX ERROR("Floating or fixed point constraint"):
   end if:
                                 -- if ranges statement
  end if:
                                 -- if bypass(token range)
  return (TRUE):
 else
  return (FALSE):
                          -- if simple expression statement
end FLOATING OR FIXED POINT CONSTRAINT:
 -- INDEX CONSTRAINT -- DISCRETE RANGE DISCRETE RANGE * )
function INDEX CONSTRAINT return boolean is
 if (DISCRETE RANGE) then
```

```
while (BYPASS(TOKEN COMMA)) loop
    if not (DISCRETE RANGE) then
     SYNTAX ERROR("Index constraint");
    end if;
                                   -- if not discrete range
   end loop:
   if (BYPASS(TOKEN RIGHT PAREN)) then
    return (TRUE);
  else
    SYNTAX ERROR("Index constraint"):
                                   -- if bypass(token right paren)
 else
  return (FALSE);
 end if:
                            -- if discrete range statement
end INDEX CONSTRAINT;
 -- RANGES --> SIMPLE EXPRESSION SIMPLE EXPRESSION?
function RANGES return boolean is
begin
 if (SIMPLE EXPRESSION) then
  if (BYPASS(TOKEN RANGE DOTS)) then
   if (SIMPLE EXPRESSION) then
     null:
    SYNTAX ERROR("Ranges"):
   end if:
                            -- if simple expression statement
  end if:
                                   -- if bypass(token range dots)
  return (TRUE):
  return (FALSE):
 end if:
                            -- if simple expression statement
end RANGES:
 -- AGGREGATE --> (COMPONENT ASSOCIATION , COMPONENT ASSOCIATION*)
function AGGREGATE return boolean is
begin
 if (BYPASS(TOKEN LEFT PAREN)) then
  if (COMPONENT ASSOCIATION) then
   while (BYPASS(TOKEN COMMA)) loop
    if not (COMPONENT ASSOCIATION) then
      SYNTAX ERROR("Aggregate");
    end if:
                                   -- if not component association
   end loop:
   if (BYPASS(TOKEN RIGHT PAREN)) then
    return (TRUE):
    SYNTAX ERROR("Aggregate"):
   end if:
                                   -- if bypass(token right paren)
  Plan
```

```
SYNTAX ERROR("Aggregate");
  end if:
                                   -- if component association statement
 else
  return (FALSE):
 end if:
                            -- if bypass(token left paren)
end AGGREGATE;
 -- COMPONENT ASSOCIATION --> CHOICE CHOICE * -- ? EXPRESSION
function COMPONENT ASSOCIATION return boolean is
begin
 if (CHOICE) then
  while (BYPASS(TOKEN BAR)) loop
   if not (CHOICE) then
    SYNTAX ERROR("Component association");
   end if:
  end loop:
  if (BYPASS(TOKEN ARROW)) then
   if (EXPRESSION) then
    null:
   else
    SYNTAX ERROR("Component association");
   end if:
                                   -- if expression statement
  end if:
                                   -- if bypass(token arrow)
  return (TRUE):
 else
  return (FALSE):
                            -- if choice statement
end COMPONENT ASSOCIATION:
 -- ALLOCATOR -- SUBTYPE INDICATION 'AGGREGATE?
function ALLOCATOR return boolean is
begin
 if (SUBTYPE INDICATION) then
  if (BYPASS(TOKEN APOSTROPHE)) then
   if (AGGREGATE) then
    null:
   0/40
     SYNTAX ERROR("Allocator").
   end if:
                                   -- if aggregate statement
  end if:
                                   -- if bypass(token apostrophe)
  return (TRUE):
  return (FALSE):
                            - if subtype indication statement
 end if.
end ALLOCATOR:
```

```
-- NAME -- > identifier NAME TAIL?
       --> character literal NAME TAIL?
       --> string literal NAME TAIL ?
function NAME return boolean is
begin
put(result file. "in name"); new line(result file);
 if (BYPASS(TOKEN IDENTIFIER)) then
   NAME POINTER := LAST RECORD:
  if (NAME TAIL) then
    null:
  end if:
  return (TRUE):
 HENRY WRITE ENABLE := TRUE:
 elsif (BYPASS(TOKEN CHARACTER LITERAL)) then
  if (NAME TAIL) then
   null:
  end if;
  return (TRUE):
 elsif (BYPASS(TOKEN STRING LITERAL)) then
  if (NAME TAIL) then
   null;
  end if:
  return (TRUE):
  return (FALSE);
 end if:
end NAME:
 -- NAME TAIL --> (LEFT PAREN NAME TAIL
          --> SELECTOR NAME TAIL *
          --> 'AGGREGATE NAME TAIL'*
          --> 'ATTRIBUTE DESIGNATOR NAME TAIL *
function NAME TAIL return boolean is
put(result file, "in name tail"); new line(result file);
 if (BYPASS(TOKEN LEFT PAREN)) then
   NAME TAIL SET := TRUE:
   HENRY WRITE ENABLE : TRUE:
   if ASSIGN STATEMENT then
   WRITE HENRY DATA/BLANK, DUMMY LEXEME, FUNCALL OR DS,
             NONE, NAME POINTER):
  else WRITE HENRY DATA(BLANK, DUMMY LEXEME, PROCALL OR DS.
               NONE. NAME POINTER):
  end if:
  if (LEFT PAREN NAME TAIL) then
   return (TRUE):
  else
   return (FALSE);
```

```
-- if left paren name tail
  end if:
 elsif (BYPASS(TOKEN PERIOD)) then
  if (SELECTOR) then
   while (NAME TAIL) loop
    null:
   end loop;
   return (TRUE);
  else
   SYNTAX ERROR("Name tail");
  end if:
                     -- if selector statement
 elsif (BYPASS(TOKEN APOSTROPHE)) then
  if (AGGREGATE) then
   while (NAME TAIL) loop
    null:
   end loop:
   return (TRUE);
  elsif (ATTRIBUTE DESIGNATOR) then
   while (NAME TAIL) loop
    null:
   end loop:
   return (TRUE):
   SYNTAX ERROR("Name tail");
  end if:
                     -- if aggregate statement
 else
  return (FALSE);
 end if:
                     -- if bypass(token left paren)
end NAME TAIL:
 -- LEFT PAREN NAME TAIL --> FORMAL PARAMETER? EXPRESSION EXPRESSION?
                       FORMAL PARAMETER? EXPRESSION EXPRESSION? *
                 ) NAME TAIL *
function LEFT PAREN NAME TAIL return boolean is
put(result file, "in left paren name tail"); new line(result file);
 if (FORMAL PARAMETER) then
                                         -- check for optional formal parameter
  null:
                            -- before the actual parameter
 end if;
                            -- if formal parameter statement
 HENRY WRITE ENABLE := TRUE:
 if (EXPRESSION) then
  if NAME TAIL SET then
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, PARAM TYPE, ACTUAL PARAM.
             LAST RECORD):
  end if:
  if (BYPASS(TOKEN RANGE DOTS)) then
   if not (EXPRESSION) then
    SYNTAX ERROR("Left paren name tail"):
   end if:
                                   -- if not expression statement
                                   -- if bypass(token range dots)
  while (BYPASS(TOKEN COMMA)) loop
```

```
if (FORMAL PARAMETER) then
    null:
   end if:
                                 -- if formal parameter statement
  HENRY WRITE ENABLE := TRUE;
   if not (EXPRESSION) then
    SYNTAX ERROR("Left paren name tail");
                                 -- if not expression statement
   if (BYPASS(TOKEN RANGE DOTS)) then
    if not (EXPRESSION) then
     SYNTAX ERROR("Left paren name tail");
                                 -- if not expression statement
    end if;
   end if:
                                 -- if bypass(token_range_dots)
   if NAME TAIL SET then
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, PARAM TYPE, ACTUAL PARAM.
             LAST RECORD);
   end if:
  end loop;
  if (BYPASS(TOKEN RIGHT PAREN)) then
   WRITE HENRY DATA(BLANK, DUMMY LEXEME, END ACTUAL PARAM,
             ACTUAL PARAM, NEXT HEN);
   CREATE NODE(NEXT HEN, LAST RECORD);
   NAME TAIL SET := FALSE:
   while (NAME TAIL) loop
    null:
   end loop;
   return (TRUE):
  else
   return (FALSE):
  end if:
                      -- if bypass(token right paren)
 elsif (DISCRETE RANGE) then
   if (BYPASS(TOKEN RIGHT PAREN)) then
    while (NAME TAIL) LOOP
     NULL:
    END LOOP:
    RETURN (TRUE);
    SYNTAX ERROR("Left paren name tail");
   end if:
 else
   return (FALSE);
  end if:
                      -- if bypass(token right paren)
end LEFT PAREN NAME TAIL;
 -- ATTRIBUTE DESIGNATOR --> identifier (EXPRESSION)?
                     --> range (EXPRESSION)?
                     -- digits (EXPRESSION)?
                     --> delta (EXPRESSION)?
function ATTRIBUTE DESIGNATOR return boolean is
 if (BYPASS(TOKEN IDENTIFIER)) or else (BYPASS(TOKEN RANGE)) then
```

Property of the second second

```
if (BYPASS(TOKEN LEFT PAREN)) then
   if (EXPRESSION) then
     if (BYPASS(TOKEN RIGHT PAREN)) then
     else
      SYNTAX ERROR("Attribute designator");
     end if;
                        -- if bypass(token right paren) statement
     SYNTAX ERROR("Attribute designator");
   end if:
                       -- if expression statement
  end if:
                       -- if bypass(token_left paren) statement
  return (TRUE);
 elsif (BYPASS(TOKEN DIGITS)) or else (BYPASS(TOKEN DELTA)) then
  if (BYPASS(TOKEN LEFT PAREN)) then
   if (EXPRESSION) then
     if (BYPASS(TOKEN RIGHT PAREN)) then
     else
      SYNTAX ERROR("Attribute designator");
     end if;
                        - .1 bypass(token right paren) statement
     SYNTAX ERROR("Attribute designator");
   end if;
                        -- if expression statement
  end if:
                        -- if bypass(token left paren) statement
  return (TRUE);
 else
  return (FALSE);
                        -- if bypass(token identifier) statement
end ATTRIBUTE DESIGNATOR:
 -- INTEGER TYPE DEFINITION --> range RANGES
function INTEGER TYPE DEFINITION return boolean is
 if (BYPASS(TOKEN RANGE)) then
  if (RANGES) then
   return (TRUE):
  else
   SYNTAX ERROR("Integer type definition");
  end if:
  return (FALSE):
 end if:
and INTEGER TYPE DEFINITION:
 -- DISCRETE RANGE --> RANGES CONSTRAINT ?
function DISCRETE RANGE return boolean is
 if (RANGES) then
```

```
if (CONSTRAINT) then
    null:
  end if:
                                   -- if constraint statement
   return (TRUE):
 else
  return (FALSE):
                            -- if ranges statement
end DISCRETE RANGE:
 -- EXIT STATEMENT --> NAME? when EXPRESSION?
function EXIT STATEMENT return boolean is
begin
 if (NAME) then
  null;
 end if:
                            -- if name statement
 if (BYPASS(TOKEN WHEN)) then
  if (EXPRESSION) then
    null:
    SYNTAX ERROR("Exit statement"):
  end if:
                                   -- if expression statement
                         -- if bypass(token when)
 if (BYPASS(TOKEN SEMICOLON)) then
  return (TRUE);
 else
  return (FALSE):
                         -- if bypass(token semicolon)
end EXIT STATEMENT:
 -- RETURN STATEMENT --> EXPRESSION?
function RETURN STATEMENT return boolean is
 if (EXPRESSION) then
  null:
 if (BYPASS(TOKEN SEMICOLON)) then
  return (TRUE);
 else
  return (FALSE):
end RETURN STATEMENT:
 -- GOTO STATEMENT -- · NAME:
function GOTO STATEMENT return boolean is
begin
 if (NAME) then
```

```
if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE):
  else
   SYNTAX ERROR("Goto statement");
  end if:
                           -- if bypass(token semicolon)
 else
  return (FALSE);
 end if:
                           -- if name statement
end GOTO STATEMENT;
 -- DELAY STATEMENT --> SIMPLE EXPRESSION:
function DELAY STATEMENT return boolean is
begin
 if (SIMPLE EXPRESSION) then
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE);
  else
   SYNTAX ERROR("Delay statement");
                           -- if bypass(token semicolon)
  end if:
 else
  return (FALSE);
 end if:
                           -- if simple expression statement
end DELAY STATEMENT;
 -- ABORT STATEMENT --> NAME : NAME *:
function ABORT STATEMENT return boolean is
begin
 if (NAME) then
  while (BYPASS(TOKEN COMMA)) loop
   if not (NAME) then
     SYNTAX ERROR("Abort statement"):
   end if;
                            -- if not name statement
  end loop:
  if (BYPASS(TOKEN SEMICOLON)) then
   return (TRUE):
   SYNTAX ERROR("Abort statement"):
  end if:
                            -- if bypass(token semicolon)
  return (FALSE);
                            -- if name statement
end ABORT STATEMENT.
 -- RAISE STATEMENT -- · NAME ? ..
function RAISE STATEMENT return boolean is
begin
```

```
if (NAME) then
  null:
 end if:
 if (BYPASS(TOKEN SEMICOLON)) then
  return (TRUE):
 else
  return (FALSE);
 end if:
end RAISE STATEMENT:
end PARSER 3:
  TITLE:
              AN ADA SOFTWARE METRIC
-- MODULE NAME: PACKAGE PARSER 4
-- DATE CREATED: 23 JUL 86
  LAST MODIFIED: 30 MAY 87
  AUTHORS:
                LCDR JEFFREY L. NIEDER
           LT KARL S. FAIRBANKS, JR.
           LCDR PAUL M. HERZIG
-- DESCRIPTION: This package contains seven functions that
     are the lowest level productions for our top-down,
     recursive descent parser. Each function is preceded
     by the grammar productions they are implementing.
with BYPASS FUNCTION, BYPASS SUPPORT FUNCTIONS, GLOBAL PARSER, GLOBAL, TEXT 1
USE BYPASS FUNCTION, BYPASS SUPPORT FUNCTIONS, GLOBAL PARSER, GLOBAL, TEXT IC
package PARSER 4 is
 function MULTIPLYING OPERATOR return boolean.
 function BINARY ADDING OPERATOR return boolean:
 function RELATIONAL OPERATOR return boolean;
 function ENUMERATION TYPE DEFINITION return boolean;
 function ENUMERATION LITERAL return boolean;
 function FORMAL PARAMETER return boolean;
 function SELECTOR return boolean:
end PARSER 4;
    .....
package body PARSER 4 is
```

```
-- MULTIPLYING OPERATOR --> *
                     --> mod
                     --> rem
function MULTIPLYING OPERATOR return boolean is
put(RESULT_FILE, "In multiplying operator"); new line(RESULT_FILE);
 if (BYPASS(TOKEN ASTERISK)) then
  return (TRUE):
 elsif (BYPASS(TOKEN SLASH)) then
  return (TRUE):
 elsif (BYPASS(TOKEN MOD)) then
  return (TRUE):
 elsif (BYPASS(TOKEN REM)) then
  return (TRUE);
 else
  return (FALSE):
 end if:
end MULTIPLYING OPERATOR:
 -- BINARY ADDING OPERATOR --> -
                      -- > &:
function BINARY ADDING OPERATOR return boolean is
put(RESULT_FILE, "In binary adding operator"); new line(RESULT_FILE);
 if (BYPASS(TOKEN PLUS)) then
  return (TRUE):
 elsif (BYPASS(TOKEN MINUS)) then
  return (TRUE):
 elsif (BYPASS(TOKEN AMPERSAND)) then
  return (TRUE):
  return (FALSE):
end BINARY ADDING OPERATOR:
 -- RELATIONAL OPERATOR --
function RELATIONAL OPERATOR return boolean is
put(RESULT_FILE, "In relational operator"); new line(RESULT_FILE);
 if (BYPASS(TOKEN EQUALS)) then
  return (TRUE):
```

```
elsif (BYPASS(TOKEN NOT EQUALS)) then
  return (TRUE):
 elsif (BYPASS(TOKEN LESS THAN)) then
  return (TRUE):
 elsif (BYPASS(TOKEN LESS THAN EQUALS)) then
  return (TRUE):
 elsif (BYPASS(TOKEN GREATER THAN)) then
  return (TRUE):
 elsif (BYPASS(TOKEN GREATER THAN EQUALS)) then
  return (TRUE):
  return (FALSE):
 end if:
end RELATIONAL OPERATOR:
 -- ENUMERATION TYPE DEFINITION --> (ENUMERATION LITERAL
                                   . ENUMERATION LITERAL *)
function ENUMERATION TYPE DEFINITION return boolean is
put(RESULT FILE, "In enumeration type definition"); new line(RESULT FILE);
 if (BYPASS(TOKEN LEFT PAREN)) then
  HENRY WRITE \bar{E}NABL\bar{E} := TRUE:
  if (ENUMERATION LITERAL) then
   while (BYPASS(TOKEN COMMA)) loop
   HENRY WRITE ENABLE := TRUE:
    if not (ENUMERATION LITERAL) then
      SYNTAX ERROR("Enumeration type definition"):
                         -- if not enumeration literal
    end if:
   end loop:
   if (BYPASS(TOKEN RIGHT PAREN)) then
    return (TRUE):
    SYNTAX ERROR("Enumeration type definition"):
   end if:
                          -- if bypass(token right paren)
  else
   SYNTAX ERROR("Enumeration type definition");
  end if:
                          -- if enumeration literal statement
 else
  return (FALSE):
                          -- if bypass(token left paren)
 end if:
end ENUMERATION TYPE DEFINITION:
 -- ENUMERATION LITERAL -- identifier
                     -- character literal
function ENUMERATION LITERAL return boolean is
put(RESULT_FILE, "In enumeration_literal"); new_line(RESULT_FILE);
 if (BYPASS(TOKEN IDENTIFIER)) then
```

```
return (TRUE):
 elsif (BYPASS(TOKEN CHARACTER LITERAL)) then
  return (TRUE);
 else
  return (FALSE);
 end if:
end ENUMERATION LITERAL;
 -- FORMAL PARAMETER --> identifier =>
function FORMAL PARAMETER return boolean is
put(RESULT FILE, "In formal parameter"); new line(RESULT FILE);
 LOOK AHEAD TOKEN := TOKEN RECORD BUFFER(TOKEN ARRAY INDEX - 1):
 if (ADJUST LEXEME(LOOK AHEAD TOKEN LEXEME,
         LOOK AHEAD TOKEN.LEXEME SIZE - 1) = "=>") then
  if (BYPASS(TOKEN IDENTIFIER)) then
   if (BYPASS(TOKEN ARROW)) then
    return (TRUE);
   else
    SYNTAX ERROR("Formal parameter");
   end if:
                                  -- if bypass(token arrow)
   SYNTAX ERROR("Formal parameter");
  end if:
                                  -- if bypass(token identifier)
 else
  return (FALSE):
 end if:
end FORMAL PARAMETER;
 -- SELECTOR --> identifier
          --> character literal
          --> string literal
          --> all
function SELECTOR return boolean is
put(RESULT_FILE, "In selector"): new line(RESULT_FILE):
 if (BYPASS(TOKEN IDENTIFIER)) then
  return (TRUE):
 elsif (BYPASS(TOKEN CHARACTER LITERAL)) then
  return (TRUE):
 elsif (BYPASS(TOKEN STRING LITERAL)) then
  return (TRUE):
```

```
elsif (BYPASS(TOKEN_ALL)) then return (TRUE); else return (FALSE); end if; end SELECTOR; end PARSER_4;
```

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